

## Menu

- Complexity Classes P and NP
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## Smileys Problem

Input: $n$ square tiles Output: Arrangement of the tiles in a square, where the colors and shapes match up, or "no, its impossible". "Not possible"

## How much work is the

 Smiley's Problem?- Upper bound: ( $O$ )
$O(n!)$ : Try all possible permutations
- Lower bound: ( $\Omega$ )
$\Omega(n)$ : Must at least look at every tile
- Tight bound: ( $\theta$ )

No one knows!

## NP Problems

- Can be solved by just trying all possible answers until we find one that is right
- Easy to quickly check if an answer is right - Checking an answer is in $\mathbf{P}$
- The smileys problem is in NP

We can easily try $n$ ! different answers We can quickly check if a guess is correct (check all $n$ tiles)

## Complexity Classes

Class $\mathbf{P}$ : problems that can be solved in polynomial time ( $O\left(n^{k}\right)$ for some constant $k$ ): Easy problems like simulating the universe are all in $\mathbf{P}$.
Class NP: problems that can be solved in polynomial time by a nondeterministic machine: includes all problems in $\mathbf{P}$ and some problems possibly outside $\mathbf{P}$ like the Smileys puzzle

$$
P=N P ?
$$

- Is P different from NP: is there a problem in NP that is not also in $P$
- If there is one, there are infinitely many
- Is the "hardest" problem in NP also in P - If it is, then every problem in NP is also in $P$
- No one knows the answer!
- The most famous unsolved problem in computer science and math
- Listed first on Millennium Prize Problems



## Quiz Responses

- Partners for PS3
- Only 4 groups worked as partners
- All partner groups got Gold stars
- Only 8 out of 20 non-partner PS got Gold stars
- Your responses:
- Learn more working alone: 7
- Finish faster working alone: 5
- Couldn't find anyone to work with: 2
- Wanted to work with PS1/PS2 partner: 4
- Started too late: 5


## Quiz Responses

- How fast:
- Way too fast: 1
- Too fast: 11
- Just about right: 11
- Too slow: 1
- Exam 1:
- Very confident: 1
- Confident: 4

AC's review session Wednesday at 7

- Concerned: 12
- Worried: 5

My office hours:
Tuesday, 3:30-4:30
Thursday, 11:30am-12:30pm
others by email request

- Terrified: 3

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## Other Response

"I decided to skip this problem set, since the lowest grade gets dropped because of conflicts with other courses"
This is not the intent of (from the Syllabus):
"For almost all students. doina the problem sets will $b_{r}$ The main point of the problem sets is to learn G the material! (Note: student who answered as this was the only one to also select "Way Too ys u: Fast" for question 2. If you don't do the PS, is $w$ the course will definitely seem way too fast! I, the weighting that is best for you is used."


## Reductions

- Problem A reduces to Problem B if:
- There is a polynomial time function $f$ such that $A(x)=B(f(x))$
- To reduce "3SAT" to "Smiley":
- $A=$ 3SAT
- $B=$ Smiley
- $f=$ tile replacements
$\rightarrow$
The Real 3SAT Problem (also can be reduced to the Smileys Puzzle)
- If A reduces B that means solving A is no harder than solving $B$ since we can use a solution to $B$ with $f$ to solve $A$


## Propositional Grammar

Sentence ::= Clause
Sentence Rule: Evaluates to value of Clause
Clause ::= Clause ${ }_{1} \vee$ Clause $_{2}$
Or Rule: Evaluates to true if either clause is true
Clause ::= Clause $_{1} \wedge$ Clause $_{2}$
And Rule: Evaluates to true iff both clauses are true

Propositional Grammar

## Clause ::= СClause

Not Rule: Evaluates to the opposite value of clause ( $\neg$ true $\rightarrow$ false) Clause ::= ( Clause )
Group Rule: Evaluates to value of clause. Clause ::= Name
Name Rule: Evaluates to value associated with Name.

## The Satisfiability Problem (SAT)

- Input: a sentence in propositional grammar
- Output: Either a mapping from names to values that satisfies the input sentence or no way (meaning there is no possible assignment that satisfies the input sentence)

Proposition Example

Sentence ::= Clause
Clause $::=$ Clause $_{1} \vee$ Clause $_{2}$ Clause $::=$ Clause $_{1} \wedge$ Clause $_{2} \quad($ and $)$ Clause $::=\neg$ Clause $\quad$ (not) Clause ::= ( Clause )
Clause ::= Name
$a \vee(b \wedge c) \vee \neg \mathrm{b} \wedge \mathrm{c}$


## The 3SAT Problem

- Input: a sentence in propositional grammar, where each clause is a disjunction of 3 names which may be negated.
- Output: Either a mapping from names to values that satisfies the input sentence or no way (meaning there is no possible assignment that satisfies the input sentence)
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## 3SAT / SAT

Is 3SAT easier or harder than SAT?

It is definitely not harder than SAT, since all 3SAT problems are also SAT problems. Some SAT problems are not 3SAT problems.

## Charge

- Wednesday's class: recap in context of everything so far
- Friday: how Lorenz was really broken
- AC's exam review is Wednesday, 7pm

