

## Menu

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Are there any non-recursive natural languages? What would happen to a society that spoke one?

Not for humans at least.
They would run out of original things to say.
Chimps and Dolphins are able to learn nonrecursive "languages" (some linguists argue they are not really "languages"), but only humans can learn recursive languages.

## Language Elements

When learning a foreign language, which elements are hardest to learn?

- Primitives: lots of them, and hard to learn real meaning
- Means of Combination
- Complex, but, all natural languages have similar ones [Chomsky] SOV (45\% of all languages) Sentence $::=$ Subject Object Verb (Korean) SVO $(42 \%) \quad$ Sentence $::=$ Subject Verb Object
 OSV (<1\%): Schemish: Tobati (New Guinea)
Expression ::= (Verb Object)
- Means of Abstraction: few of these, but tricky to learn differences across languages

English: I, we
Tok Pisin (Papua New Guinea): mi (I), mitupela (he/she and I), mitripela Tok Pisin (Papua New Guinea): mi ( I , , mitupela (he/she and I), mitripela
(both of them and I), mipela (all of them and I), yumitupela (you and I), (buth of them and I), mipela (all of them and ), yumitupela

## Running out of Ideas

"Its all been said before."

Eventually true for a non-recursive language.

Never true for a recursive language.
There is always something original left to say!
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|  | Pages in Revised ${ }^{5}$ Report on the Algorithmic Language Scheme | Pages in C++ Language Specification (1998) |
| :---: | :---: | :---: |
| Primitives | Standard Procedures 18 <br> Primitive expressions 2 <br> Identifiers, numerals 1 | Standard Procedures 356 <br> Primitive expressions 30 <br> Identifiers, numerals 10 |
| Means of Combination | Expressions 2 <br> Program structure 2 | Expressions, Statements 197 <br> Program Structure 35 |
| Means of Abstraction | Definitions $1 / 2$ | Declarations, Classes 173 |
|  | 48 pages total (includes formal specification and examples) | 776 pages total (includes no formal specification or examples) |
| C++ Core language issues list has 529 items! |  |  |
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## Why don't we just program computers using English?

- Not concise enough

English:
To find the maximum of two numbers, compare them. If the first number is greater than the second number, the maximum is the first number. Otherwise, the maximum is the second number.

Scheme:
(define (max ab) (if (> ab) ab))
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## Why don't we just program computers using English?

- Mapping between surface forms and meanings are ambiguous and imprecise
Would you rather be paid biweekly or every week?

```
Dictionary
```



```
ooj,
1.) Happening every, wow weks.
```

The exact meaning(s) of every Scheme expression is determined by simple, unambiguous rules we will learn today (and refine later in the course).

## Why don't we just program computers using English?

- Limited means of abstraction

There are only a few pronouns: he, she, it, they, these, ... (English doesn't even have a gender-neutral pronoun for a person!)
Only Webster and Oxford can make up new ones.
define allows any programmer to make up as many pronouns as she wants, and use them to represent anything.
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## Essential Scheme

Expression ::= (Expression ${ }_{1}$ Expression*)
Expression ::= (if Expression ${ }_{1}$ Expression $_{2}$ Expression $_{3}$ )
Expression ::= (define name Expression)
Expression ::= Primitive
Primitive ::= number
Primitive ::= + |-|*| ...
Primitive ::= ...

Grammar is clear, just follow the replacement rules. But what does it all mean?


## Evaluation Rule 1: Primitives

If the expression is a primitive, it is self-evaluating.

```
    >2
    2
    > #t
    #t
> +
    #<primitive:+>
```


## Evaluation Rule 3: Application

3. If the expression is an application:
a) Evaluate all the subexpressions of the combination (in any order)
b) Apply the value of the first subexpression to the values of all the other subexpressions.
(expression expression $_{1}$ expression $_{2} \ldots$ )

## Expressions and Values

- (Almost) every expression has a value
- Have you seen any expressions that don't have values?
- When an expression with a value is evaluated, its value is produced

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## Evaluation Rule 2: Names

If the expression is a name, it evaluates to the value associated with that name.
$>$ (define two 2 )
$>$ two
2

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| :---: | :---: | :---: |

## Rules for Application

1. If the procedure to apply is a primitive, just do it.
2. If the procedure is a compound procedure, evaluate the body of the procedure with each formal parameter replaced by the corresponding actual argument expression value.



## You've Already Used Lambda!

(define (closer-color? sample color1 color2)
Expr)
is a shortcut for:
(define closer-color?
(lambda (sample color1 color2)
Expr))


## Evaluating Special Forms

- Eval 4-if. If the expression is
(if Expression Expression $_{1}$ Expression ${ }_{2}$ ) evaluate Expression. If it evaluates to \#f, the value of the if expression is the value of Expression $_{2}$. Otherwise, the value of the if expression is the value of Expression.
- Eval 4-lambda. Lambda expressions selfevaluate. (Do not do anything until it is applied.)


## More Special Forms

- Eval 4-define. If the expression is (define Name Expression)
associate the Expression with Name.
- Eval 4-begin. If the expression is (begin Expression ${ }_{0}$ Expression $_{1} \ldots$ Expression $_{k}$ )
evaluate all the sub-expressions. The value of the begin expression is the value of Expressionk.

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## Special Forms

Expression ::= (lambda (Parameters) Expression)
Eval 4-lambda. Lambda expressions self-evaluate.
Parameters ::=
Parameters::= Name Parameters
Expression ::= (define Name Expression)
Eval 4-define. If the expression is (define Name Expression) associate the Expression with Name.
Expression ::= (if Expression Expression $_{1}$ Expression ${ }_{2}$ )
Eval 4-if. Evaluate Expression. If it evaluates to \#f, the value of the if expression is the value of Expression . Otherwise, the value of $^{\text {. Ot }}$ the if expression is the value of Expression
Expression ::= (begin ExpressionList Expression)
Eval 4-begin. Evaluate all the sub-expressions. The value of the begin expression is the value of Expression.

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## Charge

- PS1 Due Wednesday
-Staffed Lab hours:
today, 3:30-5pm
Tuesday, 4-5:30pm
- Reading for Friday: SICP, 1.2
- Reading for Monday: GEB, Ch 5

