


Cryptography in World War II
 Jefferson Institute for Lifelong Learning at UVa
 Spring 2006 David Evans

Class 4: Modern Cryptography



<http://www.cs.virginia.edu/jillcrypto>

Menu


- Some loose ends on WWII
- Maurice Burnett

- Modern Cryptography
 - Modern symmetric ciphers
 - Public-key cryptosystems

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British Cipher Machine

- Design based on commercial Enigma
- 5 rotor wheels (instead of 3 in Enigma)
- Multiple rings per rotor
- Last 2 rotor wheels didn't rotate
- British attempted to break it (without success)



Typex Machine

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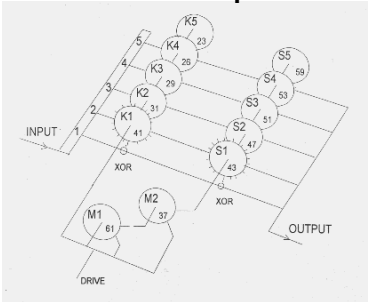
German Code-Breaking Efforts

- About 6,000 people (compare to 12,000 working at Bletchley Park)
- Decentralized: each military branch had their own, didn't share what they learned
- Effective against manual codes: broke about 50% of manually coded messages
- Didn't attempt to break rotor-based ciphers – so confident Enigma was unbreakable, didn't try to Typex and similar machines

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Lorenz Cipher



From <http://www.codesandciphers.org.uk/lorenz/fish.htm>

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Modern Symmetric Ciphers

A billion billion is a large number, but it's not that large a number.
Whitfield Diffie

- Same idea but:
 - Use digital logic instead of mechanical rotors
 - Larger keys (random bits, not rotor alignments)
 - Lorenz $\approx 5^{12} < 10^9$
 - Modern ≥ 128 bits $> 10^{37}$
 - Encrypt blocks of letters at a time

Modern Ciphers

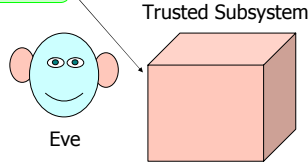
- AES (Rijndael) successor to DES selected 2001
- 128-bit keys, encrypt 128-bit blocks
- Brute force attack (around 10^{30} times harder than Lorenz)
 - Try 1 Trillion keys per second
 - Would take 10790283070806000000 years to try all keys!
 - If that's not enough, can use 256-bit key
- No known techniques that do better than brute force search

Login Process

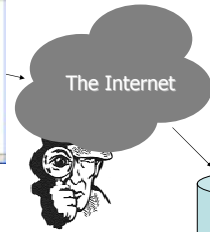
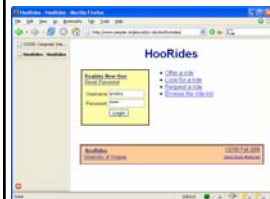
Terminal

Login: alyssa
Password: fido

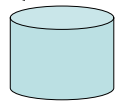
login sends
<"alyssa", "fido">



Sending Passwords

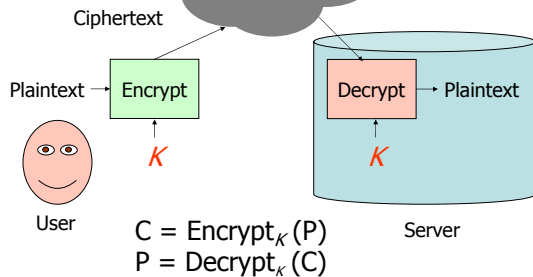


User



Server

The Internet



$$C = \text{Encrypt}_K(P)$$

$$P = \text{Decrypt}_K(C)$$

Key Agreement Demo

(Animated version at end of slides.)

Asymmetric Cryptosystems

- Need a hard problem (like symmetric cryptosystems)
- With a trap door: if you know a secret, the hard problem becomes easy

One-Way Functions

- Easy to compute, hard to invert
- Trap-door one way function:
 - $D(E(M)) = M$
 - E and D are easy to compute.
 - Revealing E doesn't reveal an easy way to compute D .
 - Hence, anyone who knows E can encrypt, but only someone who knows D can decrypt

RSA [Rivest, Shamir, Adelman 78]

One-way function:
multiplication is easy, factoring is hard
Trap-door: number theory (Euler and Fermat)



Security of RSA

- n is public, but not p and q where $n = pq$
- How much work is factoring n ?

Number Field Sieve (fastest known factoring algorithm) is:

$$O(e^{1.9223((\ln(n))^{1/3} (\ln(\ln(n)))^{2/3}))})$$

$n \sim 200$ digits – would take quintillions of years

Asymmetric Cryptosystems

- Encryption and Decryption are done with different keys
- Keep one of the keys secret, reveal the other

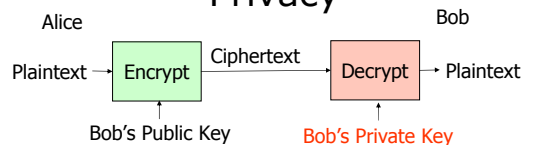
$$E_{KRA} (E_{KUA} (M)) = M$$

Alice's Public Key: KUA

Alice's Private Key: KRA

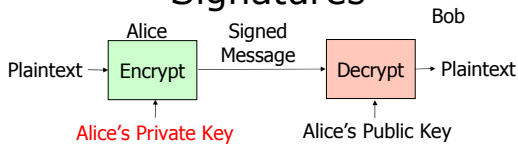
Only KRA can decrypt a message encrypted using KUA.

Public-Key Applications: Privacy

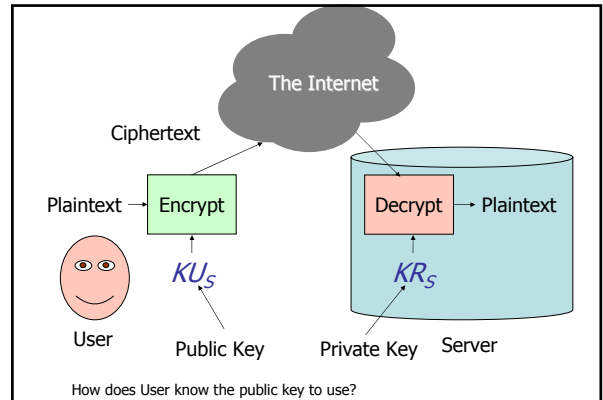


- Alice encrypts message to Bob using Bob's *Public Key*
- Only Bob knows Bob's *Private Key* \Rightarrow only Bob can decrypt message

Signatures



- Bob knows it was from Alice, since only Alice knows Alice's Private Key
- Non-repudiation: Alice can't deny signing message (except by claiming her key was stolen!)
- Integrity: Bob can't change message (doesn't know Alice's Private Key)



How does User know the public key to use?

Key Management

Approach 1: Meet Secretly

- User and Server Operator meet secretly and swap public keys
 - If you can do that, might as well agree on a secret (symmetric key) instead
 - Doesn't work for Internet transactions

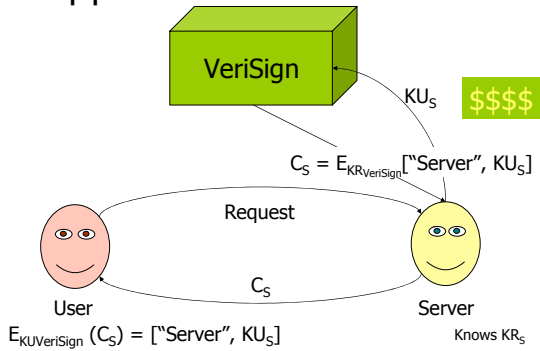
Approach 2: Public Announcement

- Publish public keys in a public forum
 - Append to email messages
 - Post on web site
 - New York Time classifieds
- Easy for rogue to pretend to be someone else
 - Forge email, alter web site, lie to New York Times

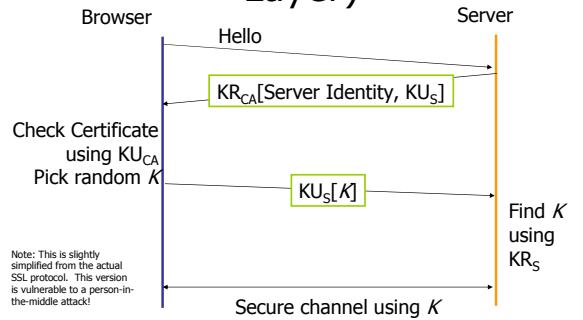
Approach 3: Public Directory

- Trusted authority maintains directory mapping names to public keys
- Entities register public keys with authority in some secure way
- Authority publishes directory
 - Print using watermarked paper, special fonts, etc.
 - Allow *secure* electronic access
 - Depends on secure distribution of directory's key

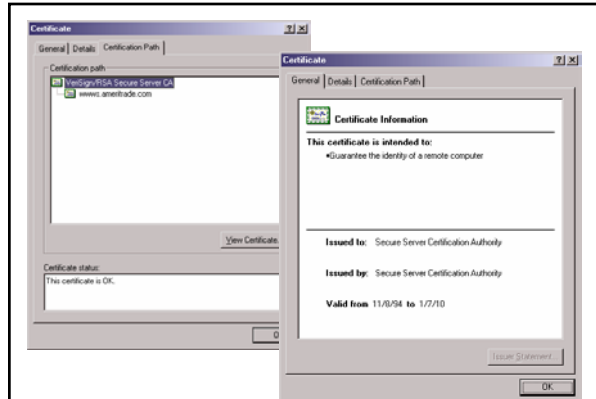
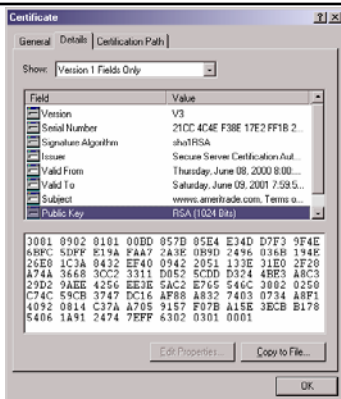
Approach 4: Certificates



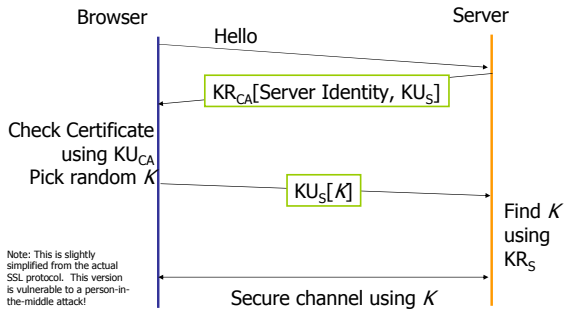
SSL (Secure Sockets Layer)



Data encrypted using secret key exchanged using some public key associated with some certificate.



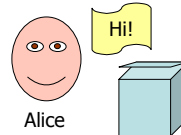
SSL Recap



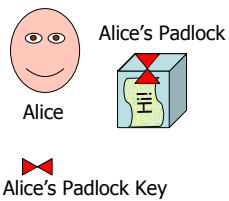
Questions?

Animated version of Asymmetric Cryptography Demo

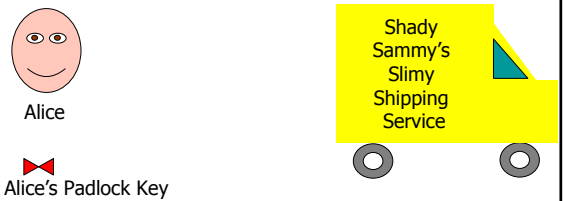
Padlocked Boxes



Padlocked Boxes



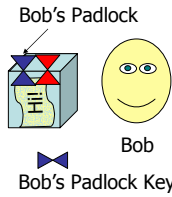
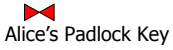
Padlocked Boxes



Padlocked Boxes



Alice

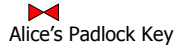


Bob

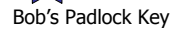
Padlocked Boxes



Alice



Bob



Padlocked Boxes



Alice



Bob



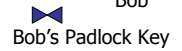
Padlocked Boxes



Alice



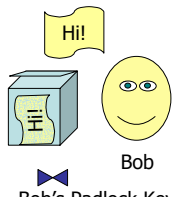
Bob



Padlocked Boxes



Alice



Bob