

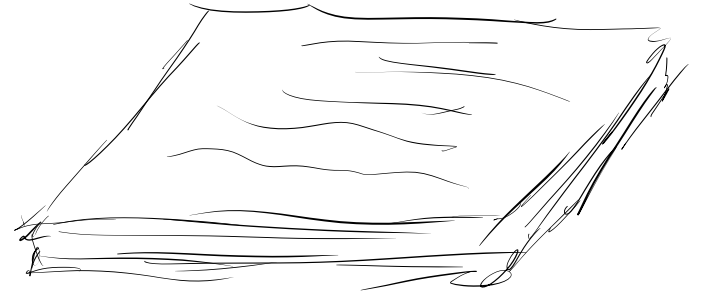
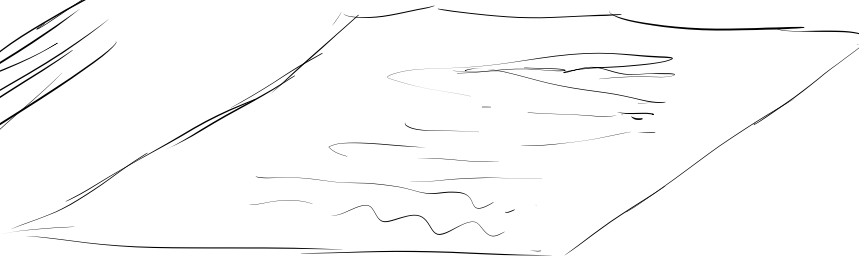
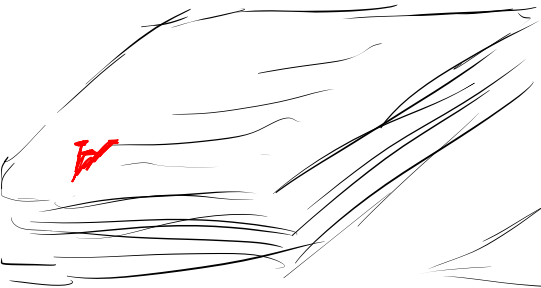
a. $\exists g \in \text{Grades}^M . \forall p \in \text{Papers}^P . E(p, g)$

b. $\forall p \in \text{Papers} . \exists g \in \text{Grades} . p \text{ gets } g$

M: grades
P: papers
E(x,y): x gets y

let $g = A$

rubber-stamping

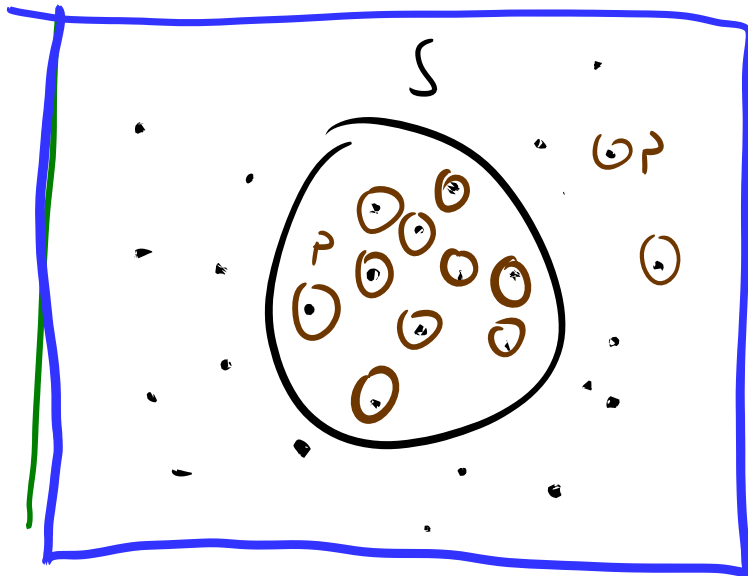
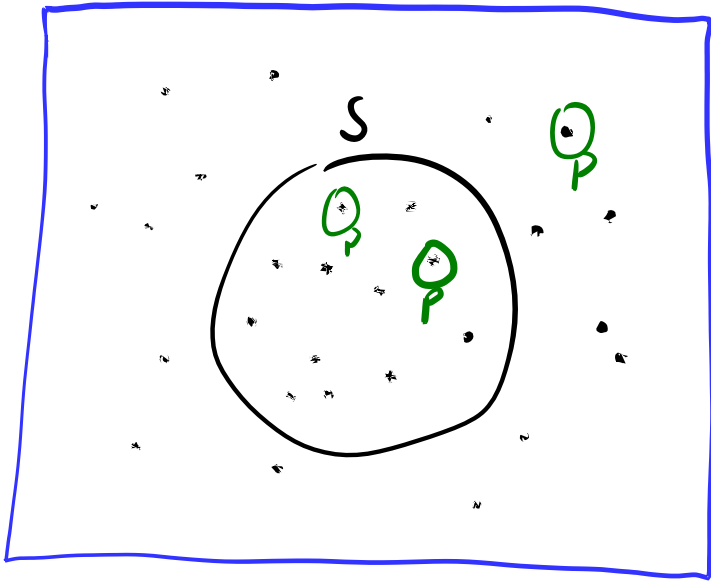




P: get a perfect score

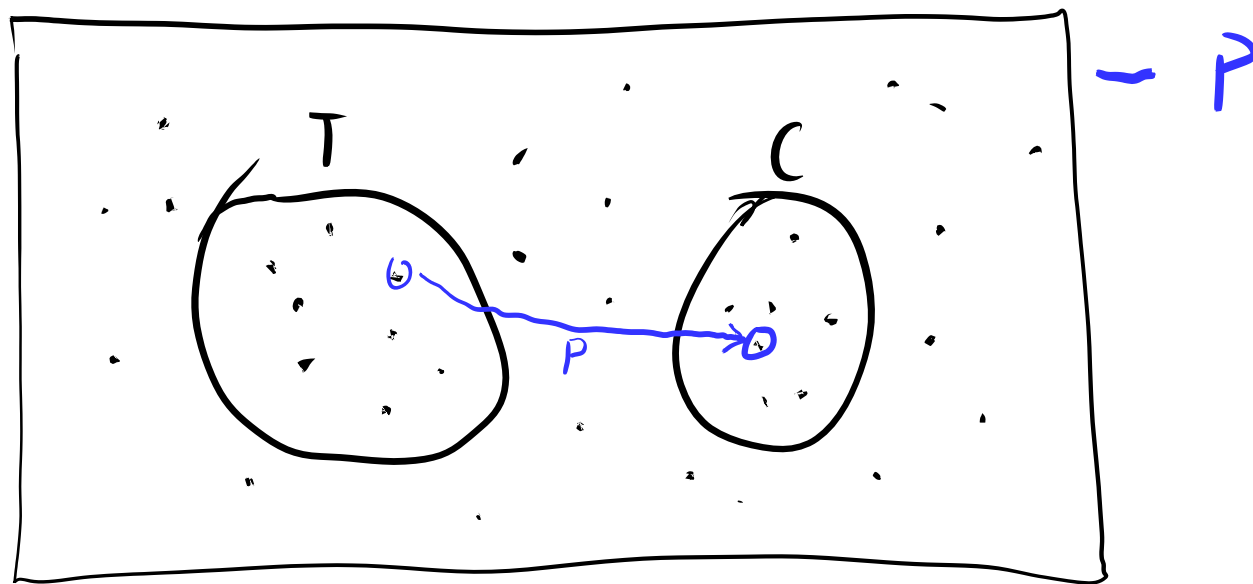
→ $\exists x \in S. P(x)$ — easy to prove : 1 example is enough

→ $\forall x \in S. P(x)$ ~ easy to disprove : 1 example



$$\exists x \in T. \exists y \in C. P(x, y)$$

T : set of all teachers
 C : set of candy types
 $P(x, y)$: x likes y



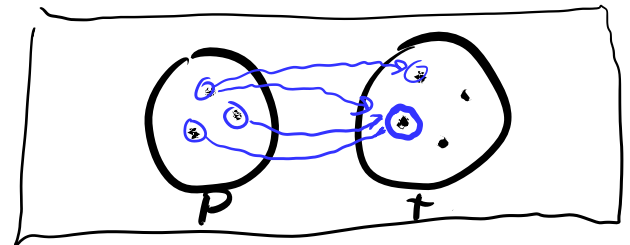
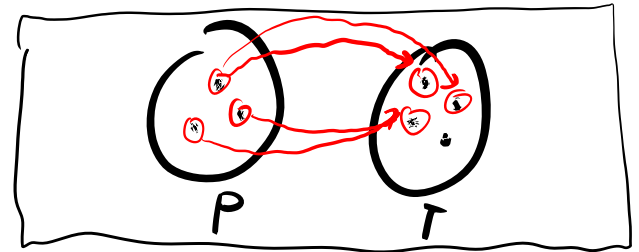
T : set of tests

P : set of programs

$M(x, y)$: program x ^{passes} meets test y

$\forall p \in P$. $\exists t \in T$. $M(p, t)$

$\exists t \in T$. $\forall p \in P$. $M(p, t)$



for any prog you pick,

I can find a test it meets

I can pick a test

such that no matter which prog
you pick,

the prog meets test

$b = F$
for each thing x
 if $P(x)$:
 $b = T$
return b

$b = T$
for each thing x
 if $\neg P(x)$:
 $b = F$
return b

$\exists x. P(x)$

$\forall x. P(x)$

$\exists x \in S. \forall y \in T. P(x, y)$

$b = F$
for each x in S

$c = T$
for each y in T
 if $\neg P(x, y)$:
 $c = F$

if c :
 $b = T$

return b

Entailment

$$X \models Y$$

$$X \rightarrow Y \equiv T$$

$$X \leftrightarrow Y \equiv T$$

$$X \equiv T$$

$$|S| > 0, \quad \forall x \in S. P(x) \models \exists x \in S. P(x)$$

$$\exists x \in S. \forall y \in T. P(x, y) \models \forall y \in T. \exists x \in S. P(x, y)$$