

## Introduction and Problem Statement

### Chemical-Mechanical Polishing (CMP)

- Rotating pad polishes each layer on wafers to achieve planarized surfaces
- Uneven features cause polishing pad to deform
- Density control is achieved by adding fill geometries into layout



#### Industry context: *fixed-dissection regime*:

- Density constraints imposed only for *fixed* set of  $w \times w$  windows
- Layout partitioned by  $r^2$  fixed dissections
- Each  $w \times w$  window is partitioned into  $r^2$  *tiles*

### The Filling Problem

- Given: design rule-correct layout in an  $n \times n$  layout region
- Find design rule-correct filled layout, such that:
  - No fill geometry is added within distance  $B$  of any layout feature, and:
  - Min-Var objective:** no fill is added into any window with density  $\geq U$ , and minimum window density in the filled layout is maximized, or:
  - Min-Fill objective:** number of filling features is minimized, and density of any window remains within given range  $(LB, UB)$

### Model for oxide planarization via CMP

$$z = \begin{cases} z_0 - \left( \frac{K_i t}{\rho(x, y)} \right) & t < (\rho_0 z_1) / K_i \\ z_0 - z_1 - K_i t + \rho_0(x, y) z_1 & t > (\rho_0 z_1) / K_i \end{cases}$$

Crucial model element: determining the effective initial pattern density  $\rho(x, y)$

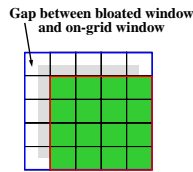
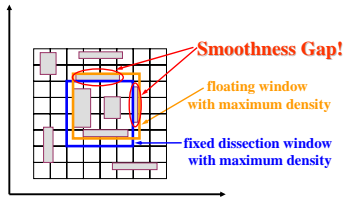
## Smoothness Gap & New Local Density Smoothness

#### Spatial Local Density:

$$\rho(W_{ij}) = \sum_{k=i}^{i+r-1} \sum_{l=j}^{j+r-1} \text{area}(T_{kl})$$

#### Effective Local Density:

$$\rho(W_{ij}) = \sum_{k=i}^{i+r-1} \sum_{l=j}^{j+r-1} \text{area}(T_{kl}) \cdot f(k - (i+r/2), l - (j+r/2))$$



### Accurate Layout Density Analysis

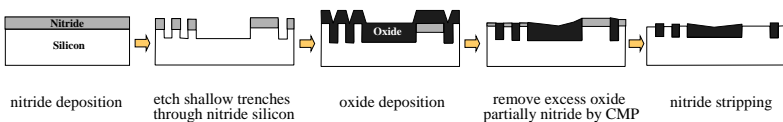
- Optimal extremal-density analysis ( $K^2$ )
  - Computational inefficiency!
- Multi-level density analysis algorithm:
  - Any window contained by *bloated* on-grid window
  - Any window contains *shrunk* on-grid window
  - Gap between max bloated and max on-grid window
    - Algorithmic inaccuracy!

### Three Types of Local Density Variation:

- Max density variation of every  $r$  neighboring windows in each fixed-dissection row
- Max density variation of every cluster of windows which cover one tile
- Max density variation of every cluster of windows which cover  $\frac{r}{2} \times \frac{r}{2}$  tiles

## STI Dual-Material Dummy Fill

### Shallow Trench Isolation (STI)



### STI CMP Model

- STI post-CMP variation:** controlled by changing the feature density distribution
- Compressible pad model:** polishing occurs on up/down areas after some step height
- Dual-material polish model:** two different materials for top & bottom surfaces

STI Fill is a non-linear programming problem!

### Previous methods

- Min-Var objective:** minimize max height variation
- Min-Fill objective:** minimize total inserted fills, while keeping given lower bound

### Drawbacks of previous work

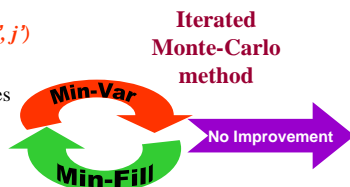
- Can not guarantee to find a global minimum since it is deterministic
- Simple termination is not sufficient to yield optimal/sub-optimal solutions

### Monte-Carlo method for STI Min-Var

- Calculate priority of tile  $(i, j)$  as  $\Delta H - \Delta H(i, j, i', j')$
- Pick the tile for next filling randomly
- If the tile is overfilled, lock all neighboring tiles
- Update tile priority

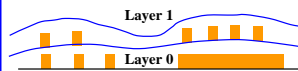
### MC/Greedy methods for STI Min-Fill

- Find a solution with **Min-Var** objective to satisfy the given lower bound
- Modify the solution with respect to **Min-Fill** objective



## Multiple-Layer Oxide CMP Dummy Fill

### Multiple-layer Oxide



### Multiple-layer density model

$$\hat{\rho}_{0(k)} = \begin{cases} [\hat{d}_k + \frac{z_{k-1}}{z_k} \hat{\rho}_{0(k-1)}] \times \hat{f} & k > 1 \\ \hat{d}_1 \times \hat{f} & k = 1 \end{cases}$$

### Multiple-layer Oxide Fill Objectives

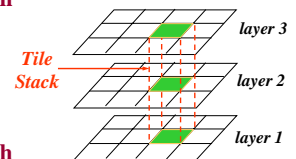
- Sum of density variations** can not guarantee the Min-Var objective on each layer
- Maximum density variation** across all layers

### Linear Programming formulation

- Min  $M$
- Subject to:  $0 \leq L_k \leq \rho_0(i, j, k) \leq H_k \leq \sum_{l=1}^k \left( \frac{z_l}{z_k} \right) i, j = 0, \dots, \frac{nr}{w} - 1, k = 1, \dots, K$
- $(H_k - L_k) \leq M$   $k = 1, \dots, K$
- $0 \leq x_{ijk} \leq \text{slack}(T_{ijk})$   $i, j = 0, \dots, \frac{nr}{w} - 1, k = 1, \dots, K$

### Multiple-Layer Monte-Carlo Approach

- Tile stack:** column of tiles
- Effective density of tile stack:** sum of effective densities of all tiles in stack



### Multiple-Layer Monte-Carlo Approach

- Compute slack area, cumulative effective density of tile stack
- Calculate tile stack's priority according to cumulative effective density
- While (sum of priorities  $> 0$ ) Do:
  - Randomly select a tile stack according to its priority
  - From bottom to top layer, check for fill insertion feasibility
  - Update slack area and priority of the tile stack
  - If no slack area is left, lock the tile stack