Moving-Target TSP and Related Problems



Computer Science University of Virginia

Charlottesville, Virginia 22904



C. S. Helvig Gabriel Robins Alexander Zelikovsky



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Introduction

- **Classical TSP:** sites to be visited are stationary
- Related work:
 - Time-dependent TSP (cost to travel between stationary sites changes with time)
- Our contribution:

Problem Formulation



- A set $S = \{s_1, s_2, ..., s_n\}$ of *targets*
- Each target s_i has constant velocity v_i
- Each target s_i starts moving from a position $p_i \in \Re^n$





- Moving-Target TSP formulation
- Algorithms for variants of Moving-Target TSP

Motivation:

- Supply ships resupply patrolling boats
- Planes intercept mobile ground units

Basic Observation

• Lemma: Optimal tours have no waiting periods



• **Corollary:** Pursuer always moves at max speed

• A *pursuer* starting at origin with maximum speed $v > |v_i|$

Find:

• A fastest tour which intercepts all targets

Variations (Moving-Target Vehicle Routing Problem):

- Multiple pursuers j = 1...k
- Pursuers have a given capacity to fill target demand

Few Moving Targets

• Few (O(log n/log log n)) moving targets \Rightarrow

- Efficient $(1+\alpha)$ -approximation algorithm
- α is performance bound of a heuristic for stationary TSP



One-Dimensional Variant

Trivial Algorithm:

- Intercept all targets on one side of the origin first
- Then intercept targets on the other side
- Does not work



Moving-Target TSP with Resupply

- Pursuer can intercept only one target before requiring resupply at the origin
 - Targets move directly away (or towards) the origin
- Corresponds to Moving-Target VRP except:
 - Single pursuer
 - Pursuer supply = target demand
- **Solution** (when targets move away from origin):
 - Intercept targets in order of increasing d_i/v_i
 - Algebraic proof for two targets
 - For n targets, swapping targets improves tour



- Targets intercept origin \Rightarrow Implicit change in target velocity
- Valid tour: No target passes through the origin
- **Theorem:** If the tour in which pursuer intercepts
 - Targets moving away in ascending order of d_i/v_i
 - Approaching targets in descending order of d_i/v_i
- is valid, then it is optimal

• **Problem:** Find the fastest *valid* tour

Theorem: The tour where the pursuer intercepts targets in descending order of d_i/v_i is always valid and the slowest

• **Theorem:** Slowest tour $\leq 2 \times (\text{optimal valid tour})$

- Lemma: Pursuer can change direction only after intercepting the fastest target
- Dynamic programming solution
- $O(n^2)$ algorithm
- Implemented and verified

Summary

- Formulation of Moving-Target TSP
- Exact algorithm for one-dimensional version
- Heuristic when few targets are moving

Approximate and exact heuristics for selected variants of Moving-Target TSP with Resupply

Multi-Pursuer Moving-Target TSP with Resupply

- Moving-Target Vehicle Routing Problem with
 - Multiple pursuers
 - Pursuer supply = target demand
 - All targets are moving away
- Minimization Objectives:
 - Makespan = when the last pursuer at the base
 - Standard for multiprocessor scheduling
 - NP-hard even for stationary targets
 - **Total time** = total time while pursuers are moving
 - Standard for classical vehicle routing problem
 - Trivial for stationary case
- **Theorem:** Total time objective is NP-hard

- If all targets have the same d_i/v_i • Total time objective is still NP-hard
 - Approximation algorithm with Performance ratio = max_i { $(1 + v_i) / (1 - v_i)$ }
- If all targets have the same speed v_i
 - Nontrivial result
 - The following strategy is optimal:
 - send next available pursuer after the closest target

Future Work

- 2-Dimensional Moving-Target TSP
- Moving-Target TSP with Resupply
- When targets can pass through origin or when finding the fastest valid tour: Prove/Disprove NP-Hardness
 - Find efficient algorithms/heuristics
- Multi-Pursuer Moving-Target TSP