

Caveat

You have three hours to complete the test. I most likely will be unavailable via email when you take test. Any appropriate question I do answer will be posted on the class website. The test is open book and class notes and closed neighbor. Computer access is limited to the class website. Every problem has a straightforward answer. Not all questions have the same point value.

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1. (5 points) The decision sorting problem determines whether its input is arranged in sorted order. Under what conditions is the decision sorting problem NP-hard (e.g., NP is a proper subset of P (not the right answer)).

2. (10 points) Suppose a tour for a Travelling Salesperson instance equals [1, 9, 3, 4, 7, 8, 5, 6, 2] (i.e., $tour[0] = 1$, $tour[1] = 9$, $tour[2] = 3$, ...). What is the tour after $interchange(2, 7)$ is executed, where the $interchange()$ method is the one described in assignment 8.

tour = [_____]

3. (15 points) The decision travelling salesperson problem takes as input a weighted graph and an amount k . The question under consideration is whether there is a tour whose length is at most k . Provide a *simple linear nondeterministic* algorithm that solves the problem, where the graph is represented using a two-dimensional array w (i.e., element $w[a][b]$ is the length of edge (a, b)).

```
DTSPSolver( int w[1..n][1..n], int k ) {  
    int[] t = new int array with n elements  
    // Attempt to assign t so that it represents a tour with the desired  
    // characteristics (i.e., a tour whose length  $\leq k$ . Depending on w it  
    // may or may not be successful
```

4. (15 points) Complete the following reduction for the CNF satisfiability problem for a formula F . The reduction is to the Halting problem. The reduction begins as follows.

- Define algorithm $p()$ in the following manner.

```
p(CNF F) {
    // perform exhaustive search on assigning values to the variables
    // in F (there are at most  $2^n$  different assignments)
    Boolean found = false;
    for each possible assignment of 0 and 1 to the variables in F do
        a = the current assignment of 0s and 1s to the variables in F
        found = F(a); // F(a) is the value of F when its variables
                    // have the values in a
        if ( found ) {
            // found an assignment that makes F satisfiable
            break;
        }
    }
    // loop forever iff a satisfiable assignment was found
    while ( found ) {
        continue;
    }
}
```

- Show that $p()$ has infinite loop if and only if F is satisfiable.

5. (10 points) What is the importance of the Covey diagram for time management?

	Urgent	Non-urgent
Important		
Non-important		

6. (10 points) Explain why simulated annealing uses a Boltzmann-like criterion $e^{-\text{delta}/t}$ for determining whether to accept a mutation that worsens the solution. Your answer should take into account the roles of *delta* (the amount by which the solution is worsened) and the current temperature *t*.

7. (10 points) A mutation method allows simulated annealing to consider different solutions in its search for an optimal answer. For what follows, the problem of interest is 0/1 Knapsack.
- Suppose the mutation method operates by replacing one item with another. When starting with a random solution is there always a way to get to an optimal solution? Explain.

 - Suppose the mutation method operates by either removing or adding one item. When starting with a random solution is there always a way to get to an optimal solution? Explain.
8. (10 points) A typical simulated annealing starting temperature is $2s$, where s^2 is an estimate of the variance of the scores in the solution space. Explain why $2s$ is a reasonable starting temperature for accepting most mutations given a Boltzmann-like criterion. (FYI, $e^{-1/2} \approx 0.606$ and $e^{-1} \approx 0.367$).

