A PageRank Based Algorithm with Anti-Collocation Constraints for Virtual Machine Placement in Cloud Datacenters

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In a cloud datacenter, one wants to satisfy demands for varying types of virtual machines while utilizing as few physical machines as possible.



Problem Statement

Model both VMs and PMs as vectors, where each dimension is a different resource.

Example: Write machine resources as

$$VMs = \left\{ \begin{pmatrix} 1\\1\\2 \end{pmatrix}, \begin{pmatrix} 2\\3\\2 \end{pmatrix} \right\} \qquad PMs = \left\{ \begin{pmatrix} 4\\4\\4 \end{pmatrix} \right\}$$

Problem Statement

Model both VMs and PMs as vectors, where each dimension is a different resource.

Example: Write machine resources as $\begin{pmatrix} CPU \\ Memory \\ Disk \end{pmatrix}$

$$VMs = \left\{ \begin{pmatrix} 1\\1\\2 \end{pmatrix}, \begin{pmatrix} 2\\3\\2 \end{pmatrix} \right\}$$

$$PMs = \left\{ \begin{pmatrix} 4 \\ 4 \\ 4 \end{pmatrix} \right\}$$





Anti-Collocation Constraints

In addition, we want to satisfy sets of *anti-collocation constraints* on certain resources.

Solution

Treat anti-collocation constrained resources as multiple dimensions in the resource vector

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Vector Bin Packing

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Issue: Vector Bin Packing is an NP-Hard problem!

Vector Bin Packing - Differences

However, the VM packing problem deviates from standard vector bin packing in a few ways

- Don't know full set of vectors to pack at a single time
- ► Vectors are pulled from a fixed, known set we decide upon
- ► Allowed to permute some dimensions of the vector (CPU cores)

Try to maximize PM usage or minimize PM variance

- First Fit Assign VM to first PM it fits into
- First Fit Decreasing Sum Assign VM to the PM whose resource utilization is highest
- CompVM (Dot Product)
- Integer Linear Programming

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Do not account for resource imbalances



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Assigns VMs to complementary PMs, but is not aware of the set of VM types or the best final PM state

Try to maximize PM usage or minimize PM variance

- ► First Fit Assign VM to first PM it fits into
- ► First Fit Decreasing Sum Assign VM to the PM whose resource utilization is highest
- CompVM (Dot Product)
- ► Integer Linear Programming

Too slow for large problem instances

Desired Solution

Goal

Minimize the number of wasted resources in each dimension.

- ► Needs to be efficient at assignment time
- Capable of handling anti-collocation constraints
- Minimizes the total number of PMs used



Profile Graph

Idea

Describe each possible state of a PM as a *profile*. Link these together into a graph based on possible transitions between states.



Let $PMs = \{[4, 4, 4, 4]\}.$











Profile Graph

Same profile!













Desired Metric Qualities



Desired Metric Qualities



Want to rank each profile based on

- how connected it is to higher-utilization profiles
- how flexible the profile is to accommodating various VM types

Desired Metric Qualities



Want to rank each profile based on

- how connected it is to higher-utilization profiles
- how flexible the profile is to accommodating various VM types

 \implies PageRank!

PageRank

Initialize PageRanks uniformly

$$PR(p_i;0)=\frac{1}{N}$$

At each time step update as



$$PR(p_i; t+1) = \frac{1-d}{N} + d\sum_{p_j \in M(p_i)} \frac{PR(p_j; t)}{L(p_j)}$$

PageRank







At assignment time

Out of your *active* PMs, assign the VM to the PM whose resulting profile has the highest PageRank.

Activate more PMs if current set of PMs gets too full



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Experiments

Simulation

- PM Usage
- VM Migrations
- Energy Consumption
- SLO Violations

Real Testbed

- ► PM Usage
- VM Migrations
- SLO Violations

Experiment

Simulation

Simulation performed on CloudSim using VM and PM types below.

VM Types	Virtual cores		Memory (GiB)	Virtual Disk	
	#	Speed (GHz)	Wiemory (OID)	#	Size (GB)
m3.medium	1	0.6	3.75	1	4
m3.large	2	0.6	7.5	1	32
m3.xlarge	4	0.6	15	2	40
m3.2xlarge	8	0.6	30	2	80
c3.large	2	0.7	3.75	2	16
c3.xlarge	4	0.7	7.5	2	40

TABLE I Description of VM types

TABLE II Description of PM types

PM Types	Physical Cores		Memory (GiB)	Physical Disk	
	#	Speed (GHz)	Wiemory (OID)	#	Size (GB)
M3	8	2.6	64	4	250
C3	8	2.8	7.5	4	250

Experiment

CloudSim

3000~VMs were allocated in batches of 1000. CPU utilization of VMs were given by two public traces. Results are averaged over 100 trials.



PM Usage





(a) PlanetLab



(b) Google Cluster

Summary

PageRankVM is a heuristic VM placement algorithm that:

- ► Reduces the total number of PMs needed to host VM demands
- ► Satisfies anti-collocation constraints by considering VM permutations
- Has very low placement-time cost

Questions

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