Collaborative bandit learning has become an emerging focus for personalized recommendation. It leverages user dependence for joint model estimation and recommendation. As such online learning solutions directly learn from users, e.g., result clicks, they bring in new challenges in privacy protection. Despite the existence of recent studies about privacy in contextual bandit algorithms, how to efficiently protect user privacy in a collaborative bandit learning environment remains unknown.

In this paper, we develop a general solution framework to achieve differential privacy in collaborative bandit algorithms, under the notion of global differential privacy and local differential privacy. The key idea is to inject noise in a bandit model’s sufficient statistics (either on server side to achieve global differential privacy or client side to achieve local differential privacy) and calibrate the noise scale with respect to the structure of collaboration among users. We study two popularly used collaborative bandit algorithms to illustrate the application of our solution framework. Theoretical analysis proves our derived private algorithms reduce the added regret caused by privacy-preserving mechanism compared to its linear bandits counterparts, i.e., collaboration actually helps to achieve stronger privacy with the same amount of injected noise. We also empirically evaluate the algorithms on both synthetic and real-world datasets to demonstrate the trade-off between privacy and utility.

CCS Concepts: • Security and privacy → Privacy protections; • Theory of computation → Sequential decision making; Online learning algorithms.

Additional Key Words and Phrases: Differential privacy; collaborative learning; contextual bandits

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1 INTRODUCTION

Recommender system is an indispensable component to improve user engagement in modern online information services, such as e-commerce, online advertisement and search engines. As a reference solution for recommendation, collaborative filtering based algorithms achieved impressive success in practice [21, 25, 31]. However, the rapid appearance of new information and new users together with the ever-changing nature of content relevance make traditional offline learning of collaborative filtering incompetent. This motivates the recent developments in online collaborative learning for recommendation, especially contextual bandit based algorithms [1, 15, 23]. Collaborative bandit algorithms provide a principled solution to the explore-exploit dilemma, and enjoy the benefits of collaborative learning paradigm, such as alleviating the cold-start challenge. Recent advances in collaborative bandits include modeling user dependency (e.g., social influence) [6, 37], online user or item clustering [16, 17, 24], and estimating a low-rank structure with latent factors (i.e., matrix factorization based collaborative filtering) [20, 35, 36].

Nevertheless, personalized recommendation is a double-edged sword: the gained utility also comes with the risk of privacy violation. Overly personalized recommendations could be a potential source of privacy vulnerability, for adversaries to take advantage of, e.g., infer users’ sensitive information. Real-world privacy breaches have been reported in Amazon’s recommendation system [5] and Facebook’s advertisement system [22], where an adversary learns considerable amount of information about a user solely based on the systems’ recommendation sequences. Comparing...
to the offline learnt models, online learning methods directly interact with sensitive user data, e.g., user clicks or purchasing history, and timely update the models to adjust their output, which makes privacy an even more serious concern [3, 30, 33, 34]. Realizing its importance, private online learning has recently attracted increasing attention in the research community, with a goal to prevent the algorithm’s sequential output from revealing a user’s private information. While there is existing research on differentially private online convex optimization [18, 33] and contextual bandits [29, 30], private collaborative bandits have not been explored yet.

The challenges regarding the risk of privacy breach in a collaborative bandit based recommender system are unique. In such a system, the algorithm recommends an item to a user, and the user provides feedback (e.g., click) based on his/her true preference. The feedback (reward) is then used to update not only the model’s reward estimation on this user, but also on other users via the imposed dependency among users. As a result, any change in one user’s feedback promptly leads to changes in the algorithm’s output, e.g., different sequences of recommended items, potentially for all users. This is originally designed to improve subsequent recommendations collectively across all users. But a user’s private information could thus be inferred and revealed simply by releasing the recommendation sequence, e.g., extraction attack, even if this user’s feedback is kept private in the system.

In this work, we propose the first study to equip collaborative bandit algorithms with privacy guarantees, under the notion of global differential privacy [11] and local differential privacy [10]. Under global differential privacy, a user is assumed to trust (or say he/she has to trust) the system and provide real engagement data to the system, and the system outputs private recommendations; while under local differential privacy, each user provides perturbed statistics to the system and is no longer required to trust the system or the communication between him/her and system. As the very first study on private collaborative bandits, we focus on algorithms that leverage known dependency (e.g., social connections) among users, such as [6, 37]. Specifically, these algorithms propagate the reward collected from one user to update his/her peers’ bandit models, according to a given and fixed user dependency structure.

One common practice to achieve privacy guarantee is to inject noise to perturb certain statistics derived from private information in the learning process, either on the server side to achieve global differential privacy or on the client side to achieve local differential privacy [8, 10, 11]. However, how to efficiently inject noise in the collaborative bandit learning setting is non-trivial, because of the inherent information sharing mechanism. Specifically, to preserve privacy in collaborative bandits, we apply the tree-based mechanism [7, 12] to add Laplace noise to the models’ statistics to guarantee privacy on each user’s reward feedback (e.g., user clicks). We conduct sensitivity analysis, to which the key is to calibrate the noise scale with respect to the structure of collaboration defined by the user dependency graph. Our insight is that a careful sensitivity analysis over the collaboration structure offers the opportunity to inject minimum amount of noise and better balance the privacy and utility trade-off. In this work, we study two popularly employed collaborative bandit algorithms, Collaborative LinUCB (CoLin) [37] and Gang of Bandits (GOBLin) [17], as the baseline algorithms, which represent two classic types of social network based collaboration structure. We develop their private versions to illustrate a general solution framework for private collaborative bandit. We prove the private algorithms reduce the added regret caused by privacy-preserving mechanism compared to its linear bandits counterparts, i.e., collaboration actually helps to achieve stronger privacy with the same amount of injected noise. We also empirically evaluate the algorithms on both synthetic and real-world public datasets to validate its effectiveness and show the improved trade-off between utility and privacy from our proposed solution framework.

2 RELATED WORK

Collaborative Bandits. Rooted in contextual bandits [1, 4, 23], collaborative bandit algorithms are recently developed to alleviate the cold-start problem in online recommendation. Various solutions have been introduced to enable joint learning across users. For example, Wu et al. [37] modeled dependency among users (e.g., social influence) through a collaborative reward generation assumption; and Cesa-Bianchi et al. [6] leveraged the structure of user dependency as model regularization, where connected users are assumed to have similar model parameters. This type of collaborative bandits require the knowledge of user dependency structure beforehand. Correspondingly, online clustering of bandits studied in [17] avoided such a requirement. Li et al. [24] extended the online clustering to both users and items for collaborative filtering. Matrix factorization based collaborative bandits have been studied in [20, 35, 36], where the collaboration is achieved via a low rank structure over user and item latent factors. In this paper, we focus on privacy guarantees for the first type of collaborative bandits, which take a known collaboration structure as input, and leave the exploration of other types of collaborative bandits as our future work.
**Differential privacy.** Differential privacy [11] provides a formal notion to quantify the amount of information an adversary could obtain by observing the algorithm’s output. The common practice is to add Laplace or Gaussian noise to the output; and the scale of noise depends on privacy budget (often denoted as $\epsilon$) and sensitivity, which is the change of an algorithm’s output caused by the change of input. Prior work has studied the problem of differential privacy for offline collaborative filtering [26, 27, 38]. For example, McSherry and Mironov [27] studied differential privacy for item-based collaborative filtering methods. Liu et al. [26] proposed a differentially private matrix factorization method based on Bayesian posterior sampling. And Jain et al. [19] proposed a private matrix factorization method that guarantees user-level joint differential privacy by perturbing the low-rank decomposition.

Differential privacy was first extended to an online setting for stream data in [7, 12]. Differentially private online learning methods have been studied for online convex optimization [2, 3, 33] and bandit problems [28–30, 34]. The key technique of these solutions is the tree-based mechanism, which was proposed in [7,12] for privately releasing sum statistics in stream data with finite time horizon $T$. Its key idea is to maintain a noisy binary tree where the $T$ leaf nodes are the data points, and the internal nodes in the tree store the sum of all the leaves in its sub-tree. Each node (which represents a partial sum) in the tree is protected with $\frac{\epsilon}{\log(T)}$-differential privacy. Since each sum statistic can be rewritten into $[\log(T)]$ partial sums, composition theorem of differential privacy [27] guarantees the sequence of output sum statistics is $\epsilon$-differentially private.

Based on this tree-based mechanism, (globally) differentially private linear bandit was first studied in [29] with guaranteed privacy in collected user reward feedback. Later, Shariff and Sheffet [30] studied a setting to make both context and reward private. But they adopted a different privacy notion of joint differential privacy. In this paper, we focus on protecting privacy on user’s reward feedback, similar as [29]. However, it is non-trivial to extend the private linear bandits to collaborative bandits setting, where one user’s reward feedback directly contributes to other users’ model update. In other words, the change of model’s input from one user can be measured by the model’s output in (potentially) all users. This propagation of information has to be carefully reflected in sensitivity analysis to avoid trivial solutions. In this paper, we study both global and local differential privacy for collaborative bandits with the key idea of calibrating the noise scale with respect to the structure of collaboration.

3 PRELIMINARIES
To prepare for the discussion of our proposed differentially private collaborative bandit solution framework, we provide a brief overview of contextual bandits, collaborative contextual bandits, and differential privacy in this section.

3.1 Contextual Bandits
In a multi-armed bandit problem, an algorithm sequentially selects an arm $a_t$ from a candidate pool $\mathcal{A} = \{a_1, ..., a_k\}$ at time $t$, and receives the corresponding reward $r_{a_t}$. The goal is to maximize its cumulative reward over a finite time horizon $T$. In a typical contextual bandit setting, each arm $a$ is associated with a $d$-dimensional context vector $x_a$ and its expected reward is governed by a conjecture of the context vector and an unknown bandit model, parameterized as $\theta^*$. For example, in a linear contextual bandit setting, it is assumed that $r_a \sim N(x_a^T\theta^*, \sigma^2)$. A bandit algorithm is evaluated by its pseudo-regret with respect to the optimal arm choice in $T$ rounds of interactions, which is defined as,

$$\text{Regret}(T) = \sum_{t=1}^{T} \left[ E[r_{a_t^*}] - E[r_{a_t}] \right]$$

where $a_t^*$ is the optimal arm to select at time $t$ according to $\theta^*$.

Since the bandit parameter $\theta^*$ is unknown to learner, a contextual bandit algorithm often proceeds by maintaining an online estimate of this parameter at each time $t$ with the observations $\{(x_{a_t}, r_{a_t})\}_{t=1}^{T}$ collected from past interactions. When selecting an arm, the algorithm needs to carefully balance the need for exploitation (trusting the current estimate of $\theta^*$) and the need for exploration (testing new hypotheses to improve the estimation of $\theta^*$).

3.2 Collaborative Contextual Bandits
When applied to personalized online recommendation settings, the unknown bandit model parameter $\theta^*$ is usually attached to each user to reflect their corresponding personalized preferences. We use $\theta_u^*$ to denote the personalized bandit model parameter for user $u$. In a vanilla contextual bandit setting, $\theta_u^*$ for $u \in \mathcal{U}$ are independently estimated based on the observations from the corresponding users. However, due to the existence of mutual influence among users,
a set of isolated bandits can hardly explain all the observed reward feedback, even for a single user, which motivates the research on collaborative contextual bandits.

In a collaborative environment, the expected reward on arm $a$ from user $u$ is assumed to be correlated with those from other users. Collaborative contextual bandit algorithms aim to leverage such dependency information for improved online model estimation and arm selection. Several different ways have been developed to utilize such dependency. The most effective ways can be roughly summarized into the following three categories: 1). additive weighted reward sharing [37]; 2). graph Laplacian based model regularization [6]; 3). online clustering based model sharing [16, 17, 24]. Among these three categories, the first two can be considered as explicit collaboration, as both of them require specific input about how information is shared across users; and the last one can be considered as implicit collaboration, since the structure of collaboration also needs to be inferred from observations by the algorithm. In this paper, we focus on the first two types of collaborative bandit algorithms, and elaborate their details later.

### 3.3 Differential Privacy

For a contextual bandit algorithm that interacts with users over time horizon $T$, denote $S = \{r_t\}_{t=1}^T$ as the reward sequence, where $r_t$ is the reward feedback from user $u_t$ at time $t$. $S'$ is considered as an adjacent neighboring sequence of $S$, if it only differs from $S$ at one point of reward $r_t$. The output of a bandit algorithm $O$ (which is observed by the adversary) is the sequence of its selected arms, i.e., $(a_t)_{t=1}^T$.

**Definition 1 (Global Differential Privacy (DP) [11]).** A randomized mechanism $M$ is $\epsilon$-differentially private if for any adjacent neighboring sequences $\langle S, S' \rangle$ and output, $\mathbb{P}(M(S) \in O) \leq e^\epsilon \mathbb{P}(M(S') \in O)$.

Global differential privacy ensures the adversary observes almost the same output from a private algorithm, in a probabilistic sense, if only one input data point is changed. The difference between the corresponding output is characterized by $\epsilon$. Laplace or Gaussian noise is commonly introduced to disguise the output, where the noise scale is related to the privacy budget $\epsilon$ and the sensitivity of $M$. We formally define sensitivity below.

**Definition 2 (Sensitivity [11]).** For any adjacent neighboring sequences $\langle S, S' \rangle$, global sensitivity of a function $f(\cdot)$ is defined as $\Delta_f = \max_{S, S'} |f(S) - f(S')|$.

Global differential privacy protects sensitive user data from an adversary who has access to the algorithm’s output. But it requires the user to send his/her authentic data to the server. Thus, the server and the communication between user and server have to be trusted. To lift the trust needed from the user, local differential privacy (LDP) is proposed [10]. The key idea is that the privacy mechanism needs to perturb the sensitive statistics on the client side before sending it to the server for further computation. Local differential privacy has been adopted in many real-world applications, such as the RAPPOR system developed by Google to collect web browsing behaviour [14], and Apple provides this privacy protection when collecting users’ usage and typing history [32]. Note that the input and output of a local differential privacy mechanism could be different from the global differential privacy mechanism, even for the same problem, as they impose different privacy requirements. Let $S_i$ be the reward sequence of user $u_i$ such that $\bigcup_i S_i = S$.

The formal definition of local differential privacy is provided below, where a user perturbs his/her private statistics $S_i$ using mechanism $M$ locally, and then send the noisy statistics to the server.

**Definition 3 (Local Differential Privacy (LDP) [10]).** A randomized mechanism $M$ is $\epsilon$-locally differentially private if for any input $\langle S_i, S'_i \rangle$ and output $O$, $\mathbb{P}(M(S_i) \in O) \leq e^\epsilon \mathbb{P}(M(S'_i) \in O)$.

The key difference between LDP and DP is that a DP mechanism takes all users’ data $S$ as input and requires the output to be indistinguishable, while LDP mechanism takes only one user’s data $S_i$ as input and generates randomized responses per user (locally) for downstream tasks.

### 4 DIFFERENTIALLY PRIVATE COLLABORATIVE BANDITS: STARTING FROM COLIN

In this work, we aim to develop a general framework to guarantee global and local differential privacy for collaborative contextual bandits. Due to the intrinsic complexity of the problem, in this section we first develop the private version for a state-of-the-art collaborative bandit algorithm Collaborative LinUCB (CoLin) [37] as an example. We discuss the trade-off between privacy and utility for CoLin via rigorous regret analysis. To note, our solution framework is general.
The recommendation sequences for all users thus have to be perturbed to obtain differential privacy. But instead of adding noise directly to the model’s output, i.e., its choice of arms, we choose to add noise \( \epsilon \) to the context vector \( \tilde{x}_{at,ut} \) at time \( t \) set to \( t = 1 \) to \( T \) do

1. Inputs: \( \delta \in \mathbb{R}_+, \lambda \in [0, 1], W \in \mathbb{R}^{N \times N}, \Delta \)
2. Initialize: \( A_t \leftarrow I_{dN \times dN}, b_t \leftarrow 0, \theta_t^\prime \leftarrow A_t^{-1} b_t \).
3. for \( t = 1 \) to \( T \) do
4. Receive user \( ut \), observe context vectors, \( x_{at,ut} \in \mathbb{R}^d \) and construct \( \tilde{x}_{at,ut} = \text{Vec}(X_{at,ut}W^T) \) for \( \forall a \in A \)
5. Take action \( a_t = \arg \max_{a \in A} \tilde{x}_{at,ut}^T \theta_t + \epsilon_t \sqrt{\tilde{x}_{at,ut}^T A_t^{-1} \tilde{x}_{at,ut}} \), where \( \epsilon_t \) is given by Lemma 2.
6. Observe payoff \( r_{at,ut} \)
7. \( A_{t+1} \leftarrow A_t + \tilde{x}_{at,ut} \tilde{x}_{at,ut}^T, b_{t+1} \leftarrow b_t + \tilde{x}_{at,ut}r_{at,ut} \)
8. Sample noise \( \eta_t \sim \text{TG} \), in which \( \text{TG} = \text{max}_1 L\|W\|_2 \)
9. \( b_{t+1}^\prime \leftarrow b_{t+1} + \eta_t, \theta_{t+1}^\prime \leftarrow A_t^{-1} b_{t+1}^\prime \)
10. end for

and can be applied to other collaborative contextual bandits algorithms. In the next section, we will provide the full picture of our framework and show how it can be applied to another collaborative contextual bandit algorithm Gang of Bandits (GOBlin) [17] with minimum modification in the procedures and analysis.

### 4.1 Global Differential Privacy for CoLin

In Collaborative LinUCB (CoLin [37]), contextual bandit models are placed on a weighted graph \( G = (V, E) \), which encodes the affinity relationship among users. Specifically, each node \( v_i \in V \) in \( G \) hosts a bandit model parameterized by \( \theta_i \) for user \( i \); and the edges in \( E \) represent the affinity relation over pairs of users. This graph is encoded as an \( N \times N \) stochastic matrix \( W \), where each element \( W_{ij} \) is nonnegative and proportional to the influence that user \( i \) has on user \( j \). \( W \) is normalized such that \( \sum_{i=1}^{N} W_{ij} = 1 \) for \( j \in \{1, ..., N\} \), and it is assumed to be time-invariant and known to the learner beforehand. Accordingly, CoLin postulates an additive reward generation assumption: the expected reward \( E[r_{at,ut}] \) is not only determined by user \( ut \)’s own preference on the arm \( at \), but also by that from the neighbors who have influence on \( ut \) as \( E[r_{at,ut}] = \sum_{j=1}^{N} w_{uj} x_{at,ut}^T \theta_j \); or equivalently this can be described as,

\[
\forall t, r_{at,ut} \sim \mathcal{N}(\text{Vec}(\tilde{x}_{at,ut}W^T)\text{Vec}(\Theta), \sigma^2)
\]

where \( \text{Vec}(\cdot) \) is the matrix vectorization operation, \( \Theta \) is a \( d \times N \) matrix consisting of parameters from all the bandits in the graph: \( \Theta = [\theta_1, ..., \theta_N] \), and \( \tilde{x}_{at,ut} \) is a \( d \times N \) matrix with only the column corresponding to user \( ut \) at time \( t \) set to \( x_{at,ut} \) and all the other columns set to zero. By defining \( x_{at,ut} = \text{Vec}(X_{at,ut}W^T) \) and \( \mathcal{N} = \text{Vec}(\Theta) \), Eq (2) can be re-written as \( r_{at,ut} \sim \mathcal{N}(x_{at,ut}^T \Theta, \sigma^2) \).

With such a collaborative reward generation assumption, CoLin appeals to ridge regression for estimating the global bandit parameter matrix \( \Theta \) over all the users at time \( t \). It has a closed-form solution \( \tilde{\Theta}_t = A_t^{-1} b_t \), in which \( A_t = I_{dN} + \sum_{j=1}^{t-1} x_{at,ut}^T x_{at,ut} \) and \( b_t = \sum_{j=1}^{t-1} x_{at,ut}^T r_{at,ut} \). \( I_{dN} \) is an identity matrix and \( \lambda \) is the trade-off parameter for the L2 regularization in ridge regression.

The required information sharing in CoLin brings unique challenges in protecting users’ reward feedback, i.e., the change in one user’s reward feedback can be effectively inferred from all users’ observed recommendation sequences. The recommendation sequences for all users thus have to be perturbed to obtain differential privacy. But instead of directly adding noise to the model’s output, i.e., its choice of arms, we choose to add noise \( \eta_t \) to the sufficient statistics \( b_t = \sum_{j=1}^{t-1} x_{at,ut} r_{at,ut} \) in CoLin, where we sample \( \eta_t \) from a tree-based mechanism [7, 12]. Because differential privacy is immune to post-processing [13], this ensures differential privacy on the algorithm’s output. We name this private derivation of CoLin as (Globally) Differentially Private CoLin (DP-CoLin), and provide its details in Algorithm 1.

The key in DP-CoLin is to derive the sensitivity of CoLin. Analyzing sensitivity in a linear bandit is straightforward [29], as the sensitivity on \( b_t \) can be directly bounded by \( \|x_{at} \|_2 \|r_{at} - r_{at}^d \| \leq L \) where the reward difference is bounded by 1 and the norm of context vector is bounded by \( L \). However, for collaborative bandits, the context vectors encode user dependency and have a higher dimension \( x_{at,ut} \in \mathbb{R}^{dN} \). A trivial bound is \( \|x_{at,ut} \|_2 \|r_{at} - r_{at}^d \| \leq NL \) but we argue this is not tight enough and unnecessarily introduces large noise. Below we analyze the privacy guarantee of DP-CoLin with a tighter sensitivity bound, which calibrates the noise with respect to the structure of collaboration embedded in \( W \).
4.1.1 Privacy Analysis of DP-CoLin. Lemma 1 provides the sensitivity of model statistics \( \mathbf{b}_t \) in CoLin, based on which we develop the privacy guarantee of DP-CoLin.

**Lemma 1 (Sensitivity of \( \mathbf{b}_t \) in CoLin).** Sensitivity of \( \mathbf{b}_t \) is \( \Delta = \max_i L \| \mathbf{W}_i \|_2 \), where \( \mathbf{W}_i \) is the \( i \)-th row of user dependency matrix \( \mathbf{W} \) and \( L \) is the norm of context vector \( \mathbf{x} \).

The proof of this lemma is provided in Appendix. Note that the sensitivity \( \Delta \) of CoLin is related to the structure of \( \mathbf{W} \) and we discuss two extreme cases of \( \mathbf{W} \) to illustrate its effect on privacy protection. Consider when \( \mathbf{W} \) is an identity matrix, the resulting sensitivity by our Lemma 1 is \( L \), which is the same as in linear bandits, since there is no influence among users. When \( \mathbf{W} \) is a uniform matrix, i.e., users have homogeneous influence among each other and \( w_{ij} = \frac{1}{N} \), Lemma 1 shows the sensitivity is \( \frac{\lambda}{\sqrt{N}} \). This result is significant: stronger user dependency in CoLin not only leads to lower regret [37], but also smaller sensitivity of \( \mathbf{b}_t \), which directly reduces the level of required noise to guarantee privacy. This result is also intuitive: when every user has uniform influence on each other, it becomes harder to tell whose action causes the observed change in the algorithm’s output. Less perturbation is thus needed to protect a single user’s privacy. This improvement can hardly be obtained by directly applying existing conclusions on linear bandits.

Based on the above sensitivity analysis, we prove privacy guarantee of DP-CoLin in the following.

**Theorem 1 (Privacy of DP-CoLin).** Algorithm 1 with global sensitivity \( \Delta \) defined in Lemma 1 is \( \epsilon \)-differentially private.

**Proof.** By applying tree-based mechanism [7, 12] with privacy budget \( \epsilon \) and sensitivity \( \Delta \) as shown in line 9–11 of Algorithm 1, the perturbed statistics \( \mathbf{b}_t^\epsilon \) is \( \epsilon \)-differentially private. Since differential privacy is immune to post-processing [13], this consequently makes the model parameter \( \mathbf{\hat{b}}_t^\epsilon \) and the sequence of recommendation \( \{a_t : t \in [1..T]\} \) produced by \( \mathbf{\hat{b}}_t^\epsilon \) also \( \epsilon \)-differentially private. \( \square \)

4.1.2 Regret Analysis of DP-CoLin. We first prove the corresponding confidence bound of parameter estimation in DP-CoLin, i.e., \( \alpha_t \) in line 5 of Algorithm 1, which governs its upper confidence bound based arm selection for online learning. In the following discussion, we use \( \| \mathbf{B} \|_A = \mathbf{vB^T} \mathbf{A} \) to denote the matrix norm of vector \( \mathbf{B} \).

**Lemma 2 (Confidence Bound of DP-CoLin).** For any \( \delta > 0 \), with probability at least \( 1 - \delta \), the estimation error of bandit parameters in DP-CoLin is bounded by:

\[
\| \mathbf{\hat{b}}_t^\epsilon - \mathbf{\hat{b}}^\epsilon \|_{A_t} \leq \sqrt{\frac{dN \log \left( 1 + \frac{\sum_{t' = 1}^{T} \sum_{j=1}^{N} w_{t'j}^2}{\lambda dN} \right) - 2 \log(\delta)}{\epsilon}} + \frac{\Delta}{\lambda^2} \sqrt{\frac{T \log(T) \log(\frac{1}{\delta})}{T}}
\]

The proof is provided in Appendix. The right-hand side of the inequality in Lemma 2 gives us \( \alpha_t \) that is used in line 5 of Algorithm 1 for arm selection. We notice that in order to maintain a private bandit model \( \mathbf{\hat{b}}_t^\epsilon \), the parameter estimation error of DP-CoLin suffers from an additional term \( \frac{\Delta}{\lambda^2} \sqrt{\frac{T \log(T) \log(\frac{1}{\delta})}{T}} \) comparing to that in CoLin due to the added noise \( \eta_t \). Based on Lemma 2, we have the following theorem about the upper regret bound of the DP-CoLin algorithm, which shows the trade-off between privacy budget \( \epsilon \) and regret.

**Theorem 2 (Regret of DP-CoLin).** With probability at least \( 1 - \delta \), the cumulative regret of DP-CoLin algorithm satisfies,

\[
R(T) \leq 2 \sqrt{2dNT \log \left( 1 + \frac{\sum_{t' = 1}^{T} \sum_{j=1}^{N} w_{t'j}^2}{\lambda dN} \right) \frac{\Delta}{\lambda^2}} \sqrt{\frac{T \log(T) \log(\frac{1}{\delta})}{T}} + \frac{\max_i L \| \mathbf{W}_i \|_2}{\epsilon} \log^{1.5} T \log \left( \frac{1}{\delta} \right)
\]

Specifically, the added regret of DP-CoLin comparing to the CoLin is the last term, i.e.,

\[
\frac{2 \max_i L \| \mathbf{W}_i \|_2}{\epsilon} \log^{1.5} T \log \left( \frac{1}{\delta} \right)
\]

We illustrate the proof details in Appendix. From Theorem 2, we can find that the dependency structure plays an important role in the added regret, and again we discuss those two extreme cases of \( \mathbf{W} \) to explain its effect. If \( \mathbf{W} \) is an identity matrix, DP-CoLin algorithm is equivalent to running \( N \) independent private LinUCB [29] for each user and the
added regret is in the order of $O(\frac{N^2}{\epsilon^2} \log^{1.5} T \sqrt{\log \frac{1}{\epsilon} \sqrt{T} \log \frac{1}{\delta}})$. If $W$ is a uniform matrix, the added regret is in the order of $O(\frac{1}{\epsilon^2} T \sqrt{\log \frac{1}{\delta} \sqrt{T} \log \frac{1}{\delta}})$. It is important to note that the collaboration structure also helps reduce the added regret by a factor of $\frac{1}{\sqrt{N}}$. In the meanwhile, the regret reduction from collaboration in the original CoLin is still preserved in the first part of Eq (3) in DP-CoLin.

4.2 Local Differential Privacy for CoLin

Global differential privacy for CoLin requires each user to send true reward (e.g., clicks) to the server, which then aggregates the data, injects noise, and publishes a privacy preserving output. Local differential privacy lifts the trust on the server by asking each user to perturb his/her data locally, before any disclosure to non-trustful server or the communication. Intuitively, this stronger privacy guarantee is at the cost of worse utility.

We present the Locally Differentially Private CoLin algorithm (LDP-CoLin) in Algorithm 2 in Appendix due to the space limit. LDP-CoLin requires a different communication mechanism: instead of directly sending reward $r_{u,i}$ to the server, each user $u$ maintains $b_{u,t} = \sum_{i=1}^{t-1} \hat{x}_{a_{u,i}r_{u,i},u}$ locally as shown in line 8 of Algorithm 2. Each user perturbs their own $b_{u,t}$ by a tree-based mechanism, where noise scales with per-user sensitivity $\Delta_u$ (line 8-9), and then sends it to the server. The server aggregates the received statistics to get $\hat{b}_t^u$ as shown in line 12, and uses it for model estimation and subsequent recommendations. Again in LDP-CoLin the key is to analyze the sensitivity, which controls the minimum amount of noise needed for privacy protection.

4.2.1 Privacy Analysis of LDP-CoLin. We first analyze the sensitivity $\Delta_u$ of $b_{u,t}$ for each user $u$, and then show that Algorithm 2 is locally differentially private using this per-user sensitivity.

**Lemma 3 (Sensitivity of $b_{u,t}$ in CoLin).** Sensitivity of $b_{u,t}$ for user $u$ is $\Delta_u = |W_u|_2$.

The proof is similar to Lemma 1 and the details are provided in Appendix. The main difference is that sensitivity $\Delta_u$ is for a specific user $u$, which only relies on his/her dependent neighbors, i.e., $W_u$.

**Theorem 3 (Privacy of DP-CoLin).** Randomized response $b_{u,t}^\rho$ in Algorithm 2 with sensitivity $\Delta_u$ defined in Lemma 3 is $\epsilon$-locally differentially private.

The proof is similar to DP-CoLin but works in the local setting: as shown in line 8-9 of Algorithm 2, each user $u$ maintains his/her own tree-based mechanism with privacy level $\epsilon$ and sensitivity $\Delta_u$ locally. The local statistics $b_{u,t}$ are perturbed by the tree-based mechanism thus is $\epsilon$-locally differentially private, and thus are $\hat{\theta}_t^u$ and the resulting recommendation sequence.

4.2.2 Regret Analysis of LDP-CoLin. Due to local noise injection, the server’s arm selection strategy has to be revised accordingly, which can be guided by the following lemma.

**Lemma 4 (Confidence Bound of LDP-CoLin).** Let $t_i$ be the number of times where user $i$ interacts with the system up to time $t$, i.e., $\sum_1^{t_i} t_i = t$. For any $\delta > 0$, with probability at least $1 - \delta$, the estimation error of bandit parameters in LDP-CoLin is bounded by,

$$\|\hat{\theta}_t^u - \theta^*\|_{\mathcal{A}_i} \leq \sqrt{dN \log \left(1 + \frac{\sum_{i=1}^{t_i} \sum_{j=1}^{N} w_{u,j}^2}{\lambda dN} \right) - 2 \log(\delta) + \sqrt{\delta} \|\theta^*\|} + \frac{1}{\epsilon} \log \frac{1}{\delta} \sum_{i=1}^{t_i} \log t_i (\lambda_i \log T_i)^2$$

The proof detail is shown in Appendix. Similarly, the right-hand side of the inequality gives us $a_t$ which is used in line 5 of Algorithm 2. Based on it, we have the following theorem about the upper regret bound of LDP-CoLin.

**Theorem 4 (Regret of LDP-CoLin).** With probability at least $1 - \delta$, the cumulative regret of LDP-CoLin algorithm (Algorithm 2) satisfies,

$$R(T) \leq 2 \sqrt{2dNT \log \left(1 + \frac{\sum_{i=1}^{t_i} \sum_{j=1}^{N} w_{u,j}^2}{\lambda dN} \right) \sqrt{\delta} \|\theta^*\|} + \sqrt{dN \log \left(1 + \frac{\sum_{i=1}^{t_i} \sum_{j=1}^{N} w_{u,j}^2}{\lambda dN} \right) - 2 \log(\delta) + \frac{1}{\epsilon} \log \frac{1}{\delta} \sum_{i=1}^{t_i} \|W_i\|^2 \log^3 T_i}$$

Specifically, the added regret of LDP-CoLin comparing to the non-private CoLin is the last term.
Due to space limit, we omit the details of this proof. Note that Theorem 4 is in a general form in which we do not make any assumption about the users’ arriving frequency or order. To better illustrate the added regret, we discuss a special case where the frequency of each user interacting with the system is the same, i.e., $T_i = \frac{T}{N}$. The added regret can thus be simplified as,

$$
\frac{2}{e^2} \log^{1.5} \frac{T}{N} \log \frac{1}{\delta} \sum_{i=1}^{N} \|W_i\|^2 \left[ 2dNT \log \left( 1 + \frac{\sum_{j=1}^{T} \sum_{i=1}^{N} w_{ij}^2}{\lambda dN} \right) \right].
$$

Consider the best case scenario where $W$ is a uniform matrix, e.g., maximum collaboration, the added regret in LDP-CoLin is in the order of $O\left( \frac{\sqrt{N}}{\epsilon} \log^2 \frac{T}{N} \sqrt{T} \log \frac{1}{\delta} \right)$, while DP-CoLin only has the added regret of $O\left( \frac{1}{\epsilon} \log^{1.5} T \sqrt{\log \frac{T}{N} \sqrt{T} \log \frac{1}{\delta}} \right)$.

In fact, in both cases of the illustrative dependency structure, e.g., no collaboration and uniform collaboration, the added regret of LDP-CoLin is roughly $\sqrt{N}$-times larger compared with DP-CoLin, and increases when the number of users grows. This is the inevitable cost to protect privacy in the local (user) level. We verified this relationship between the number of users and regret in our empirical evaluations later as well.

## 5 General Framework and Application to GOBLIN

### 5.1 A General Framework for Differentially Private Collaborative Bandits

Although having different ways of realizing user dependency, the estimation of user preference $\theta$ in a collaborative bandit algorithm can be unified by following: $\theta_t = A_t^{-1} b_t$, with $A_t = I_d + \sum_{r=1}^{t-1} x_{ar, ur}' x_{ar, ur}$, and $b_t = \sum_{r=1}^{t-1} x_{ar, ur}' x_{ar, ur}'$. While the matrix $A_t$ and vector $b_t$ take the same form as that in CoLin, the projected feature vector $x_{ar, ur} \in \mathbb{R}^{dN}$ takes different forms in different collaborative bandit algorithms, because of their unique ways of user dependency modeling.

As discussed in previous section, CoLin [37] enables additive weighted reward sharing through projected feature vector $\tilde{x}_{ar, ur} = \text{Vec}(X_{ar, ur} W^T)$. As another example, GOBLin [6] encodes collaboration through graph regularization in the projected feature vector $\tilde{x}_{ar, ur} = G_g^{-1/2} \text{Vec}(X_{ar, ur})$, where $G = I_d + L$, $L$ is the graph Laplacian of the network and $G_g = G \otimes I$ is the Kronecker product between two matrices $G$ and $I_d$.

We now describe our general framework for equipping collaborative bandit algorithms with differential privacy. The first key step is to add noise $\eta_t$ to the sufficient statistics $b_t$, at each time $t$. Since $b_t$ can be treated as a sum statistic, we sample the noise $\eta_t$ from a tree-based mechanism [7, 12] based on Laplace noise to avoid adding unnecessary noise. However, it is non-trivial to decide the scale (variance) of $\eta_t$, where the scale is proportional to the sensitivity of the private statistics $b_t$. In a collaborative learning framework, the challenge comes from the information sharing in collaborative learning, i.e., the change in one user’s reward feedback promptly leads to changes in all dependent users’ observed recommendation sequences. The trivial sensitivity analysis over $b_t$ in private linear bandits results in a bound of $N\lambda$ and is unnecessarily large. To reduce the amount of noise, our key insight is to conduct tight sensitivity analysis over $b_t$ with respect to the structure of collaboration. Similar to Lemma 1 of DP-CoLin, the tight sensitivity analysis in a private collaborative bandit algorithm makes it possible to inject less noise to guarantee same privacy level comparing to private LinUCB. The significance of this framework lies in that collaboration helps to achieve stronger privacy with the same amount of injected noise. The second key step is to derive the corresponding confidence bound for exploration considering the existence of the Laplace noise in the model (similar to Lemma 2). After the two steps, we obtain a private collaborative bandit algorithm and can analyze its privacy-utility trade-off (similar to Theorem 2).

### 5.2 Global Differential Privacy for GOBLIN

We now show that our solution framework can be easily extended to another state-of-the-art collaborative bandit solution GOBLin [6]. Due to space limit and similarity between GOBLin and CoLin, we do not list the complete algorithm of DP-GOBLin.

In a nutshell, GOBLin models dependency among users by requiring connected users in a network to have similar projected feature vectors $x_{ar, ur} = G_g^{-1/2} \text{Vec}(X_{ar, ur})$ and $G = I_d + L$, $L$ is the graph Laplacian and $I_d$ is a $d \times d$ identity matrix. The parameter $\theta$, matrix $A_t$ and vector $b_t$ take the same form as that in CoLin. To realize the first step of our framework, we add noise $\eta_t$ to $b_t = \sum_{r=1}^{t-1} x_{ar, ur}' x_{ar, ur}'$ to achieve global differential privacy. The noise $\eta_t$ is sampled from a tree-based mechanism and the scale depends on the privacy budget $\epsilon$ and sensitivity $\Delta$ of GOBLin. And to realize the second step of our framework,
we can derive the confidence bound $\alpha_T$ of DP-CoLin as

$$\alpha_T = \sqrt{\frac{2 \log \left( |A_T| \right)}{\delta} + L(\theta_1, \ldots, \theta_N) + \frac{\max_i L\|G_i^{-1/2}\|_2}{\epsilon} \log^{1.5} T \log \frac{1}{\delta}}$$

(5)

where $L(\theta_1, \ldots, \theta_N) = \sum_{i=1}^N \|\theta_i\|_2 + \sum_{(i,j) \in E} \|\theta_i - \theta_j\|_2$ and $|A_T|$ is the determinant of matrix $A_T$.

We now show that our analysis for DP-CoLin can be seamlessly generalized to DP-GOBLin.

**Lemma 5 (Sensitivity of $b_i$ in GOBLin).** Sensitivity of $b_i$ is $\Delta = \max_i L\|G_i^{-1/2}\|_2$, where $G_i^{-1/2}$ is the $i$-th row of the square root of user dependency $G_i$’s graph Laplacian and $L$ is the bound of context vector $x$’s norm.

Proof details of this lemma are provided in the appendix. Before we study the relationship between sensitivity $\Delta$ and the structure of collaboration, specified by $G$, we first rewrite the sensitivity as $\Delta = \max_i L\|G_i^{-1/2}\|_2 = \max_i L\sqrt{G_i^{-1}}$, which is easier to perceive. We again use two extreme cases to explain the derived sensitivity. When $G$ is an identity matrix, i.e., users are independent and disconnected, the graph Laplacian is a zero matrix, $G_i^{-1} = 1$, and the sensitivity $\Delta$ is $L$. When the graph is fully connected, we have $G_{ii} = N$ and $G_{ij} = -1$ for $\forall i, j \in [1..N], i \neq j$. Then its inverse is

$$G^{-1} = \frac{1}{N+1} \begin{bmatrix} 2 & 1 & \ldots & 1 \\ 1 & 2 & \ldots & 1 \\ \vdots & \vdots & \ddots & \vdots \\ 1 & 1 & \ldots & 2 \end{bmatrix}$$

and the sensitivity $\Delta$ is $\frac{2L}{N+1}$. Here again we observe that stronger user connectivity leads to smaller sensitivity of $b_i$.

Based on this analysis, we prove the following privacy guarantee of DP-GOBLin.

**Theorem 5 (Privacy).** DP-GOBLin with global sensitivity $\Delta$ defined in Lemma 5 is $\epsilon$-differentially private.

The proof is similar as the privacy theorem in DP-CoLin using post-processing invariant property of differential privacy, and we omit the proof details. Correspondingly, we have the following theorem about the upper regret bound of the DP-GOBLin algorithm.

**Theorem 6 (Regret of DP-GOBLin).** With probability at least $1 - \delta$, the cumulative regret of DP-GOBLin algorithm satisfies,

$$R(T) \leq 2 \sqrt{T(1 + L^2) \log |A_T| \left( \frac{2 \log \left( |A_T| \right)}{\delta} + L(\theta_1, \ldots, \theta_N) + \frac{\max_i L\|G_i^{-1/2}\|_2}{\epsilon} \log^{1.5} T \log \frac{1}{\delta} \right)}$$

(6)

Specifically, the added regret of DP-GOBLin comparing to the non-private GOBLin is the last term in Eq (6). Similarly as in DP-CoLin, the structure of collaboration, specified by $G$, greatly affects the regret bound in terms of $\|G_i^{-1/2}\|_2$ and $A_T$. For example, larger regret reduction is expected when users are more closely related. Due to space limit, we omit the detailed results here.

### 5.3 Local Differential Privacy for GOBLin

Similar to LDP-CoLin, LDP-GOBLin works in the local setting where each user $u$ maintains $b_{u,T} = \sum_{r=1}^{t-1} x_{ar_{u,T}} \cdot r_{ar_{u,T}}$ locally and perturbs it by a tree-based mechanism with noise scales by per-user sensitivity $\Delta_u$ of GOBLin. Below we first show its local privacy guarantee, and then analyze the trade-off between privacy and regret.

**Lemma 6 (Sensitivity of $b_{u,T}$ in GOBLin).** Sensitivity of $b_{u,T}$ for user $u$ is $\Delta_u = L\|G_u^{-1/2}\|_2$.

**Theorem 7 (Privacy of LDP-GOBLin).** Randomized response $b_{u,T}^\rho$ in LDP-GOBLin with sensitivity $\Delta_u$ defined in Lemma 6 is $\epsilon$-locally differentially private.

**Theorem 8 (Regret of LDP-GOBLin).** With probability at least $1 - \delta$, the cumulative regret of LDP-GOBLin algorithm satisfies,

$$R(T) \leq 2 \sqrt{T(1 + L^2) \log |A_T| \left( \frac{2 \log \left( |A_T| \right)}{\delta} + \max_i L\|G_i^{-1/2}\|_2 \log^{1.5} T \right)}$$

(7)

where $L(\theta_1, \ldots, \theta_N) = \sum_{i=1}^N \|\theta_i\|_2 + \sum_{(i,j) \in E} \|\theta_i - \theta_j\|_2$. 

Specifically, the added regret of LDP-GOBLin comparing to the non-private GOBLin is the last term. Similar to the discussion of LDP-CoLin, we can also compare DP-GOBLin and LDP-GOBLin in the scenario where each user interacts with the system at the same frequency. And the conclusion is also similar: in those two extreme cases, i.e., totally isolated or related users, the added regret of LDP-GOBLin is approximately $\sqrt{N}$-times larger than DP-GOBLin. This again illustrates the required cost for stronger privacy guarantee.

6 EXPERIMENT
We performed empirical evaluations of our developed private collaborative bandit algorithms against several baseline algorithms including the non-private collaborative bandit algorithms CoLin [37] and GOBLin [6], non-private LinUCB [23] and private LinUCB [29]. The datasets include a synthetic dataset from simulation, and two real-world datasets for music recommendation and bookmark recommendation.

6.1 Evaluation Datasets
• Synthetic dataset. To build a synthetic dataset, we follow the settings in [6, 37] to simulate a collaborative online recommendation environment. Specifically, we generate $N$ users, each of which is associated with a $d$-dimensional parameter vector $\Theta^*$, i.e., $\Theta^* = (\Theta^*_1, \ldots, \Theta^*_N)$. Each dimension of $\Theta^*_i$ is drawn from a uniform distribution $U(0, 1)$ and normalized to $\|\Theta^*_i\|_2 = 1$. $\Theta^*$ is treated as the ground-truth bandit parameters for reward generation, and they are withheld from bandit algorithms. We construct the golden relational stochastic matrix $W$ for the graph of users by the users, each of which is associated with a $d$-dimensional feature vector $x_a$, each dimension of which is also drawn from $U(0, 1)$. We normalize $x_a$ by its L2 norm.

To simulate the collaborative reward generation process among users, we compute the reward of arm $a$ for user $i$ at time $t$ as $r_{a_{i,t}} = \text{Vec}(X_{a_{i,t}}W^T)^T\text{Vec}(\Theta^*) + y_t$ following Eq (2), where $y_t \sim N(0, \sigma^2)$. To increase the learning complexity, we delete the edges where $w_{ij}$ is smaller than a predefined threshold, and get the final user graph $G$ by normalizing each column of $W$ by its L1 norm. Note that since $w_{ij}$ is generated proportionally to the similarity between $\Theta^*_i$ and $\Theta^*_j$, the resulting graph naturally satisfies the collaborative assumption in GOBLin [6], i.e., connected users share similar $\Theta^*$. The resulting user graph $G$ represented by the relational matrix $W$ is disclosed to the bandit algorithms. In the end, we generate a size-$K$ arm pool $A$. Each arm $a$ in $A$ is associated with a $d$-dimensional feature vector $x_a$, each dimension of which is also drawn from $U(0, 1)$. We normalize $x_a$ by its L2 norm.

• LastFM and Delicious datasets. The LastFM dataset is extracted from the music streaming service Last.fm, and the Delicious dataset is extracted from the social bookmark sharing service Delicious. The two datasets are created by the HetRec 2011 workshop with the goal of investigating the usage of heterogeneous information in recommender systems.

1The LastFM dataset contains 1,892 users and 17,632 items (artists). The Delicious dataset contains 1,861 users and 69,226 items (URLs). To make these two datasets suitable for evaluating collaborative contextual bandit algorithms, necessary pre-processing is needed. We followed the same pre-processing steps and experiment settings in [6, 37]. To make this paper self-explanatory, we provide a brief description about the pre-processing steps on the two datasets. More details of the pre-processing can be found in [6, 37].

Reward: On LastFM dataset, information about “listened artists” of each user is used to create reward for the bandit algorithms: if a user listened to an artist at least once, the reward is 1; otherwise 0. On Delicious dataset, the reward for bookmarked URLs is set to 1; otherwise 0.

Context features and candidate arms: On both datasets, all tags associated with a particular item are used to create a TF-IDF feature vector, which uniquely represents the content of that item. PCA is used to reduce the dimensionality of the feature vectors to $d = 25$. For a particular user $i$, we generate the candidate arm pool with size $K = 25$ by first selecting one item from those non-zero reward items in user $i$ based on the observations in the dataset, and then randomly selecting the other 24 from those zero-reward items for user $i$.

User relation information: Both datasets contain users’ social network graph, which makes them a suitable testbed for collaborative bandit algorithms. User relation graph is directly extracted from the available social network in the

datasets. In order to make the graph denser and the algorithms computationally feasible, we used graph-cut [9] to cluster users into 100 clusters. Users in the same cluster are assumed to have the same bandit model. After user clustering, a weighted graph can be generated: the nodes are the clusters from the original graph; and the edges between different clusters are weighted by the number of inter-cluster edges in the original graph. Then the relational matrix $W$ in CoLin is obtained by setting $w_{ij} \propto c(i,j)$, where $c(i,j)$ is the number of edges between cluster $i$ and $j$.

### 6.2 Experiment Results

- **Regret comparison.** On the synthetic dataset, cumulative regret is used to evaluate the performance of the compared algorithms. In the real-world datasets, since we do not have an oracle policy, we instead use each learning algorithm’s cumulative reward for evaluation. The cumulative regret (the lower the better) on the synthetic dataset and cumulative reward (the higher the better) on real-world datasets are reported in Figure 1 (a) and Figure 2 respectively. We set the privacy budget $\epsilon = 2$ for all private algorithms in our experiments by default.

In both synthetic and real-world datasets, the non-private collaborative bandits performed better than their globally and locally private counterparts, which is surely expected. We also observe that compared with the globally differentially private collaborative bandit algorithms, i.e., DP-CoLin and DP-GOBLin, the locally differentially private algorithms have significantly worse regret (smaller cumulative reward). This is also expected as local differential privacy is a stronger privacy definition on the user side, and more model perturbation has to be introduced to achieve so. Specifically, as our analysis in Section 4.2 and Section 5.3 suggested, the added regret of LDP collaborative bandit algorithms are roughly $\sqrt{N}$-times larger than their DP counterparts.

We also notice that DP-CoLin and DP-GOBLin performed better than DP-LinUCB in both synthetic and real-world datasets. The improvement comes from two sources: 1) collaborative learning, which improves the convergence rates of model parameter estimation as discussed in [6, 37]; and 2) privacy mechanism under the collaborative environment, which adds less noise than DP-LinUCB when users are not all independent or disconnected. Accordingly to Figure 1 (a), it is obvious that comparing to the regret difference between LinUCB and GOBLin or CoLin, the regret difference between DP-LinUCB and DP-GOBLin or DP-CoLin is much larger. This confirms that the main reason of regret reduction is the calibrated privacy mechanisms developed in this paper.

- **Parameter estimation quality.** To better illustrate the performance of different bandit algorithms, we also studied their parameter estimation quality, which directly measures the algorithms’ online learning convergence. Specifically, we reported the L2 difference between the estimated bandit parameter $\hat{\theta}_t$ and the ground-truth parameter $\theta^*$ in Figure 1 (b). We observe that private collaborative bandit algorithms have a slower model convergence than their non-private counterparts. Moreover, local differential privacy clearly imposes a much larger estimation error comparing to their counterparts with global differential privacy (note that the y-axis is on a log-scale), which further confirms the required cost to guarantee privacy in the local setting.

### 6.3 Detailed Algorithm-level Analysis

To better understand the trade-off between privacy and utility in collaborative bandit learning, we varied the privacy parameter $\epsilon$ and number of users in our evaluation.

- **Effect of privacy budget $\epsilon$.** In Table 1, we reported the cumulative regret of the collaborative bandit algorithms with global and local differential privacy under different privacy parameter $\epsilon$. We vary $\epsilon$ from 0.5 to 10. We run each
experiment for \( T = 10,000 \) iterations and report the average regret of 5 repeated runs. From the results, we notice a clear trade-off between the required privacy level \( \epsilon \) and the resulting regret. Stronger privacy requirement (i.e., a smaller \( \epsilon \)) requires the privacy mechanism to introduce more noise, which directly inflates regret. This result also supports our theoretical analysis that the added regret of the private collaborative bandit algorithms is in the order of \( O(\frac{1}{\epsilon}) \).

- **Effect of number of users** \( N \). In Figure 1 (c), we show the cumulative regret of the collaborative bandit algorithms with global and local differential privacy under different number of users \( N \). We run \( T = 10,000 \) iterations and all users are evenly served for \( T/N \) times. We vary \( N \) from 5 to 50. From the result we observe that the regret increases with the number of users. By looking at the difference between the regret of non-private algorithms and their private versions, we can notice that the added regret increases with number of users \( N \). This also validates our theoretical analysis that the added regret for LDP collaborative bandit algorithms is roughly \( \sqrt{N} \) times larger than their DP versions, which is the inevitable cost to protect privacy at the local level.

7 CONCLUSION

In this work, we studied the problem of protecting global and local differential privacy for collaborative bandits. Our solution framework allows the privacy mechanism to calibrate the noise scale with respect to the user dependency graph. Our theoretical analysis proves the desired privacy guarantee under both settings in two well-studied collaborative bandits. We also rigorously proved the corresponding upper regret bound of the derived private algorithms. Most importantly, we showed the added regret caused by differential privacy mechanism is still sublinear and benefits from the collaboration structure. Extensive experiments on both synthetic and real-world public datasets verified the effectiveness of the private collaborative bandit algorithms, especially the improved trade-off between utility and privacy requirement.

As the first private collaborative bandits research, we explored a specific type of collaborative bandits with explicit knowledge of user dependency structure. In future, we plan to study privacy for other types of collaborative bandit algorithms, such as online clustering-based bandits [17, 24] and matrix factorization based bandits [20, 35, 36]. We also note that the lower regret bound of a DP collaborative bandit algorithm is yet unknown, and it is important to investigate this lower bound to show the optimality of the upper bound of a private collaborative bandit algorithm.

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APPENDIX
Details of Locally Differentially Private CoLin algorithm in Section 4.2:

Algorithm 2 Locally Differentially Private CoLin (LDP-CoLin)

1: Inputs: δ ∈ ℝ+, λ ∈ [0, 1], W ∈ ℝ^{N×N}, A_{1:N}
2: Initialize: A_1 ← \lambda_d I_{N×N}, b_{u,1} ← 0 ∀ u, \hat{\theta}_1 ← A_1^{-1}b_1,
3: for t = 1 to T do
4: // Server side:
5: for \hat{x}_{a_t, u_t} ∈ ℝ^d, and construct \hat{x}_{a_t, u_t} = Vec(\hat{X}_{a_t, u_t} W^T) for ∀ u ∈ \mathcal{U}
6: Take action a_t = arg\max_{a \in \mathcal{A}} \hat{X}_{a_t, u_t}^T \hat{\theta}_t + \alpha_t \sqrt{\hat{x}_{a_t, u_t} A_1 \hat{x}_{a_t, u_t}}, where \alpha_t is given by Lemma 4.
7: // User side:
8: Observe r_{a_t, u_t}
9: Receive user r_{a_t, u_t}
10: Send perturbed statistics b_{a_t, u_t} ← b_{a_t, u_t} + r_{a_t, u_t}
11: Sample noise \eta_{a_t, u_t} ∼ \text{TreeMechanism}_{a_t} (\Delta_{a_t}, \epsilon), in which \Delta_{a_t} = L∥W_{a_t}∥_2
12: Send perturbed statistics b_{a_t, u_t} + \eta_{a_t, u_t} to server
13: \hat{x}_{a_t, u_t} ← t_{a_t, u_t} + 1
14: // Server side:
15: A_{t+1} ← A_t + \hat{x}_{a_t, u_t} \hat{x}_{a_t, u_t}^T, b_{t+1} ← b_{t+1} + \hat{x}_{a_t, u_t} \hat{x}_{a_t, u_t}\hat{\theta}_{t+1} + A_{t+1}^{-1} b_{t+1}
16: end for

Proof of Lemma 1. Consider two reward sequences S_{1:T} and S_{1,T}' that only differ at time i, i.e., r_{a_i, u_i} ≠ r'_{a_i, u_i} and r_{a_i, u_i} = r'_{a_i, u_i} for ∀ j ≠ i. We use L2-sensitivity [8] to measure vector-valued output:

\[ \delta = \max_{S, S'} \left\| \sum_{t} \hat{x}_{a_t, u_t} r_{a_t, u_t} - \sum_{t} \hat{x}_{a_t, u_t} r'_{a_t, u_t} \right\|_2 \leq \max_{S, S'} \left\| \text{Vec}(\hat{X}_{a_t, u_t} W^T)(r_{a_t, u_t} - r'_{a_t, u_t}) \right\|_2 \leq L∥W∥_2 \]

The first inequality holds because the reward r_a for any arm a is bounded in [0, 1], and the last inequality holds because the norm of context vector x_{a_t, u_t} is bounded by L.

Proof of Lemma 2. We decompose the estimation error into two terms, ∥\hat{\theta}_{t} - \theta^{*}∥_{A_t} ≤ ∥\hat{\theta}_{t} - \theta^{*}∥_{A_t} + ∥\hat{\theta}_{t} - \hat{\theta}_{t}∥_{A_t}, where \hat{\theta}_{t} is the model parameter of non-private CoLin, i.e., \hat{\theta}_{t} = A_t^{-1}b_t.

The first term is the estimation error of non-private CoLin, and can be directly bounded using Lemma 1 in [37]. The second term is the estimation difference between DP-CoLin and non-private CoLin caused by privacy mechanism, and can be bounded as follows:

\[ \Delta = \max_{t} \left\| \text{Vec}(\hat{X}_{a_t, u_t} W^T)(r_{a_t, u_t} - r'_{a_t, u_t}) \right\|_2 \leq \max_{t} |\eta_t| L \leq \frac{L}{\epsilon} \log T \sqrt{\log T} \log \frac{1}{\delta}, \]

where the last inequality is the bound of noise added by tree-mechanism from Theorem 3.6 of [7], and \Delta is the sensitivity derived in Lemma 1.

Proof of Theorem 2. According to the definition of regret in Eq (1), the regret of DP-CoLin at time t can be written as:

\[ R_t = r_{a_t, u_t} - r_{a_t, u_t} = \text{Vec}(\hat{X}_{a_t, u_t} W^T) \hat{\theta}_{t-1} - \text{Vec}(\hat{X}_{a_t, u_t} W^T) \theta^{*} \]

\[ \leq \text{Vec}(\hat{X}_{a_t, u_t} W^T) \hat{\theta}_{t-1} + \alpha_t \text{Vec}(\hat{X}_{a_t, u_t} W^T) \|\hat{X}_{a_t, u_t} W^T\|_{A^{-1}_t} - \text{Vec}(\hat{X}_{a_t, u_t} W^T) \theta^{*} \]

\[ \leq \text{Vec}(\hat{X}_{a_t, u_t} W^T) \hat{\theta}_{t-1} + \alpha_t \text{Vec}(\hat{X}_{a_t, u_t} W^T) \|\hat{X}_{a_t, u_t} W^T\|_{A^{-1}_t} - \text{Vec}(\hat{X}_{a_t, u_t} W^T) \theta^{*} \]

\[ \leq 2\alpha_t \text{Vec}(\hat{X}_{a_t, u_t} W^T) \|\hat{X}_{a_t, u_t} W^T\|_{A^{-1}_t} \]
in which $\sigma_T$ is the upper bound of $\|\hat{\theta}_T - \theta^*\|_{\mathcal{A}_T}$ and it can be explicitly calculated based on Lemma 2. The first inequality is based on the definition of confidence bound and the second inequality holds because of the UCB-type arm selection strategy (line 6 of Algorithm 1).

The cumulative regret at time $T$ in DP-CoLin can be bounded by,

$$R(T) \leq \sqrt{\frac{T}{2} \sum_{t=1}^{T} R_t^2} \leq \sqrt{T4\sigma_T^2 \sum_{t=1}^{T} \|\text{Vec}(\hat{X}_{a_t}W^T)\|_{\mathcal{A}_T^{-1}}^2} \leq \sqrt{T8\sigma_T^2 \ln \left(\frac{|\mathcal{A}_T|}{\lambda dN}\right)}$$

where the third inequality is based on Lemma 11 in [1]. Substitute $\sigma_T$ with the bound from Lemma 2 gives the final result.

PROOF OF LEMMA 3. Similar to the proof of Lemma 1, for user $u$ we consider two neighbouring reward sequence $S_{u,1:T_u}$ and $S'_{u,1:T_u}$ that only differs at time $j$, i.e., $r_{a_j,u} = r'_{a_j,u}$ for $\forall j \neq i$.

$$\Delta = \max_{S_u, S'_u} \|b_{u,t} - b'_{u,t}\|_2 = \max_{S_u, S'_u} \|\sum_{t'=1}^{t-1} b_{a_{t'},u} r_{a_{t'},u} + \sum_{t'=1}^{t-1} b'_{a_{t'},u} r'_{a_{t'},u}\|_2$$

$$= \max_{S_u, S'_u} \|\text{Vec}(\hat{X}_{a,u}W) (r_{a,u} - r'_{a,u})\|_2 \leq \max_{S_u, S'_u} \|\text{Vec}(\hat{X}_{a,u}W)\|_2$$

PROOF OF LEMMA 4. Similar to the analysis in Lemma 2, we separate the estimation error into two terms, $\|\hat{\theta}_T^l - \theta^*\|_{\mathcal{A}_T} \leq \|\hat{\theta}_T^l - \theta^*\|_{\mathcal{A}_l} + \|\hat{\theta}_T^l - \hat{\theta}_l\|_{\mathcal{A}_T}$, where the second term is the estimation difference between LDP-CoLin and non-private CoLin caused by local privacy mechanism. This term can be bounded as follows:

$$\|\hat{\theta}_T^l - \hat{\theta}_l\|_{\mathcal{A}_T} = \|\hat{\theta}_T^l (\sum_{i=1}^{N} \eta_i)\|_{\mathcal{A}_l} = \|\sum_{i=1}^{N} \eta_i\|_{\mathcal{A}_l} \leq \|\sum_{i=1}^{N} \eta_i\|_{\mathcal{A}_T}$$

Since each $\eta_i$ is sampled from a tree-based mechanism, and is the sum of at most $\log t_i$ Laplace noise $Lap\left(\frac{A_l \log T_i}{e}\right)$. To further bound the term, we use Corollary 2.9 of [7] regarding sum of independent Laplace noise and get

$$\|\sum_{i=1}^{N} \eta_i\|_{2} \leq \sum_{i=1}^{N} \sum_{t=1}^{\log t_i} \left(\frac{\log t_i}{\epsilon}\right)^2 \leq \frac{1}{\epsilon} \log \frac{1}{\delta} \sum_{t=1}^{N} \log t_i (A_l \log T_i)^2$$

and it holds with probability at least $1 - \delta$.

Details of Globally and Locally Differentially Private GOBLin in Section 5:

PROOF OF LEMMA 5. Consider two neighbouring reward sequence $S_{1:T}$ and $S'_{1:T}$ that only differs at time $i$, i.e., $r_{a_i,j} = r_{a_i,j'}$ for $\forall j \neq i$.

$$\Delta = \max_{S, S'} \|b_{i,j} - b'_{i,j}\|_2 = \max_{S, S'} \|\sum_{t'=j}^{i-1} b_{a_{t'},u} r_{a_{t'},u} + \sum_{t'=j}^{i-1} b'_{a_{t'},u} r'_{a_{t'},u}\|_2$$

$$= \max_{S, S'} \|\text{Vec}(\hat{X}_{a,u}G^{-1/2})(r_{a,u} - r'_{a,u})\|_2 \leq \max_{S, S'} \|\text{Vec}(\hat{X}_{a,u}G^{-1/2})\|_2$$

Where the fourth equality is based on the vector trick of Kronecker product.
REFERENCES