Email ID: \_\_\_\_\_

# CS 3330 Exam 2 – Spring 2016

Name: \_\_\_\_\_

# Computing ID:

**Letters** go in the boxes unless otherwise specified (e.g., for **C** 8 write "C" not "8"). **Write Letters clearly**: if we are unsure of what you wrote you will get a zero on that problem. **Bubble and Pledge** the exam or you will lose points.

Assume unless otherwise specified:

- little-endian 64-bit architecture
- %rsp points to the most recently pushed value, not to the next unused stack address.
- questions are single-selection unless identified as select-all

**Multiple-select**: are all clearly marked; put 1 or more letters in the box, or write "None" if no options are correct.

**Mark clarifications**: If you need to clarify an answer, do so, and also add a **\*** to the top right corner of your answer box.

.....

Question 1: Which of the following caches has the largest blocks?

- A A full-associative cache with 6-bit tags and 128 lines
- **B** A direct-mapped cache with 8-bit tags and 64 lines
- C A direct-mapped cache with 6-bit tags and 128 lines
- **D** A full-associative cache with 8-bit tags and 64 lines
- **E** All of the above have the same block size
- **F** Which of the above depends on the address size

#### Information for questions 2–4

Given a 12-bit address jklmnopqrstu, where j is the high-order bit, what are the parts of the address for a **direct-mapped** cache with 32 lines and 8-byte blocks?

Answer by putting one or more letters in the box; e.g., klmn would be the second-, third-, fourth-, and fifth-highest-order bits

**Question 2:** (see above) The index is

Question 3: (see above) The tag is

Question 4: (see above) The offset is

Answer:
Answer:
Answer:

Answer:	

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#### Information for questions 5–9

Suppose we expand our pipeline to ten stages as follows:

- Fetch has three stages; the PC is needed at the beginning of  $F_1$ ; valP (that is, old PC + offset) is available at the end of  $F_1$ ; and the parts of the instruction (icode, valC, etc) are available at the end of  $F_3$ .
- Decode is still one stage.
- Execute is two stages; operands are needed at the beginning of E<sub>1</sub> and results available at the end of E<sub>2</sub>. Condition code checking for conditional jumps and moves is performed in E<sub>2</sub>.
- Memory is three stages; the address and whether to read or write is needed at the beginning of M<sub>1</sub>; the value to write is needed at the beginning of M<sub>2</sub>; and a value read from memory is available at the end of M<sub>3</sub>.
- Writeback is still one stage.

In the following questions, **assume we forward** instead of stalling wherever possible. Answer each question with an integer between 0 and 10.

**Question 5:** (see above) How many stages of speculative work would be erased on a branch misprediction?

Answer:	

**Question 6:** (see above) We said valP was available in  $F_1$ . Why would a version where valP was not available until  $F_2$  be a very bad idea? Write a short answer here:

<b>Question 7:</b> (see above) How many cycles of stalling would be needed to resolve a (memory) load-use hazard?	Answer:
Question 8: (see above) How many cycles of stalling would be needed to run addq %rax, %rbx addq %rbx, %rcx	Answer:
<b>Question 9:</b> (see above) How many cycles of stalling would a ret hazard create?	Answer:

**Question 10:** Of all of the Y86-64 instructions, only popq uses both write ports E and M in the register file, one for %rsp and one for storing the value popped from the stack. Each of the other instructions use at most one register-file write port.

ret is similar to popq in many ways; why does ret not need two register file write ports?

- A It does not update the stack pointerB It updates the stack pointer, but not a register in the register file
- **C** The value it pops from the stack is stored, but not in the register file
- **D** The value it pops from the stack is not stored



### Information for questions 11–13

The following questions are grouped because they have the same set of possible answers.

**Question 11:** (see above) Which of the following optimizations can change the big-O of the code?

- **A** loop unrolling
- **B** function inlining
- **C** multiple accumulators
- **D** moving work outside of loops
- **E** none of the above

**Question 12:** (see above) Which of the following optimizations is designed to take advantage of pipelines?

- A loop unrolling
- B moving work outside of loops
- **C** multiple accumulators
- **D** function inlining
- **E** none of the above

**Question 13:** (see above) Which of the following optimizations is designed to take advantage of caches?

- **A** function inlining
- **B** loop unrolling
- **C** multiple accumulators
- **D** moving work outside of loops
- **E** none of the above

#### Information for questions 14–15

In the following diagrams, indicate the control signals to give each pipeline register by putting a single letter in each box; use N for normal, B for bubble, and S for stall.

**Question 14:** (see above) Assume that  $i_4$  and  $i_5$  resulted from incorrect speculative execution and should not be allowed to continue but all other instructions are OK and may continue to execute normally.



**Question 15:** (see above) Assume that  $i_3$  requires some information not yet available and needs another cycle in the E stage, but all other instructions are OK and may continue to execute normally.



Answer:	

Answer:

Answer:	

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**Question 16:** When running Y86-64 instructions on the single-stage SEQ implementation described in the textbook and implemented in the first HCL homework, we may encounter

- **A** dependencies
- **B** hazards
- **C** both of the above
- **D** neither of the above

**Question 17:** Longer pipelines tend to achieve what compared to shorter pipelines? **Select all that apply** 

- **A** lower throughput
- **B** higher latency
- **C** lower latency
- **D** higher throughput

**Question 18:** Assume we have a 4-line direct-mapped cache. If we access the following addresses (split into parts for you) in the order listed, which will hit and which will miss? Answer by putting an H or M in each box.

_	Tag	Index	Offset
	0	0	0
	0	0	8
	1	0	0
	0	3	8
	0	0	0
	0	3	0

#### Information for questions 19–21

Given a 12-bit address jklmnopqrstu, where j is the high-order bit, what are the parts of the address for a **set-associative** cache with 8 sets of 4 lines each and 8-byte blocks?

Answer by putting one or more letters in the box; e.g., klmn would be the second-, third-, fourth-, and fifth-highest-order bits

Question 19:	(see above) The index is	Answer:
Question 20:	(see above) The tag is	Answer:
Question 21:	(see above) The offset is	Answer:

1103:	
Answer:	

Answer:

### Information for questions 22–23

The following questions ask about how specific characteristics of caches.

**Question 22:** (see above) Suppose I know that memory accesses in my program will be purely sequential, with addresses being accessed in order.

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Which of the following will have the largest impact on memory performance?

- **A** high levels of associativity
- **B** total cache size
- **C** number of individual blocks
- **D** individual block size

**Question 23:** (see above) Suppose I know that memory accesses in my program will be completely random, with no temporal or spatial locality.

Which of the following will have the largest impact on memory performance?

- **A** total cache size
- **B** number of individual blocks
- **C** high levels of associativity
- **D** individual block size

**Question 24:** Which of the following code snippets exhibits the most spatial locality?

```
A for(c=0; c<n; c+=1)
      for(r=0; r<n; r+=1)</pre>
        a[r][c] += b[r][c];
B for(r=0; r<n; r+=1)
      for(c=0; c<n; c+=1)</pre>
        a[r][c] += b[c][r];
C for(c=0; c<n; c+=1)
      for(r=0; r<n; r+=1)</pre>
        a[r][c] += b[c][r];
D for(r=0; r<n; r+=1)
      for(c=0; c<n; c+=1)</pre>
        a[r][c] += b[r][c];
```

**Question 25:** Suppose that for a particular chip the multiplication operation is pipelined and takes three cycles to produce a value (the book would call this "latency 3, issue 1"). This would mean that computing ((a\*b)\*c)\*d would require 9 cycles. Using reassociation, what is the minimal number of cycles needed for that same product?

Α	7	
В	8	<b>A</b>
С	4	Answer:
D	9	
Е	6	
F	3	
G	5	

Answer:	

Answer:

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Answer:

**Question 26:** For our five-stage pipeline we had to implement prediction for conditional jumps in order to avoid extra stalls in the pipeline, but we did not need to do that for conditional moves.

If the following instructions were also conditional (e.g., if halt had conditional forms like haltge and haltl), which ones would need prediction to avoid stalls?

### Select all that apply

- A halt
- **B** OPq
- **C** call
- $\boldsymbol{\mathsf{D}} \quad \texttt{popq}$
- E irmovq

## **Pledge:**

On my honor as a student, I have neither given nor received aid on this exam.

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