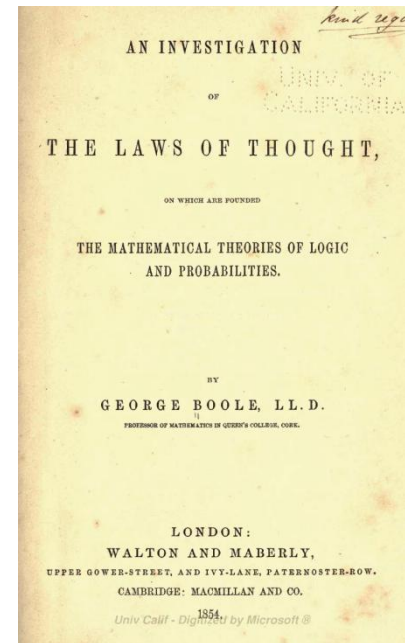
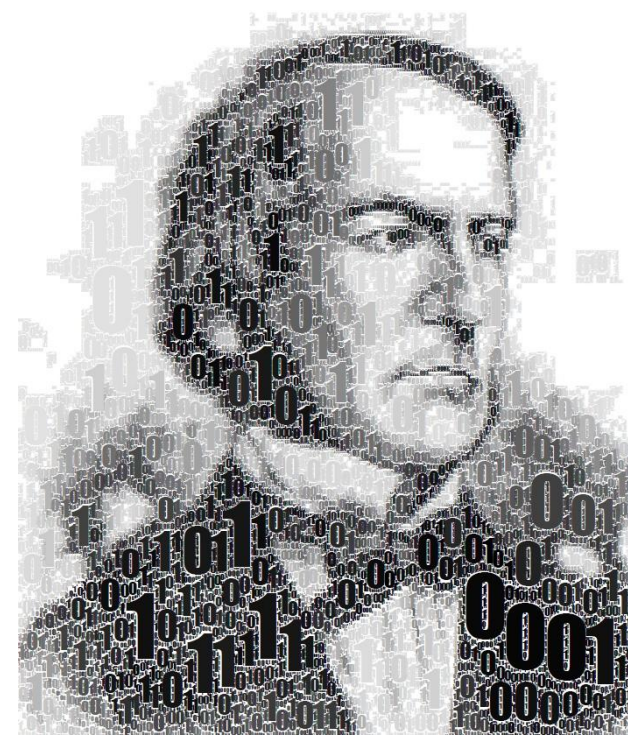



Historical Perspectives

George Boole (1815-1864)

- Mathematician and philosopher
- Invented symbolic / **Boolean logic**
- Invented **Boolean algebra**, i.e. “calculus of reasoning”
- A **founder of computer science**
- “An Investigation into the Laws of Thought”
- **Influenced** De Morgan, Schröder, Shannon
- All modern computers, electronics, phones, data transmission, rely on **Boolean principles**



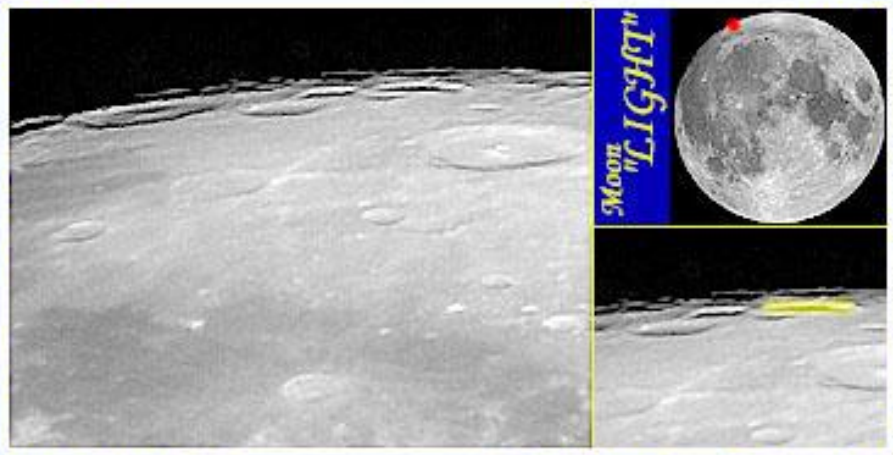

GEORGE BOOLE
 LL.D., D.C.L., F.R.S.
 1815 - 1864
 GEORGE BOOLE, FATHER OF MODERN ALGEBRA, AUTHOR OF THE LAWS OF THOUGHT AND FIRST PROFESSOR OF MATHEMATICS AT UNIVERSITY COLLEGE, CORK, WAS BORN IN LINCOLN AND ESTABLISHED AN ACADEMY IN THIS HOUSE C. 1810.

BOOLE

63 km

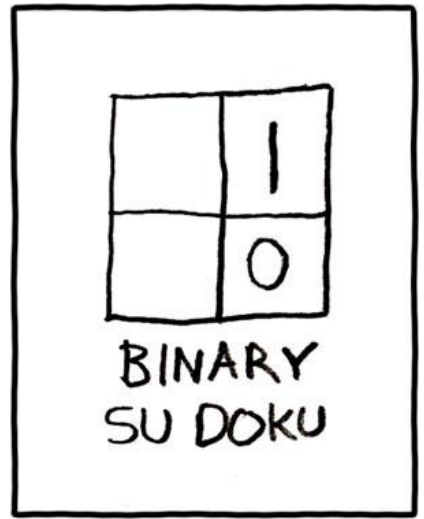
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97 / 10 / 15 D=254mm f/D=10



B/W QuickCam

a.cidadao@mail.telepac.pt



All cats have four legs.
 I have four legs.
 Therefore, I am a cat.



BOOLE ORDER'S LUNCH

NO, NO, YES, NO, NO, YES,
YES, NO, NO, NO, YES...



**Boolean
humor**

There are only 10 types
of people in the world:
Those who understand binary
and those who don't.

01000100
01000001
01000100

s.hastis

001010 0010,0010

00101010 0010001110:

00100100100001 00100100001110101 0001001 00100101
0101000010010 00 00100 00100010010 001010010 (100100
0001010 0001010010010 001000100100 0010010)0010010
001000100010001 0010 00100100. 001000100 00100100 010
0001000100 0001100100 0010001001000010 0010001 0001.
001001001 01001001 0100010 ? 0010 1001001 0000100 0100
0010001, "1001000 0010 10001 00101010 001 00100100 0010
001001 - 100100 100 001000 001 001000 101001." 010010
001001 000100 01000 001000 00100 001000 10001111 10001
10001001 01000100 00101100010 0001000 01101010 00
01001.

1001000 1000 01000 001001 1010001 000 1000 001
0100100 0100100 01000 1010100110 010001110001
0010011110001 0100101 10100010010010.

10011000101000100100,

100110101

BINARY LETTER FROM GRANDMA

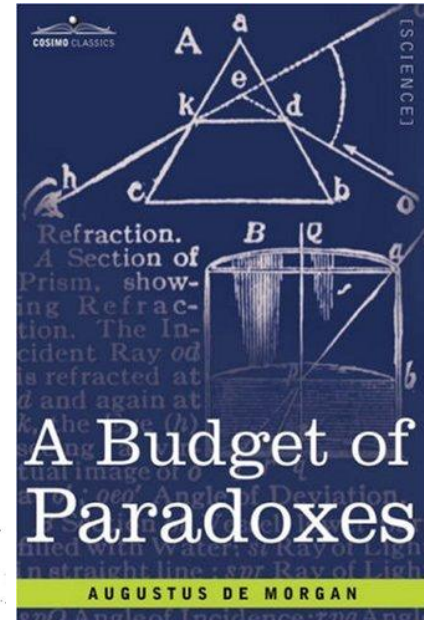
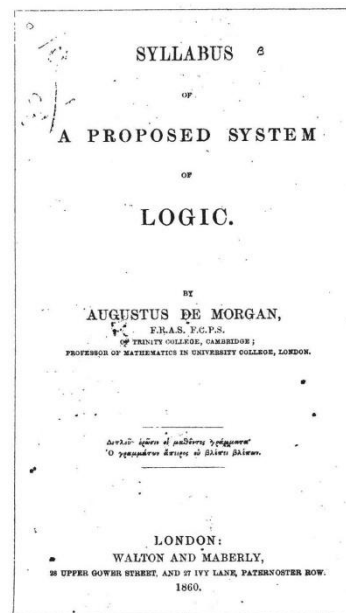
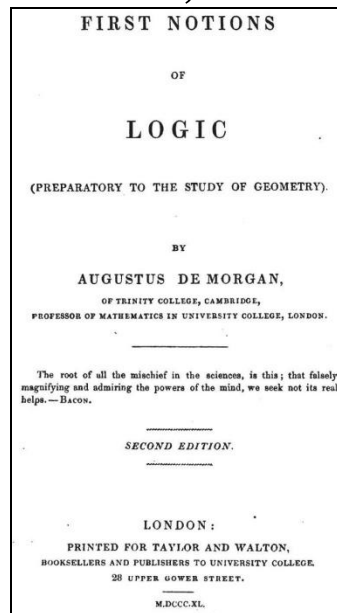
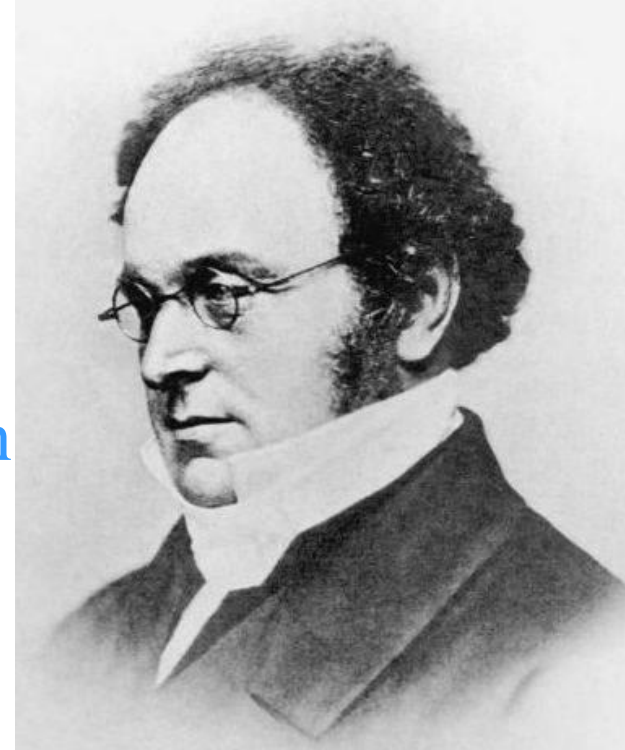


Mozart writing the digital version of his symphony No. 38
in D major.

Historical Perspectives

Augustus De Morgan (1806-1871)

- Mathematician and logician
- Developed logic & mathematical induction
- De Morgan's Laws in logic & set theory
- Invented relational algebra
- Corresponded extensively with Hamilton
- Influenced Russell, Whitehead, and Tarski
- Studied paradoxes



Historical Perspectives

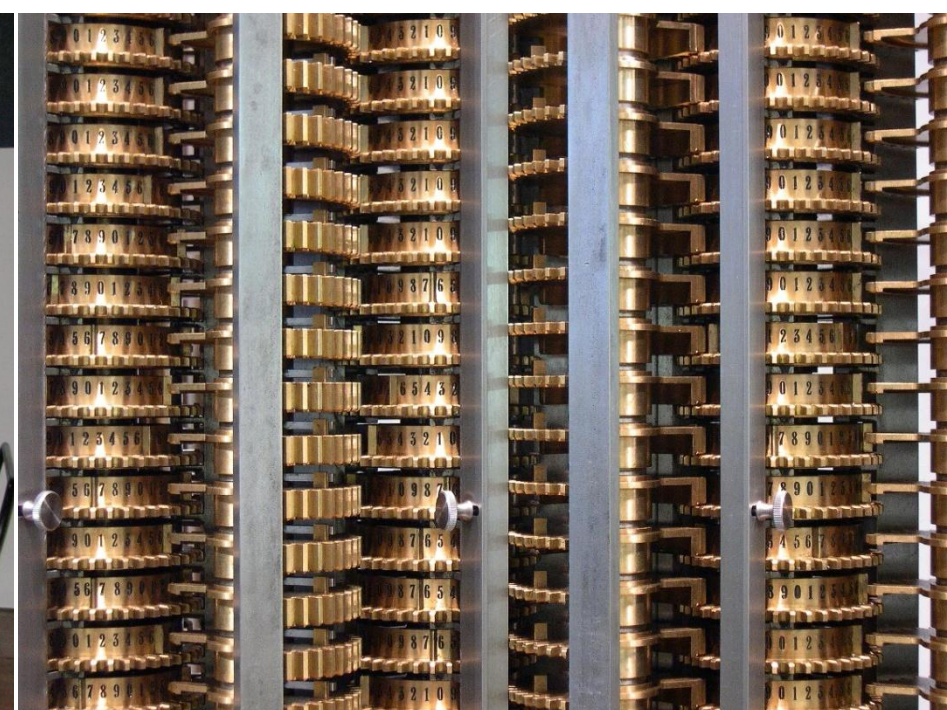
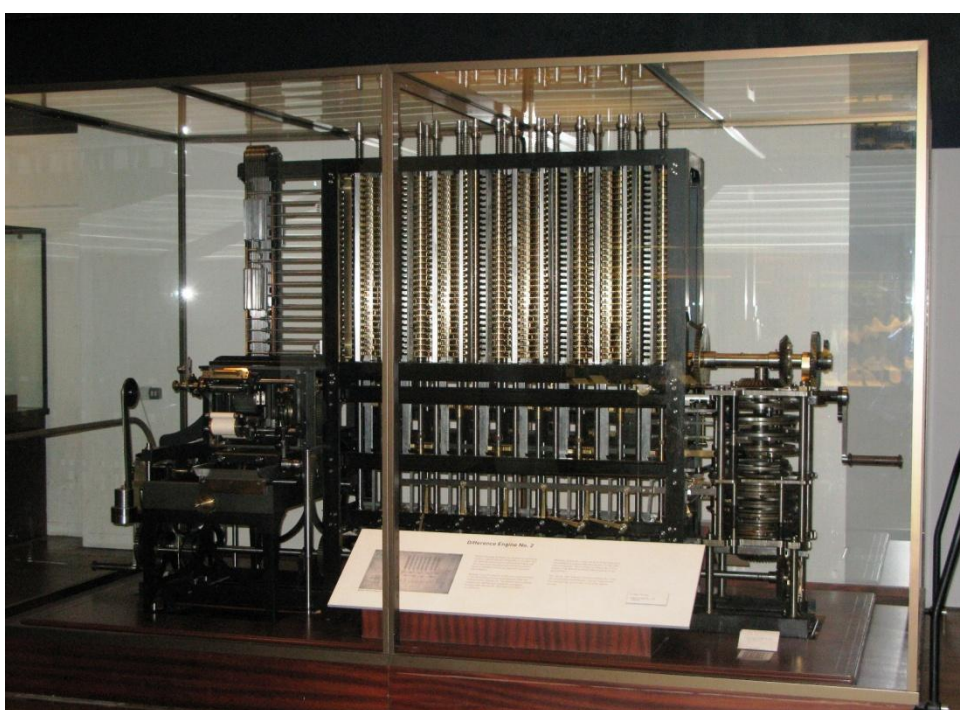
Charles Babbage (1791-1871)

- Mathematician, philosopher, inventor mechanical engineer, and economist
- The **father of computing**
- Built world's **first mechanical computer**
 - the “**difference engine**” (1822)
- Originated the **programmable computer**
 - the “**analytical engine**” (1837)
- Worked in **cryptography**
- Developed **Babbage's principle** of division of labor

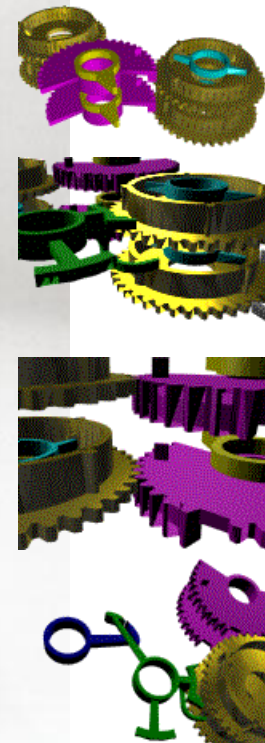
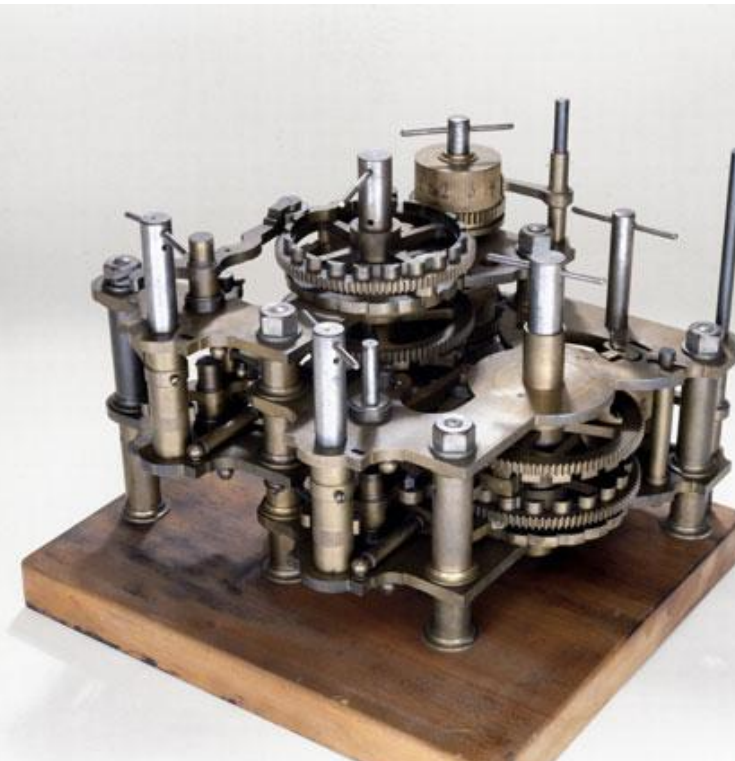
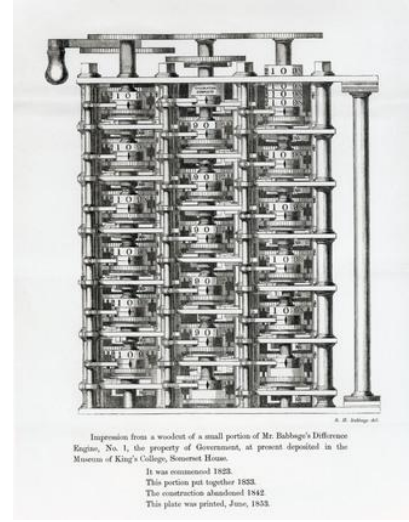
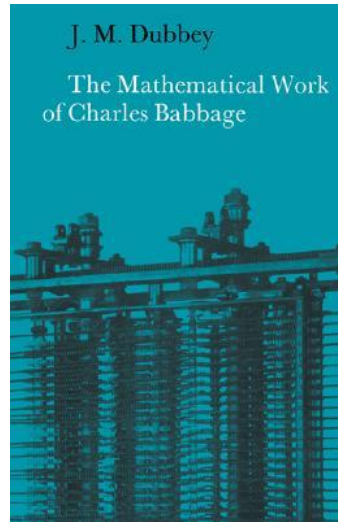
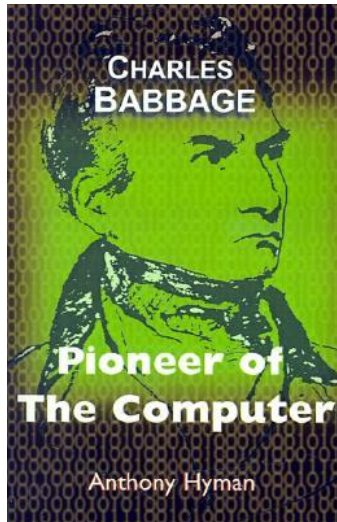
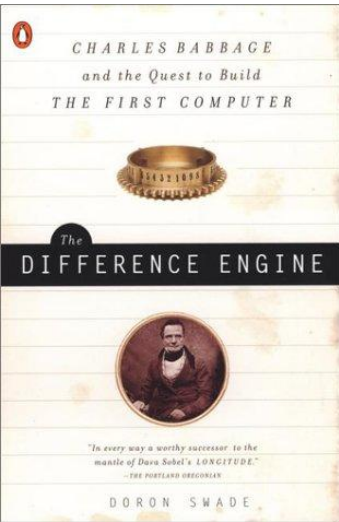


Babbage's Difference Engine

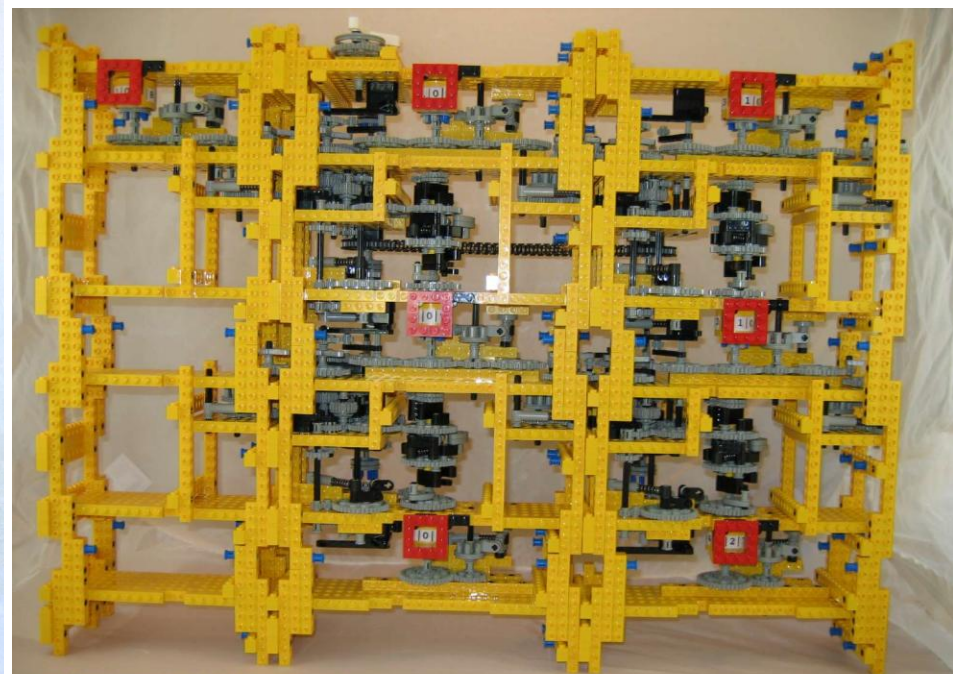
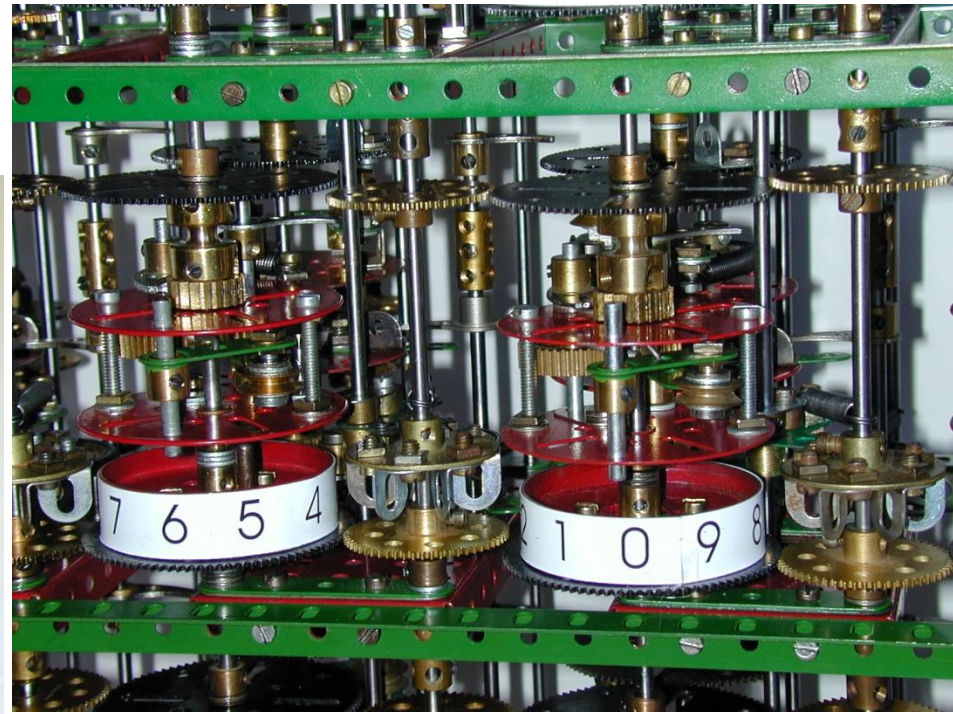
- World's **first mechanical computer**
- Designed in **1822**, redesigned in 1847-1849
- **25,000 parts**, 15 tons, 8ft tall, 31 digits of precision
- Tabulated polynomial functions, used **Newton's method**
- **Approximated** logarithmic and polynomial functions
- Used **decimal number system** and hand-crank



Babbage's Difference Engine

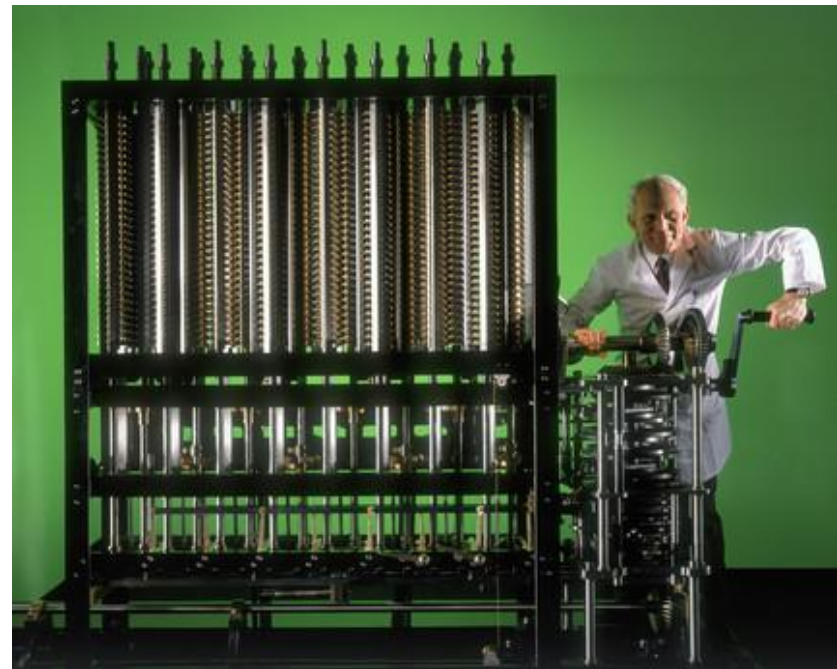
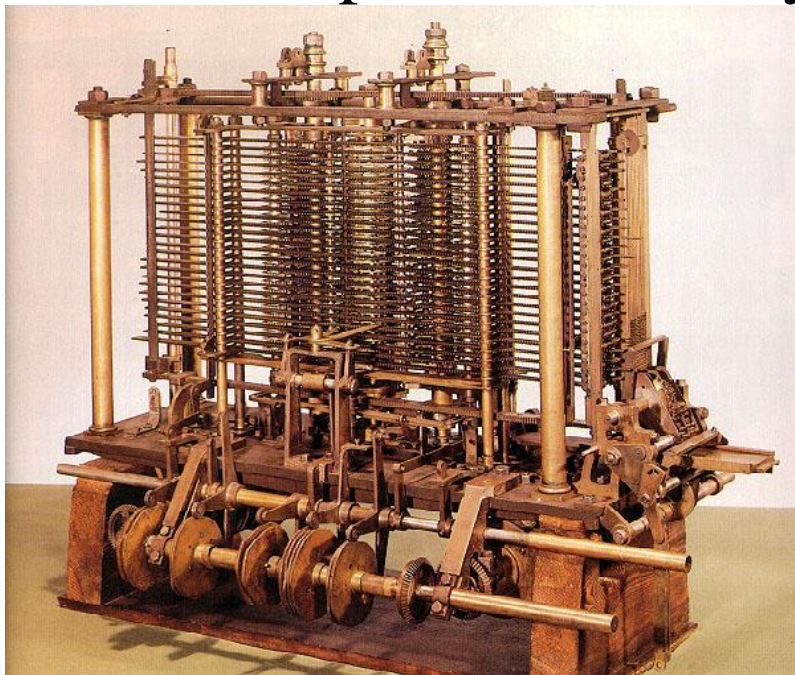


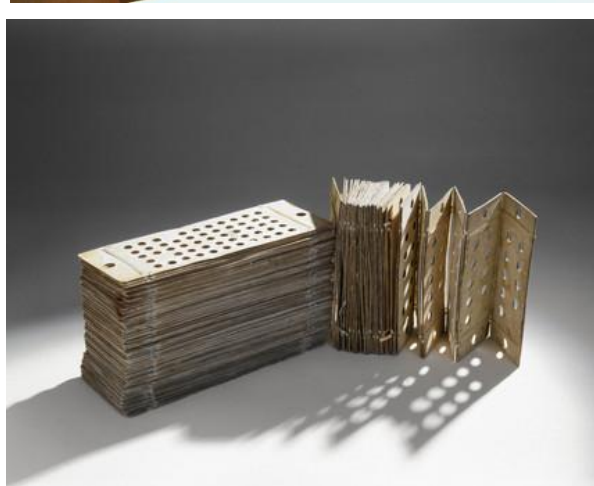
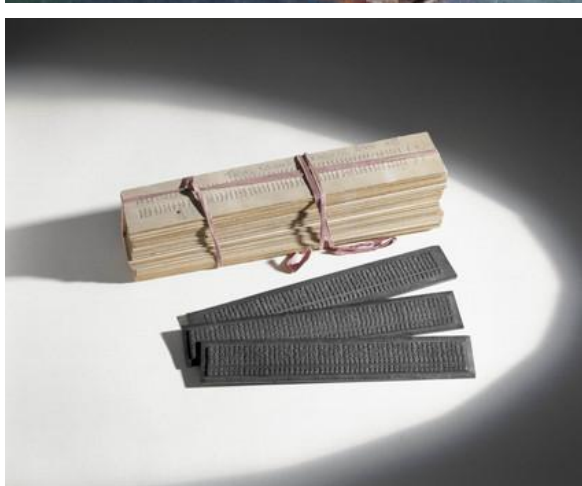
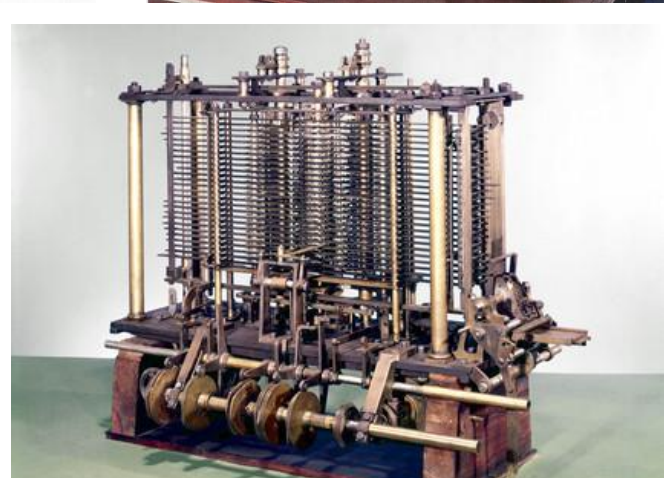
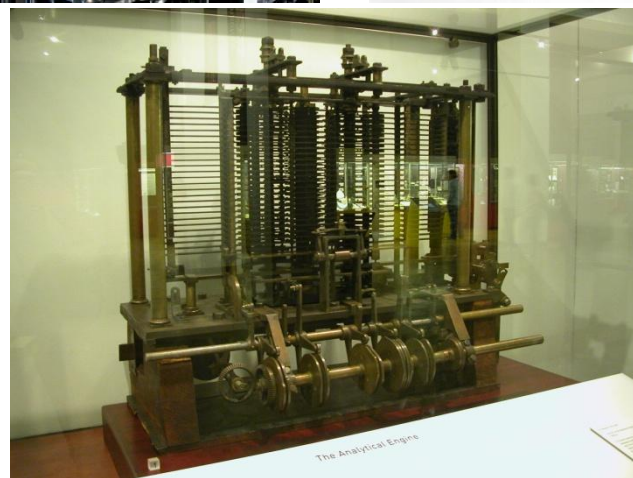
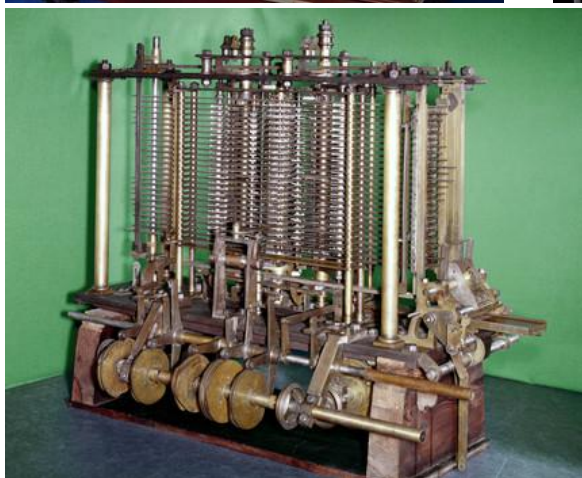
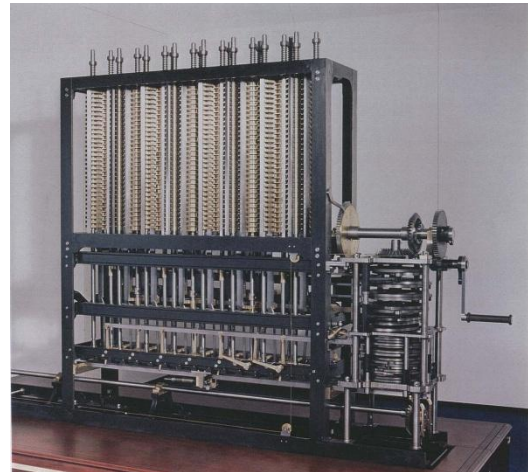
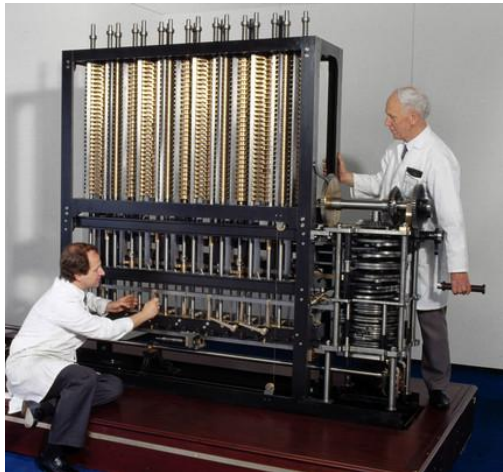
Babbage's difference engine built from Mechano and Lego

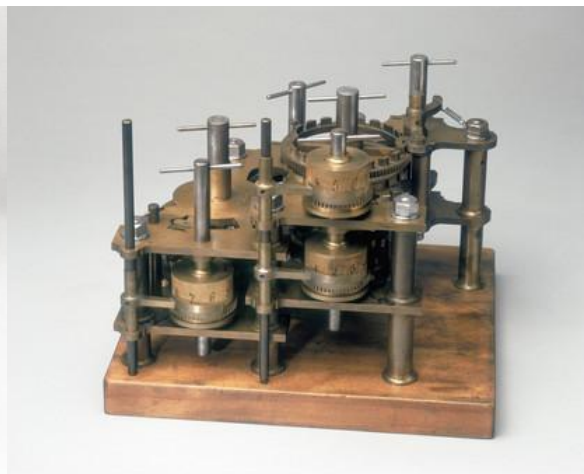
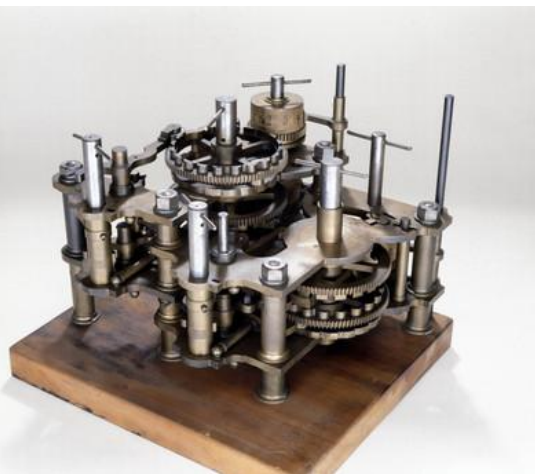
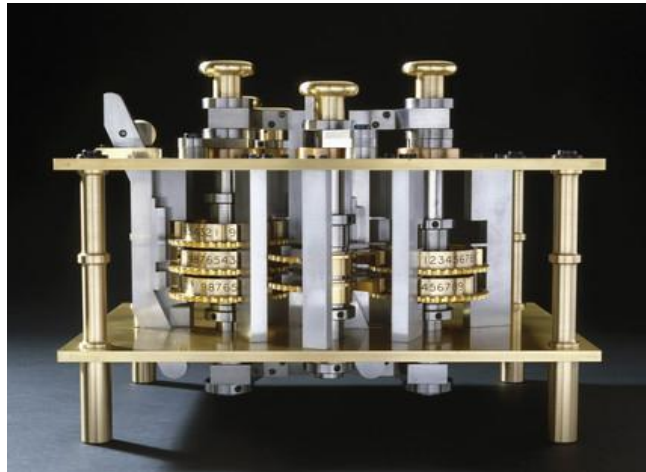


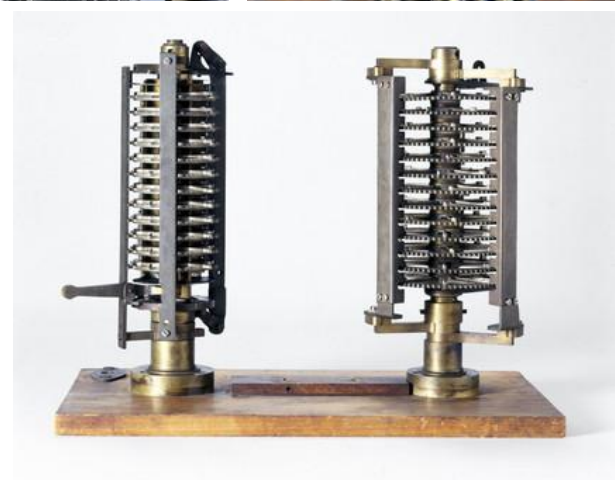
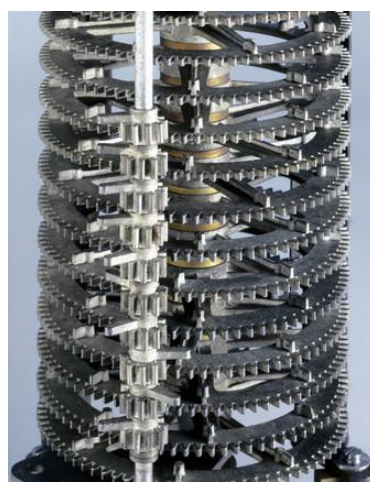
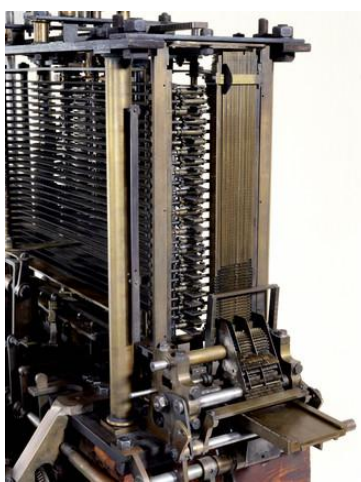
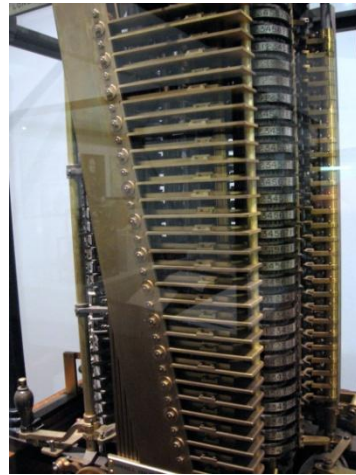
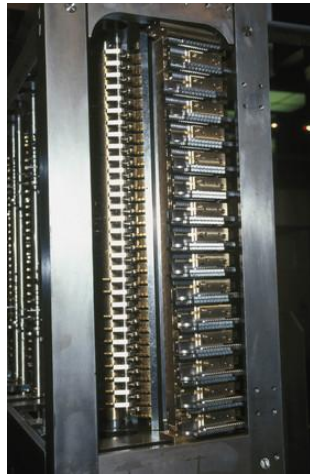
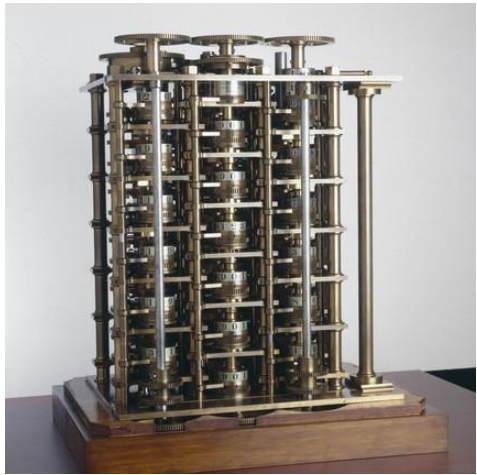
Babbage's Analytical Engine

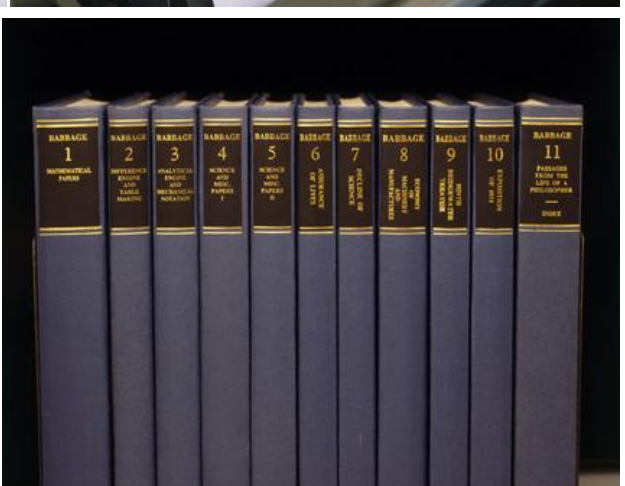
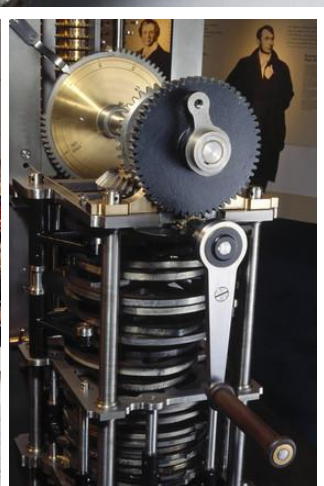
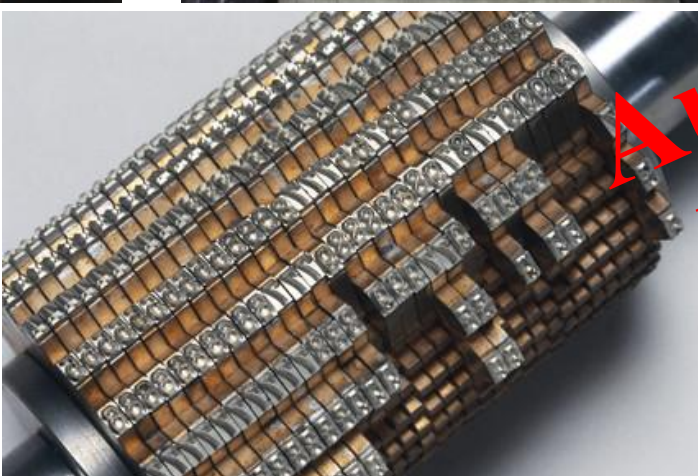
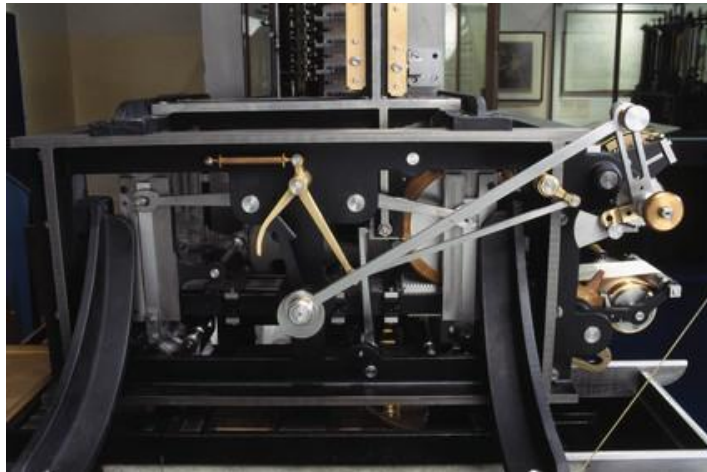
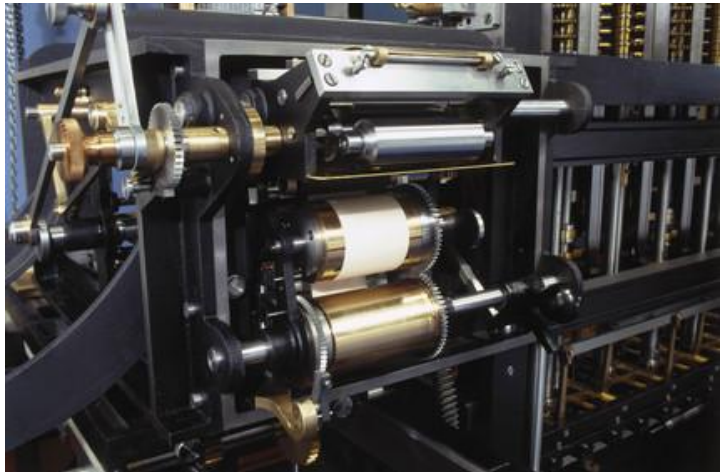
- World's **first general-purpose computer**
- Designed in **1837**, redesigned throughout Babbage's life
- **Turing-complete**, memory: 1000x50 digits (21 kB)
- **Fully programmable** "CPU", used punched cards
- Featured **ALU**, "**microcode**", **loops**, and **printer!**
- Could **multiply** two 20-digit numbers in **3 min**
- Few components built by Babbage; constructed in 1991

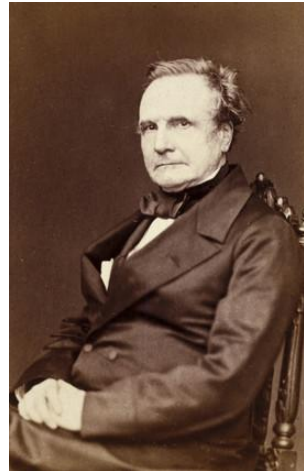
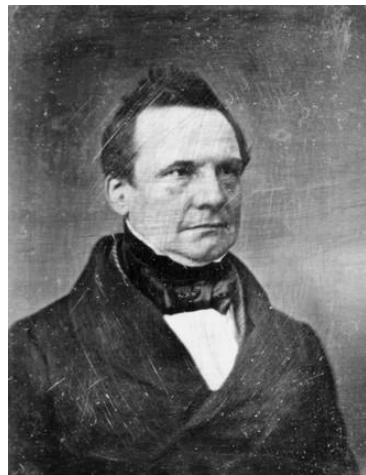
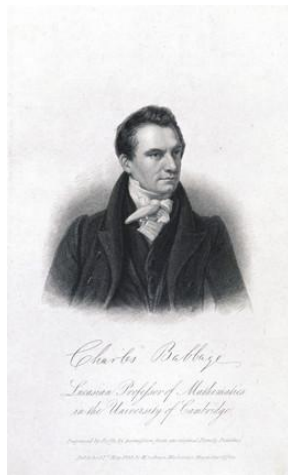


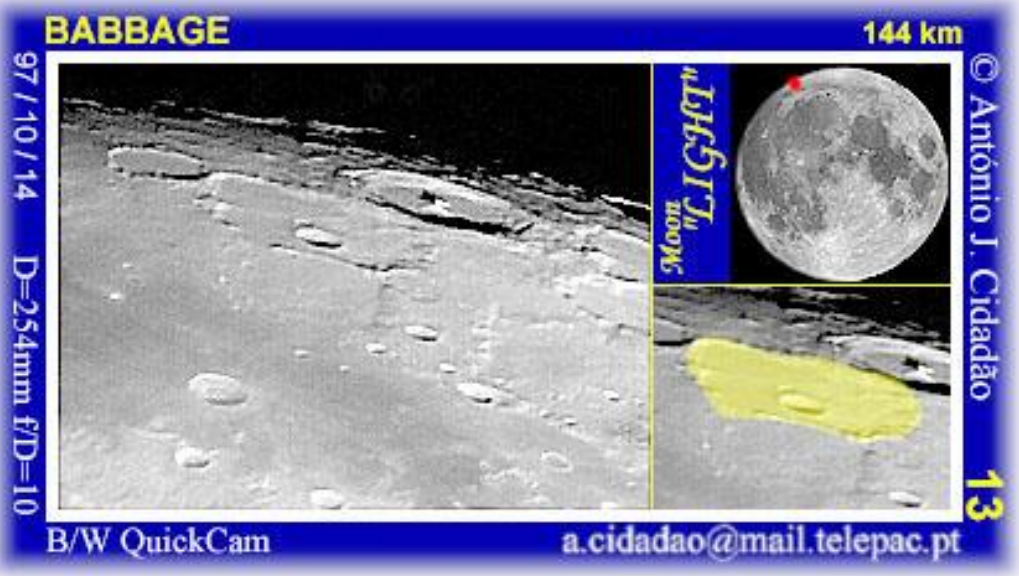
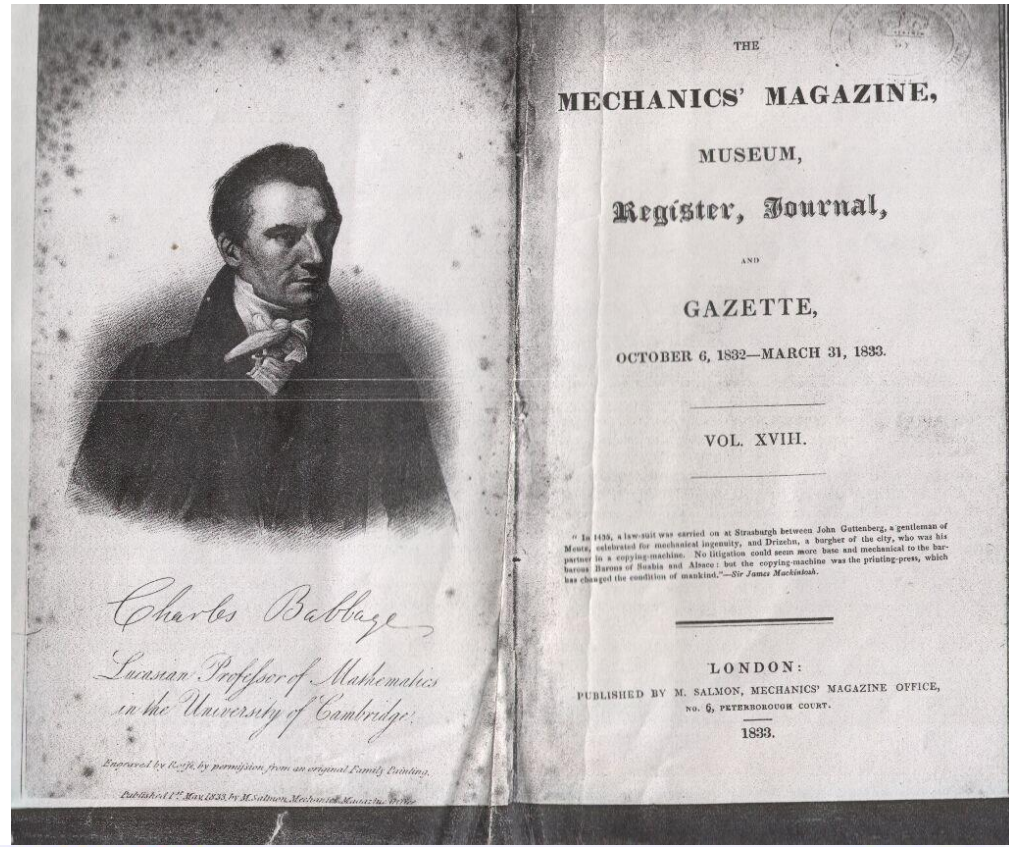
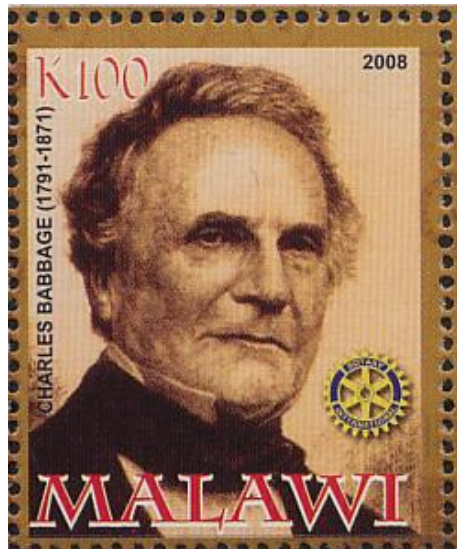












Antônio J. Cidadão 13

a.cidadao@mail.telepac.pt



CHARLES BABBAGE INSTITUTE

CENTER FOR THE HISTORY OF INFORMATION TECHNOLOGY



INSTITUTE OF TECHNOLOGY

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WELCOME TO THE CHARLES BABBAGE INSTITUTE

The Charles Babbage Institute (CBI) is an archives and research center dedicated to preserving the history of information technology and promoting and conducting research in the field.

Primary support for CBI is provided by the University of Minnesota, through the Institute of Technology and the University Libraries. Additional support is provided by corporate donors and individuals through the Friends of CBI.



SPOTLIGHT

- May 20th MHHC: IBM's Blue Gene
- New *CBI Newsletter* (Spring 2009, Vol. 31:1)
- McDonald Named 2009-2010 Tomash Fellow
- 2009 Norberg Travel Award Recipients

THE CBI ARCHIVES

The CBI Archives collects, preserves and provides access to rich archival collections and rare publications documenting the history of technology. Detailed [archival finding aids](#) are available. Researchers can also access digitized images ([Burroughs Corporation Image Database](#)) and one of the world's largest collections of research grade oral history interviews ([CBI Oral History Database](#)) through the CBI Web site. [More »](#)

THE CBI RESEARCH PROGRAM

CBI's historical research program identifies areas in which to collect archival materials; fosters new understanding of developments in the history of computing, software, and networking; supports the work of scholars outside the Institute (Tomash Fellowship and Norberg Travel Grant); and works collaboratively with individuals and organizations throughout the world. [More »](#)

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HAVE A QUESTION?

Ask a CBI archivist your questions about collections and services through instant message during regular business hours.

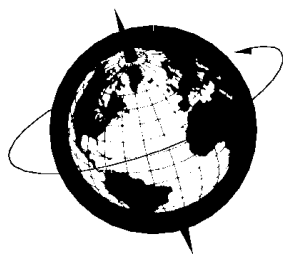
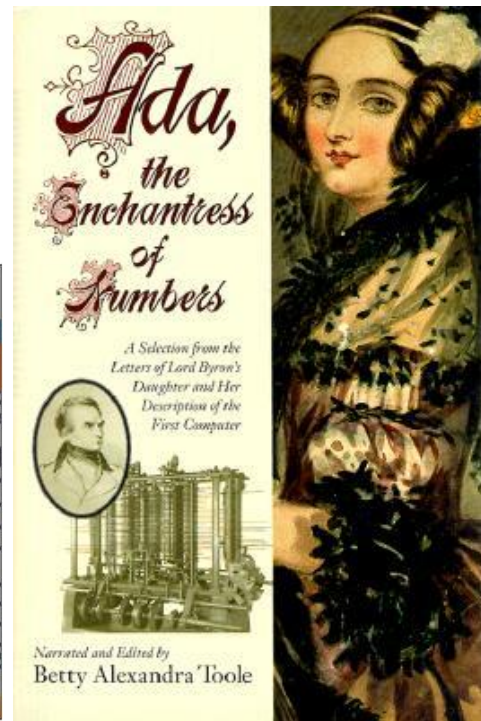
IM an Archivist

CBI Archivist is offline
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Historical Perspectives

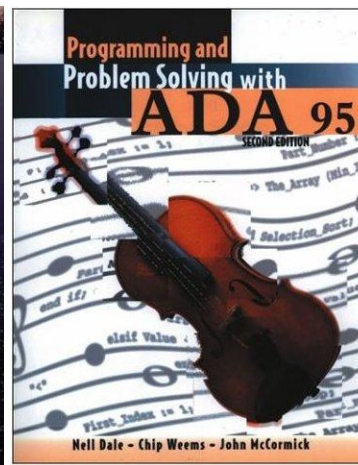
Countess Ada Lovelace (1815-1852)

- Daughter of Lord Byron
- Tutored in math and logic by De Morgan
- Wrote the “**manual**” for **Babbage’s analytical engine**, as well as **programs** for it
- **World’s first computer programmer!**
- **Foresaw** the vast **potential** of computers
- Babbage: “**The Enchantress of Numbers**”
- DoD’s **Ada language** “**MIL-STD-1815**”



Ada

*The International Language
for Software Engineering*





Ada Byron, Lady Lovelace
1815 - 1852



TILDA SWINTON TIMOTHY LEARY KAREN BLACK FRANCESCA FARIDANY JOHN PERRY BARLOW

CONCEIVING

Ada

A film by Eryn Hershman Leeson

"One of the Year's 10 Best!"
-B. Ruby Rich, San Francisco Bay Guardian

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IN THIS ISSUE



Will IBM buy Sun?

If IBM buys Sun Microsystems how will the diverse product portfolios fit together?

NEWS ANALYSIS 12

OGC 'secret' out

The Office of Government Commerce finally publishes two ID card Gateway reviews

NEWS 8

Tech terms banned

IT professionals react with hostility to a list of words council leaders want to ban

NEWS ANALYSIS 10

Beware of SaaS risk

The cost benefits of software-as-a-service should not blind companies to potential hazards

NEWS ANALYSIS 14

Web past to present

We celebrate 20 years of the internet by looking back at key events in its development

THIS WEEK ON THE WEB 20

Leadership lessons

CW500 Club president shares his insights on challenges and opportunities facing IT leaders

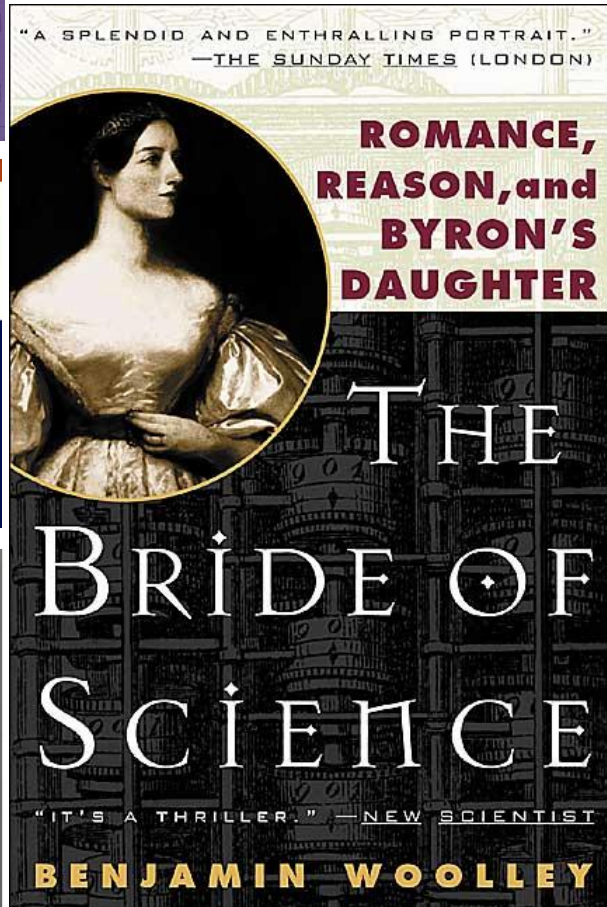
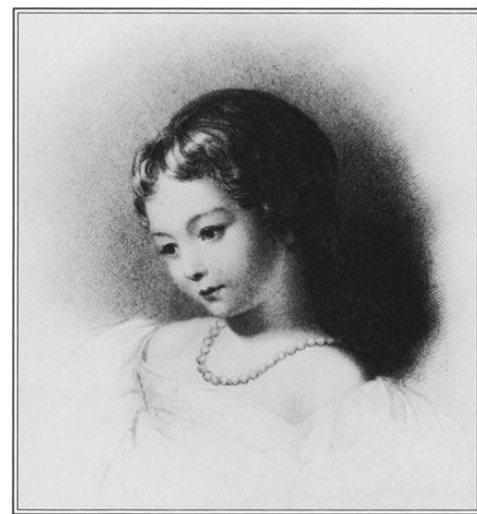
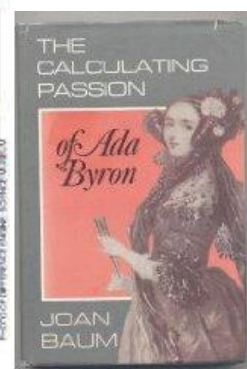
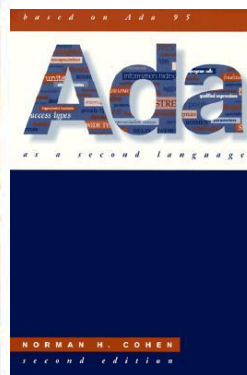
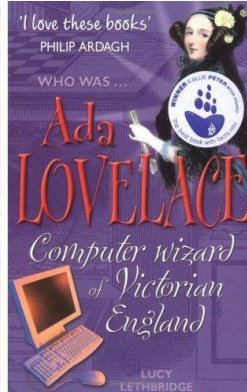
STRATEGY 22



Female role models in IT

ADA LOVELACE DAY AIMS TO RAISE AWARENESS OF WOMEN'S ACHIEVEMENTS IN THE TECHNOLOGY SECTOR PAGE 24

LATEST JOBS
IT VACANCIES
START ON
PAGE 23



"A SPLENDID AND ENTHRALLING PORTRAIT."
—THE SUNDAY TIMES (LONDON)

ROMANCE,
REASON, and
BYRON'S
DAUGHTER

THE
BRIDE OF
SCIENCE

"IT'S A THRILLER."
—NEW SCIENTIST

BENJAMIN WOOLLEY



Lovelace Medal

Lovelace Medal Home

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Lovelace Medal

The Lovelace Medal is presented to individuals who have made a contribution which is of major significance in the advancement of Information Systems or which adds significantly to the understanding of Information Systems.

About the medal

Lovelace Medal 2009

2009 winner

The winner of the 2009 Lovelace Medal is Professor Yorick Wilks.

Previous Lectures

Video: A tribute to Karen Spärck Jones

The 2008 BCS Lovelace Medal lecture was a very special event dedicated to the memory of Karen Spärck Jones who was presented the award just weeks before she died last year. The lecture was delivered by Dr Ann Copestake and is now available to watch online.



2007 Lovelace Lecture - Sir Tim Berners-Lee

The Web is a technical and social creation, dependent on both technical protocols and social conventions. The origins and potential futures of this large scale, emergent phenomena were discussed by Sir Tim Berners-Lee in this year's BCS Lovelace Lecture - now available to watch via this website.



Previous winners

[Previous winners](#) of the Lovelace Medal have included:
2008 - Dr Tony Storey
2007 - Karen Sparck-Jones
2006 - Sir Tim Berners-Lee
2005 - Dr Nicholas McKeown

Ada Lovelace notes on “Sketch of the Analytical Engine Invented by Charles Babbage”, by L. F. Menabrea, 1843

Her notes (three times longer than the paper itself!) contain the world’s first computer program (for calculating Bernoulli numbers):

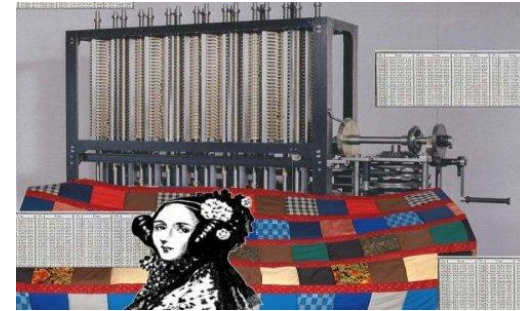
Number of Operations Nature of Operations		Variables for Data						Working Variables									Variables for Results	
		¹ V ₀	¹ V ₁	¹ V ₂	¹ V ₃	¹ V ₄	¹ V ₅	⁰ V ₆	⁰ V ₇	⁰ V ₈	⁰ V ₉	⁰ V ₁₀	⁰ V ₁₁	⁰ V ₁₂	⁰ V ₁₃	⁰ V ₁₄	⁰ V ₁₅	⁰ V ₁₆
		+	+	+	+	+	+	+	+	+	+	+	+	+	+	+	+	+
		0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
		0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
		0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
		0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
		m	n	d	m'	n'	d'										$\frac{dn' - d'n}{mn' - m'n} = x$	$\frac{d'm - dm'}{mn' - m'n} = y$
1	×	m	n'	mn'										
2	×	n	m'	$m'n$										
3	×	d	dn'									
4	×	0	d'	$d'n$									
5	×	0	0	$d'm$								
6	×	0	0	dm'							
7	-	0	0	$(mn' - m'n)$						
8	-	0	0	$(dn' - d'n)$					
9	-	0	0	$(d'm - dm')$				
10	÷	$(mn' - m'n)$	0	$\frac{dn' - d'n}{mn' - m'n} = x$		
11	÷	0	0	$\frac{d'm - dm'}{mn' - m'n} = y$	

Quotes from the Ada Lovelace notes on

“Sketch of the Analytical Engine Invented by Charles Babbage”, 1843

“We may say most aptly, that the Analytical Engine *weaves algebraical patterns* just as the Jacquard-loom weaves flowers and leaves.”

“Again, it might act upon *other things besides number*, were objects found whose mutual fundamental relations could be expressed by those of the *abstract science of operations*, and which should be also susceptible of adaptations to the action of the operating *notation* and mechanism of the engine. Supposing, for instance, that the fundamental relations of pitched sounds in the science of harmony and of musical composition were susceptible of such expression and adaptations, the engine might compose elaborate and scientific pieces of *music of any degree of complexity or extent.*”



Quotes from the Ada Lovelace notes on

“Sketch of the Analytical Engine Invented by Charles Babbage”, 1843

“Many persons who are not conversant with mathematical studies, imagine that because the business of the engine is to give its results in *numerical notation*, the *nature of its processes* must consequently be *arithmetical* and *numerical*, rather than *algebraical* and *analytical*. This is an error. The engine can **arrange and combine** its numerical quantities exactly **as if they were *letters* or any other *general symbols***; and in fact it might bring out its results in algebraical *notation*, were provisions made accordingly.”

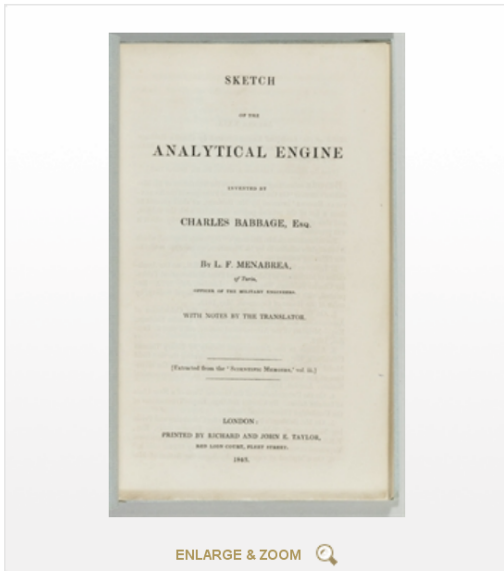
“But it would be a **mistake** to suppose that because its *results* are given in the *notation* of a more restricted science, its *processes* are therefore **restricted to those of that science**. The object of the engine is in fact to give the *utmost practical efficiency* to the resources of *numerical interpretations* of the **higher science of analysis**, while it uses the processes and combinations of this latter.”



LOT 21 / SALE 2013

EMAIL PRINT

[BABBAGE]. -- MENABREA, Luigi Federico (1809-1896). *Sketch of the Analytical Engine invented by Charles Babbage... with notes by the translator.* Offprint from: *Scientific Memoirs*. Translated by Augusta Ada King, Countess of Lovelace (1809-1896). Volume 3. London: Richard and John E. Taylor, 1843.



ENLARGE & ZOOM

Price Realized (Set Currency)
\$170,500

Price includes buyer's premium

Estimate
\$10,000 - \$15,000

Sale Information

Sale 2013
Important Scientific Books: The
Richard Green Library
17 June 2008
New York, Rockefeller Plaza

Lot Description

[BABBAGE]. -- MENABREA, Luigi Federico (1809-1896). *Sketch of the Analytical Engine invented by Charles Babbage... with notes by the translator.* Offprint from: *Scientific Memoirs*. Translated by Augusta Ada King, Countess of Lovelace (1809-1896). Volume 3. London: Richard and John E. Taylor, 1843.

LOTS IN THIS SALE

NEW YORK, ROCKEFELLER PLAZA | 17 JUNE 2008

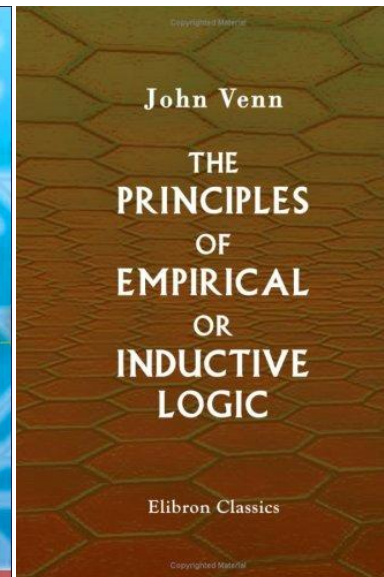
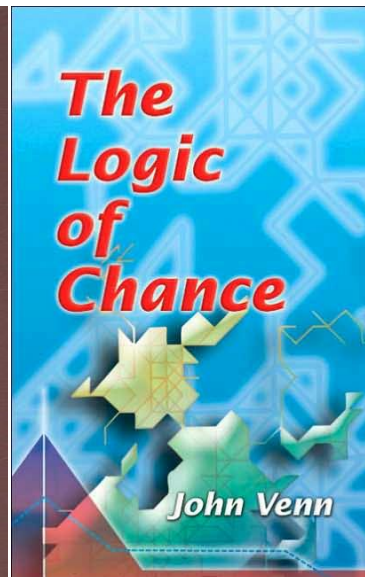
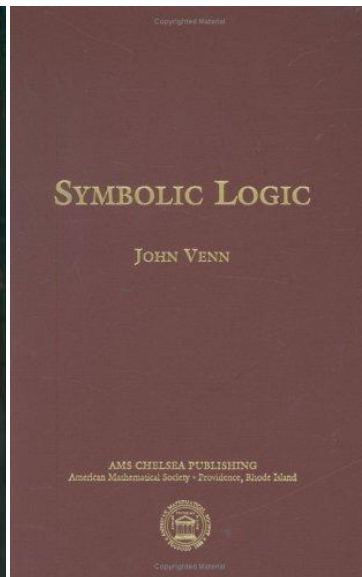
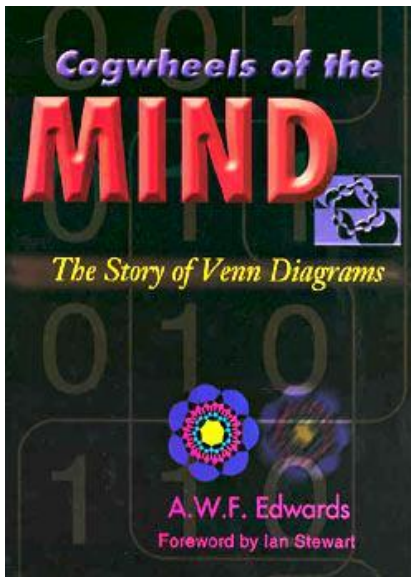
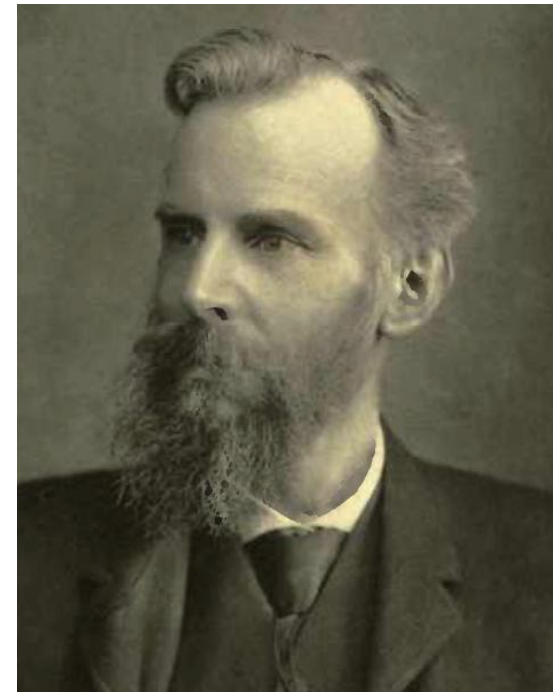
Important Scientific Books:
The Richard Green Library

- LOT #21
[BABBAGE]. -- MENABREA, Luigi Federico...
- LOT #22
BABBAGE, Charles. *Passages from the...*
- LOT #23
BABBAGE, Charles. *The Ninth...*
- LOT #24
[BALLISTICS]. *Une merveille du génie...*
- LOT #25
[BALLISTICS]. BRITISH INFORMATION...
- LOT #26
[BALLISTICS]. *The United States...*
- LOT #27
BAYER, Johann (1572-1625)....
- LOT #28
BEAUMONT, William (1785-1853)....
- LOT #30

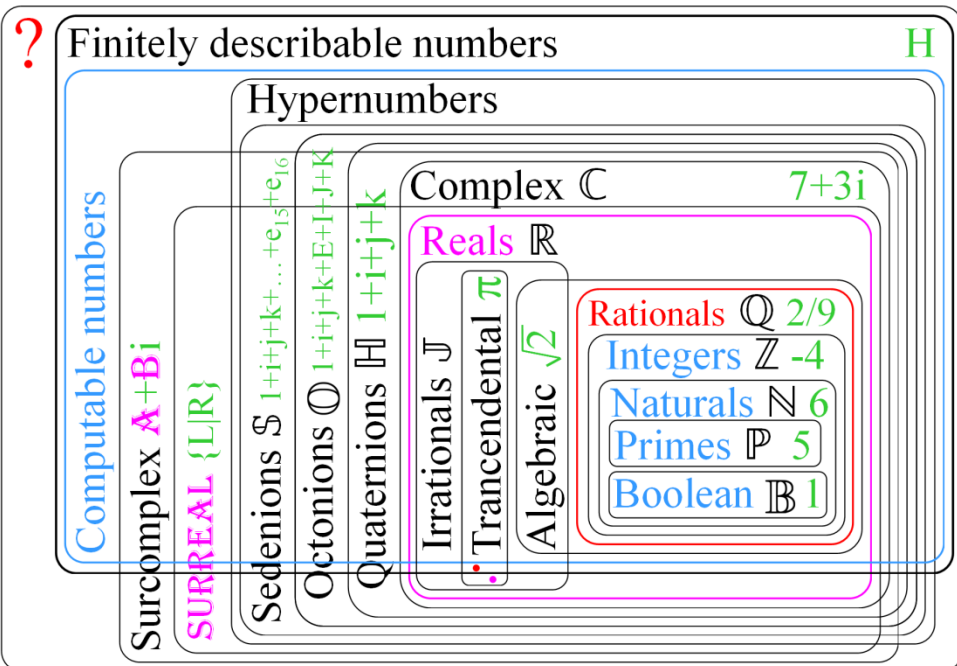
Historical Perspectives

John Venn (1834-1923)

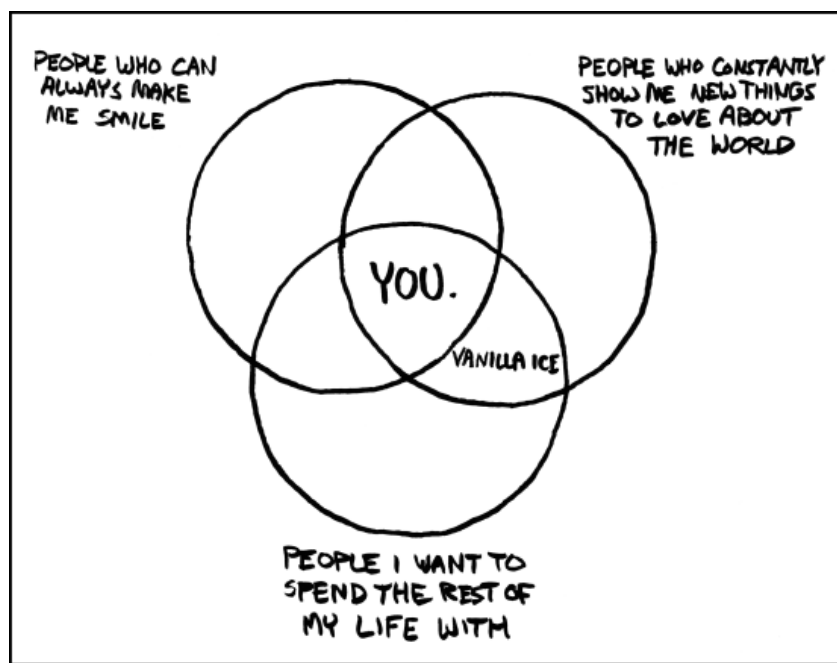
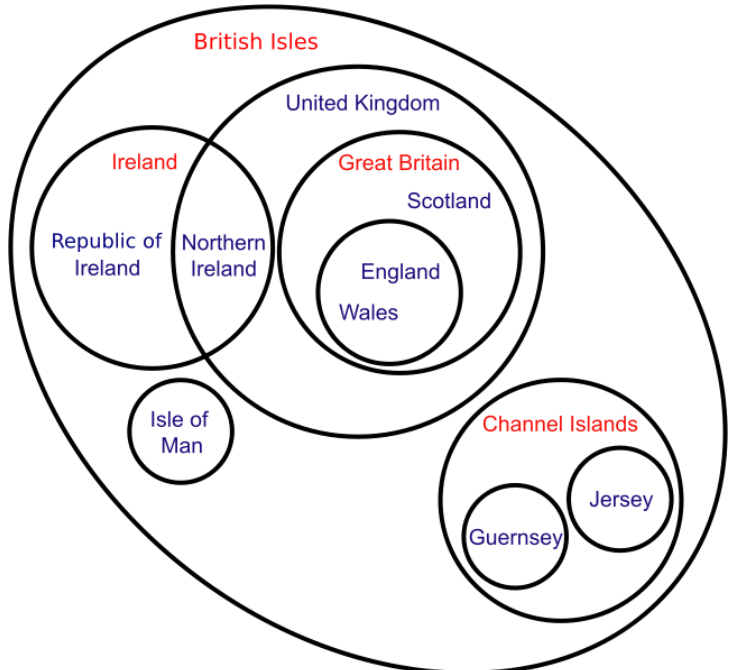
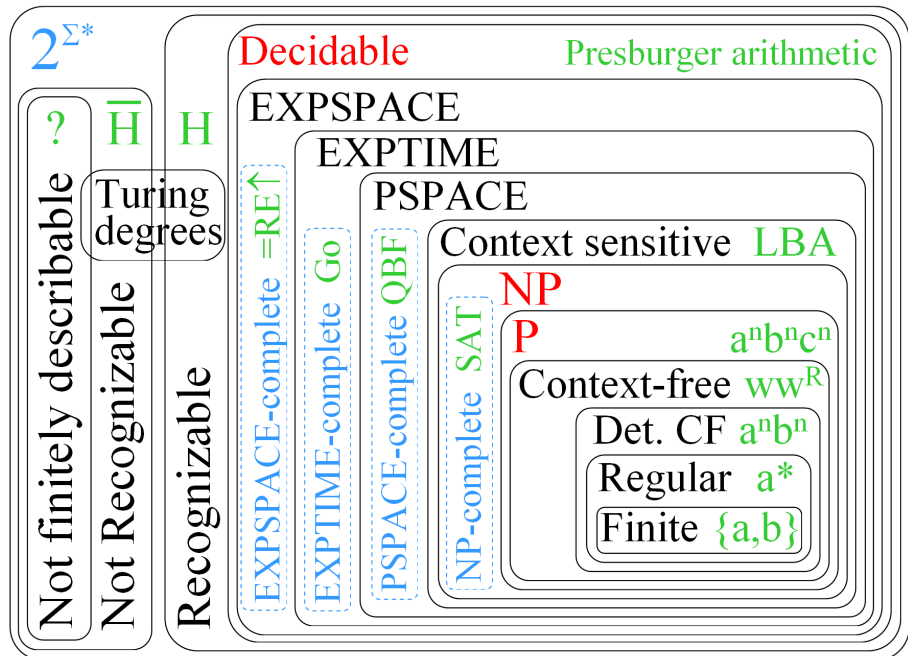
- Logician and philosopher
- Worked in logic, probability, set theory
- Introduced the “**Venn diagram**” (1880)
 - Very widely used, **many applications**
 - **Ties together** fundamental concepts from logic, geometry, combinatorics, knot theory

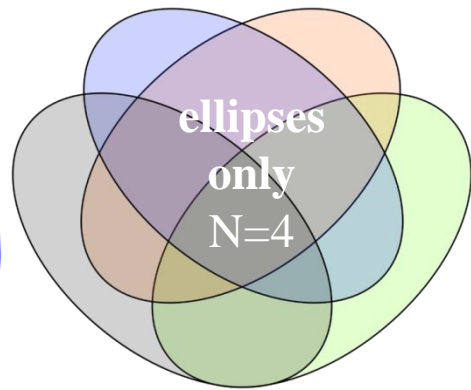
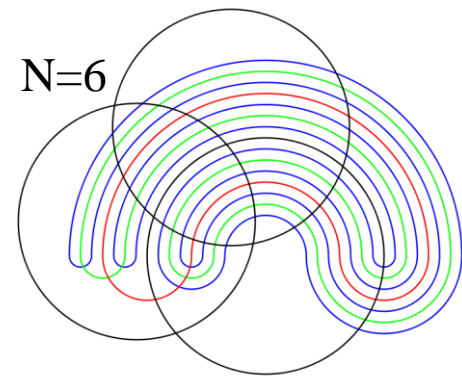
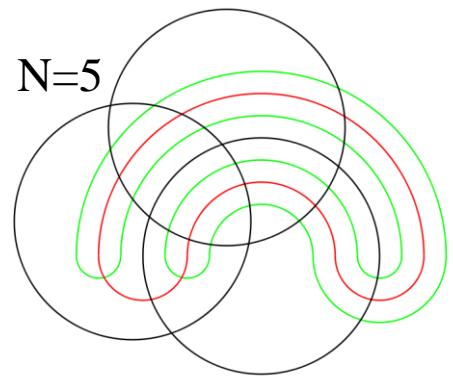
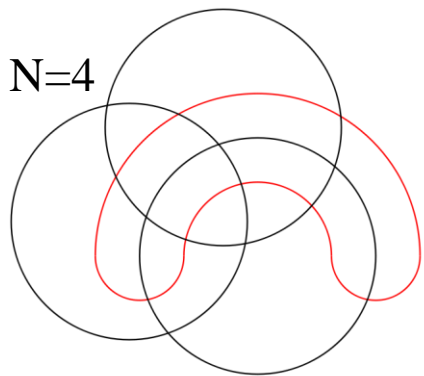


Generalized Numbers

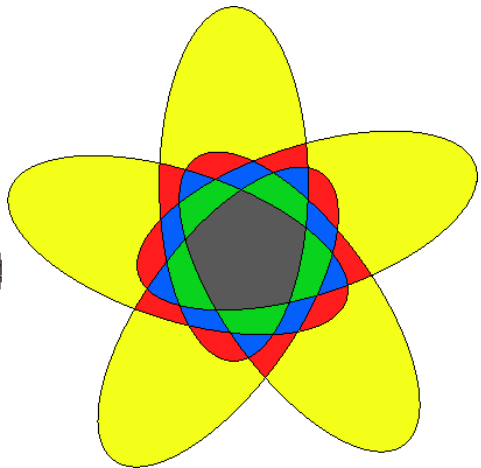
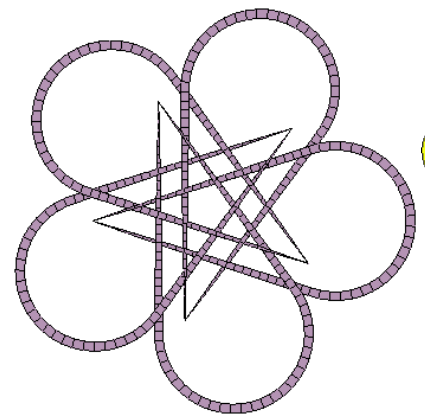
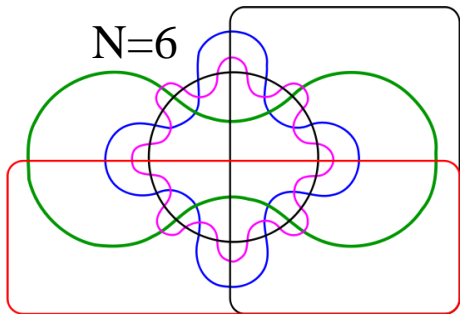
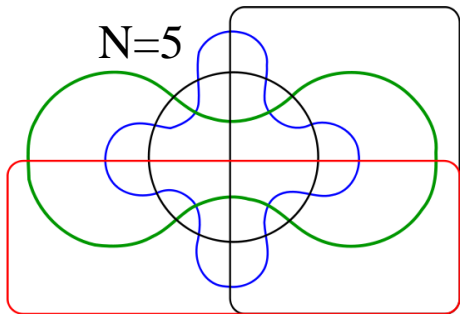
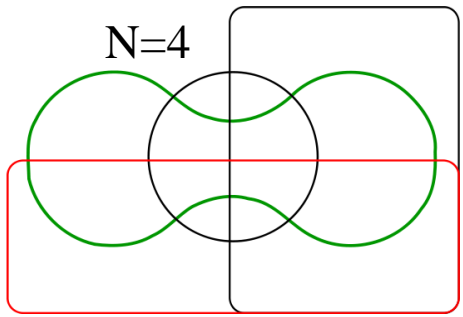
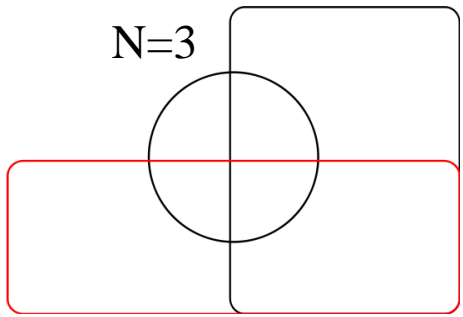


The Extended Chomsky Hierarchy

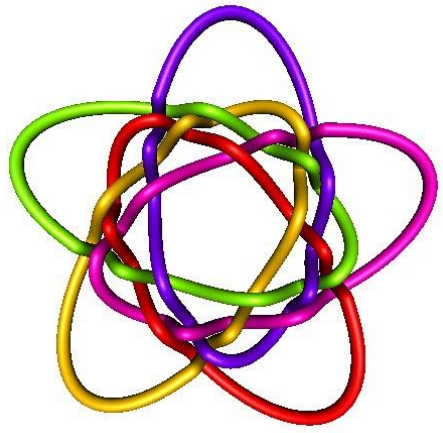




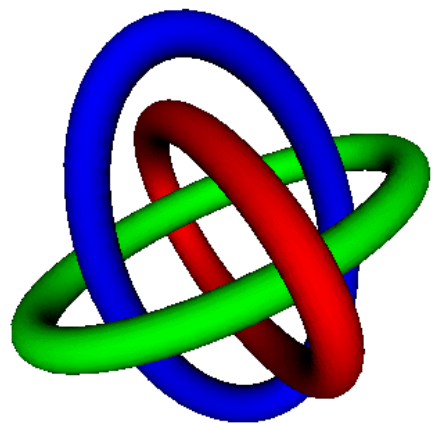
Generalized Venn diagrams [John Venn, 1880]



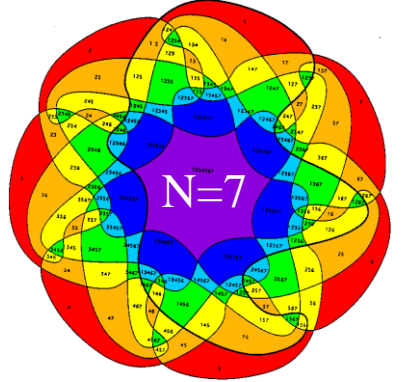
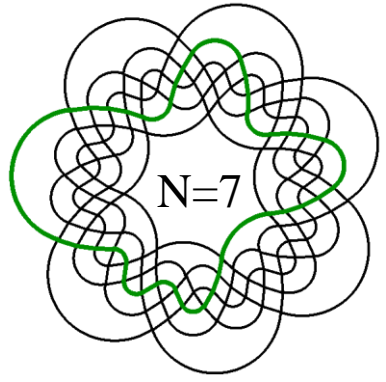
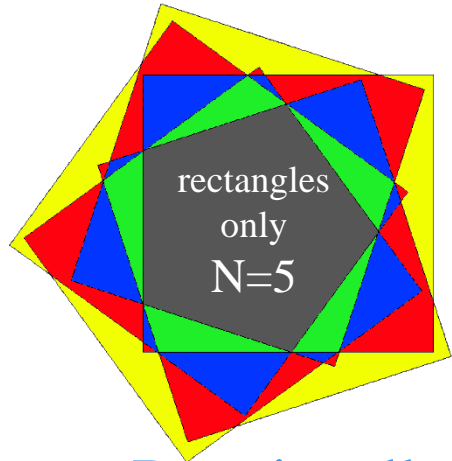
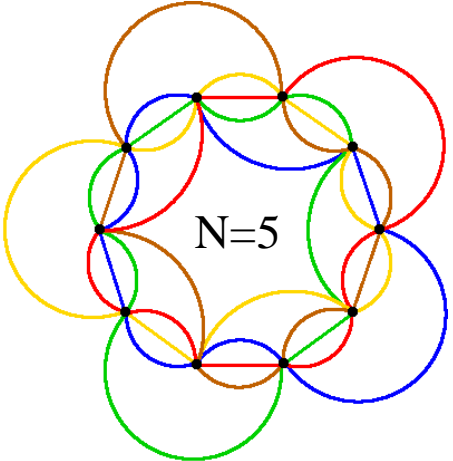
Ellipses only
N=5



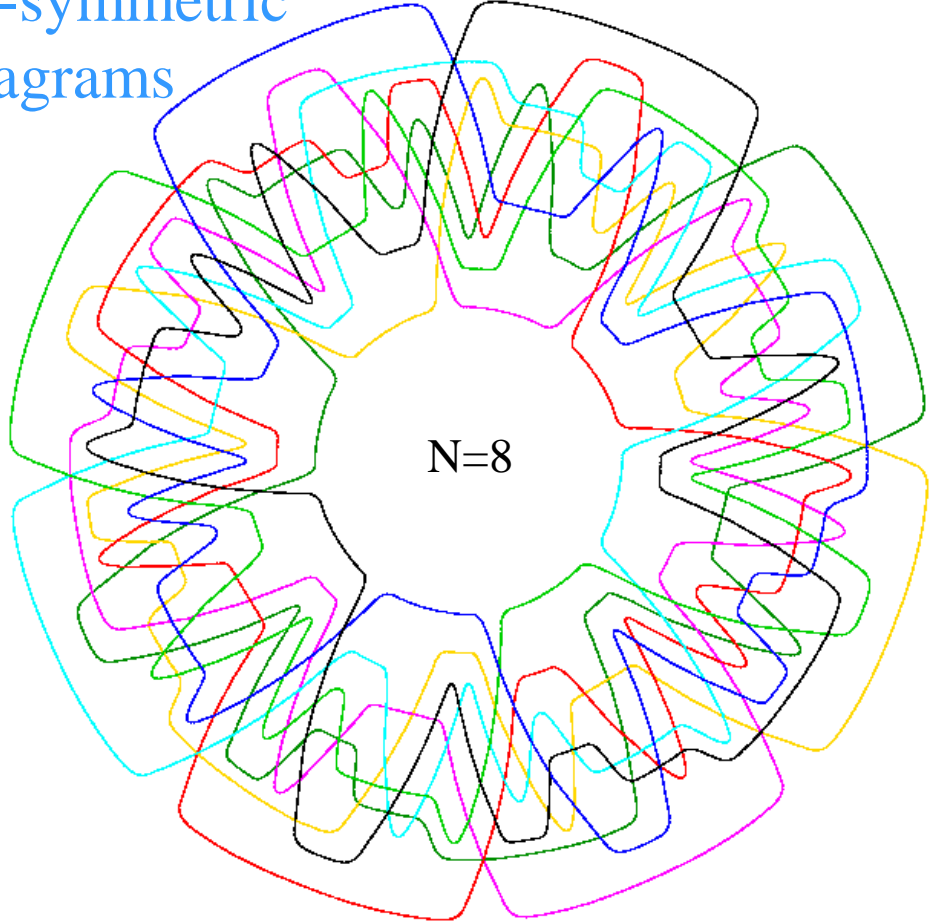
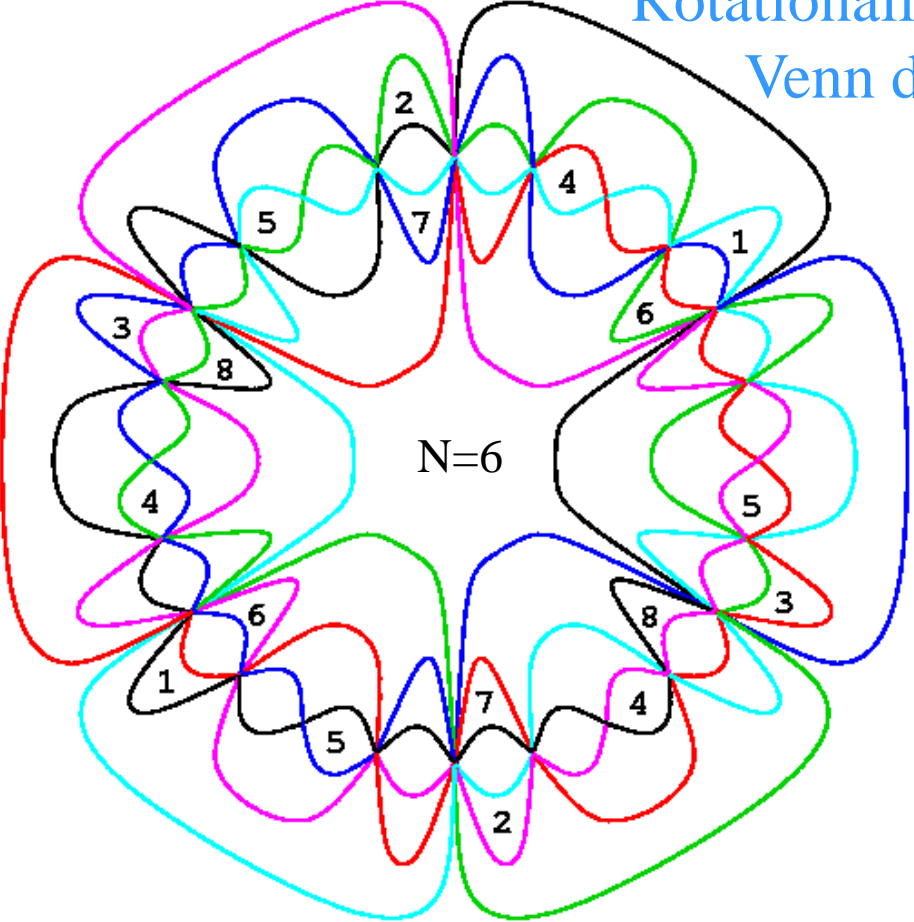
Borromean rings
analogue N=5

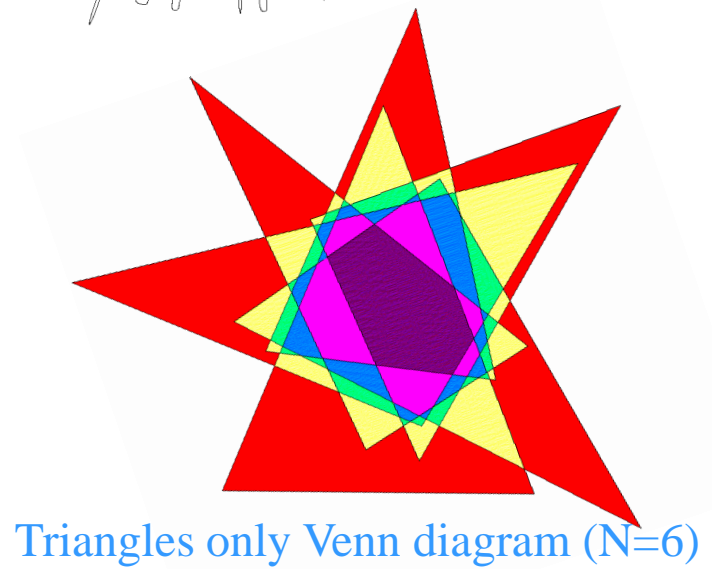
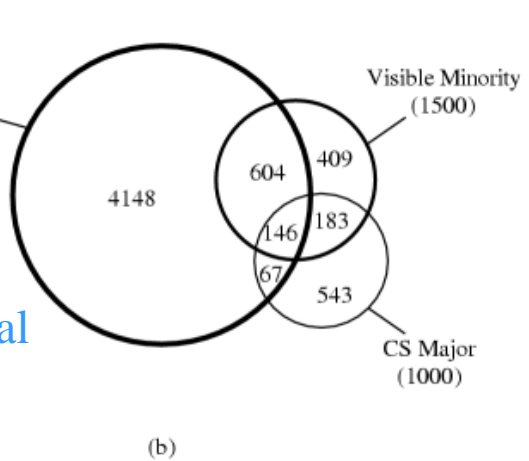
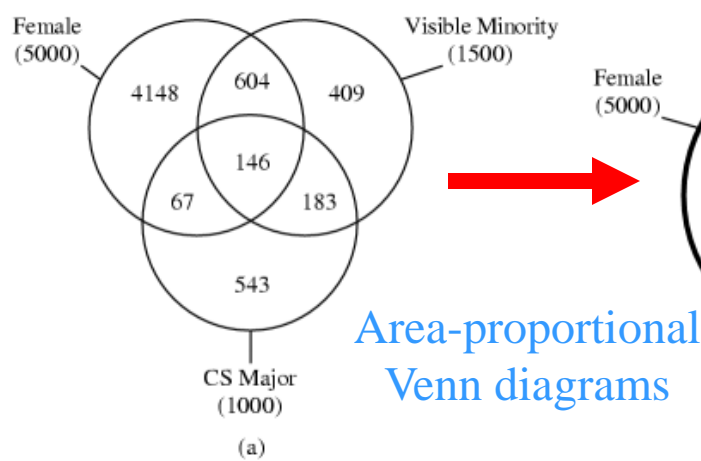
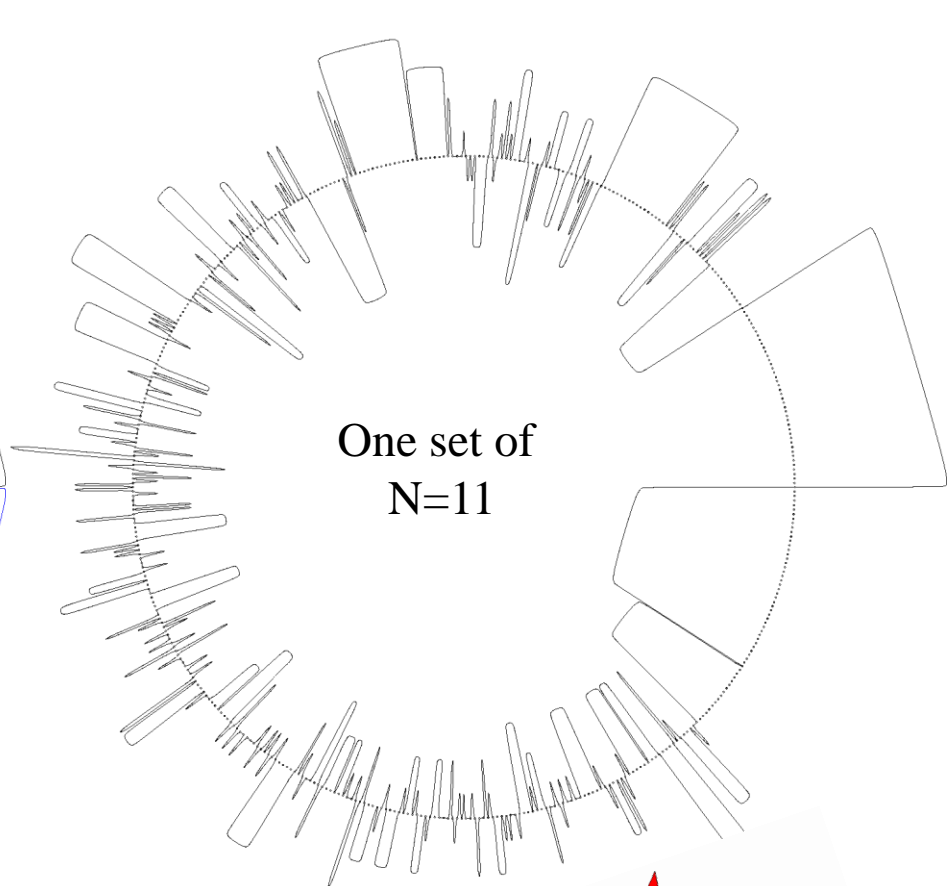
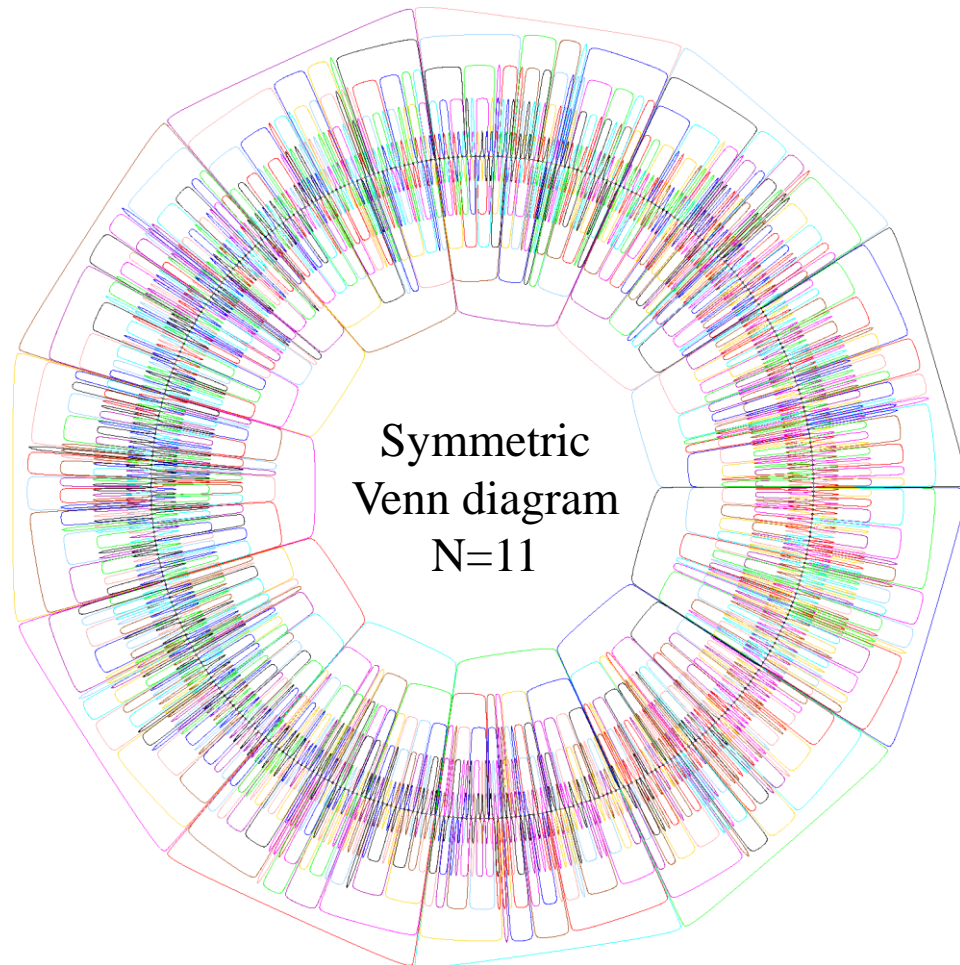


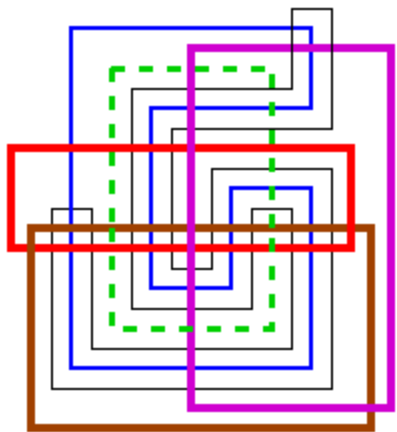
Borromean rings
N=3



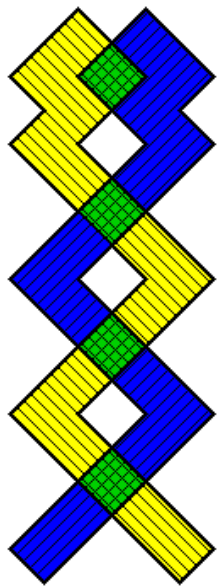
Rotationally-symmetric
Venn diagrams



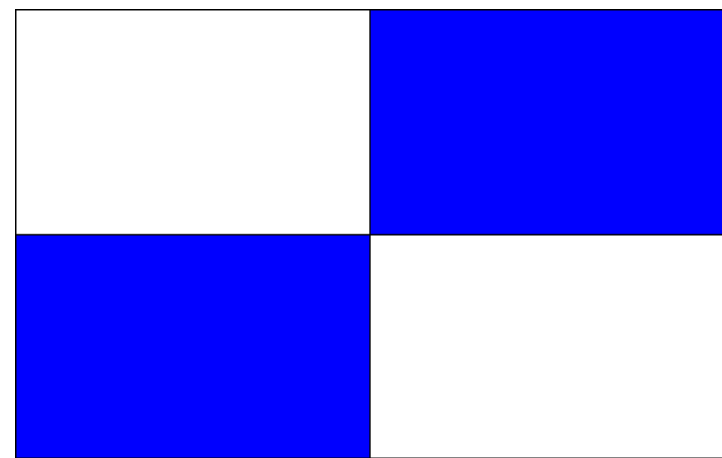
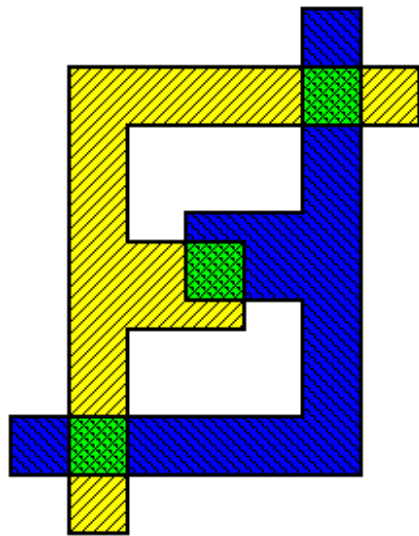




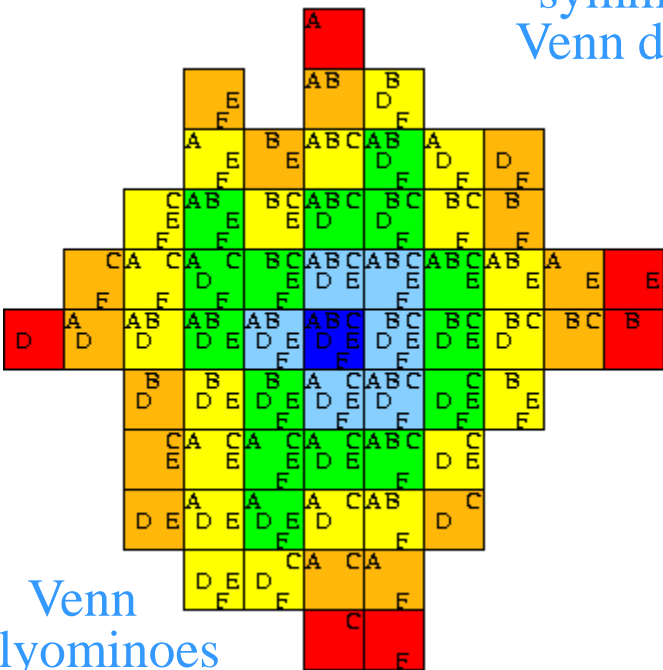
exposed Venn diagrams $n=5$



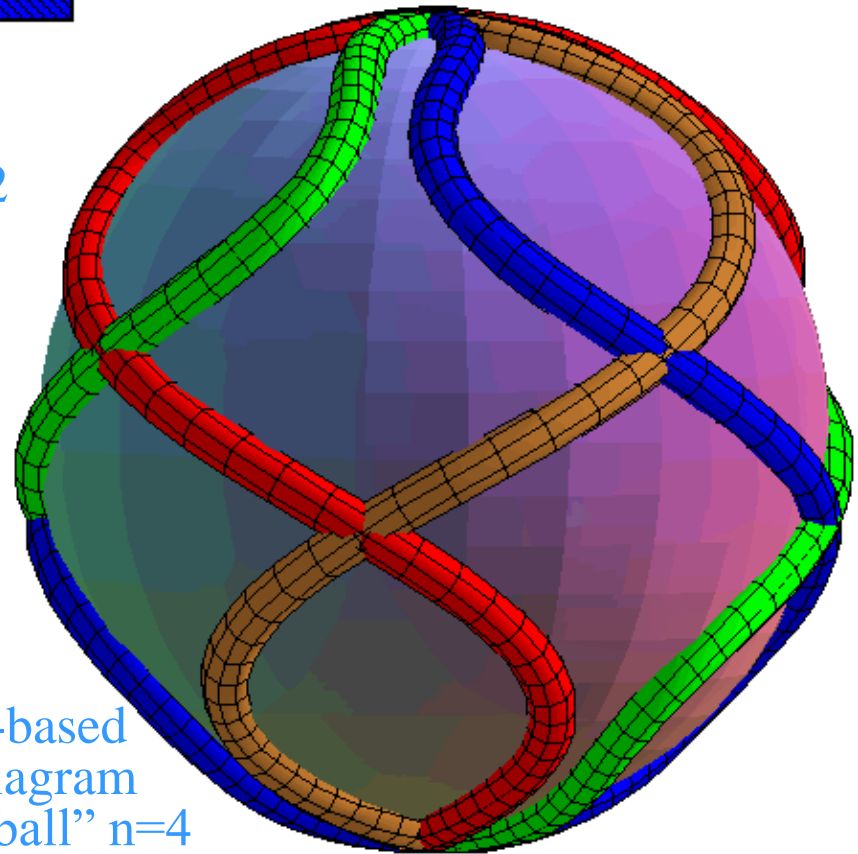
symmetric k -fold Venn diagrams $n=2$



$x = 0.$

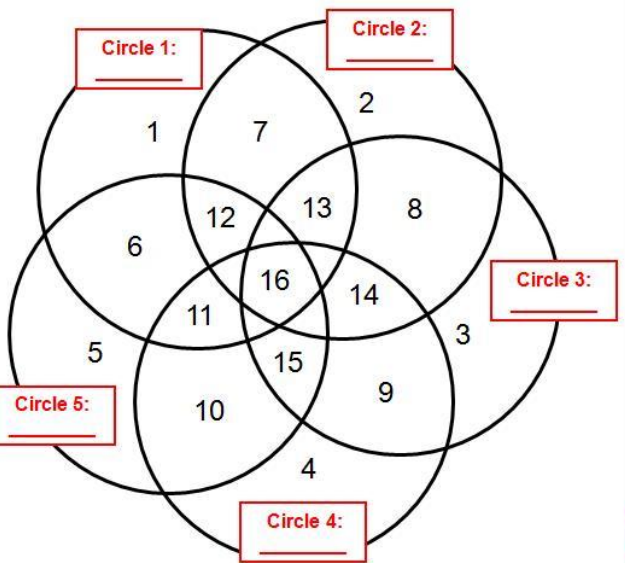


Venn polyominoes

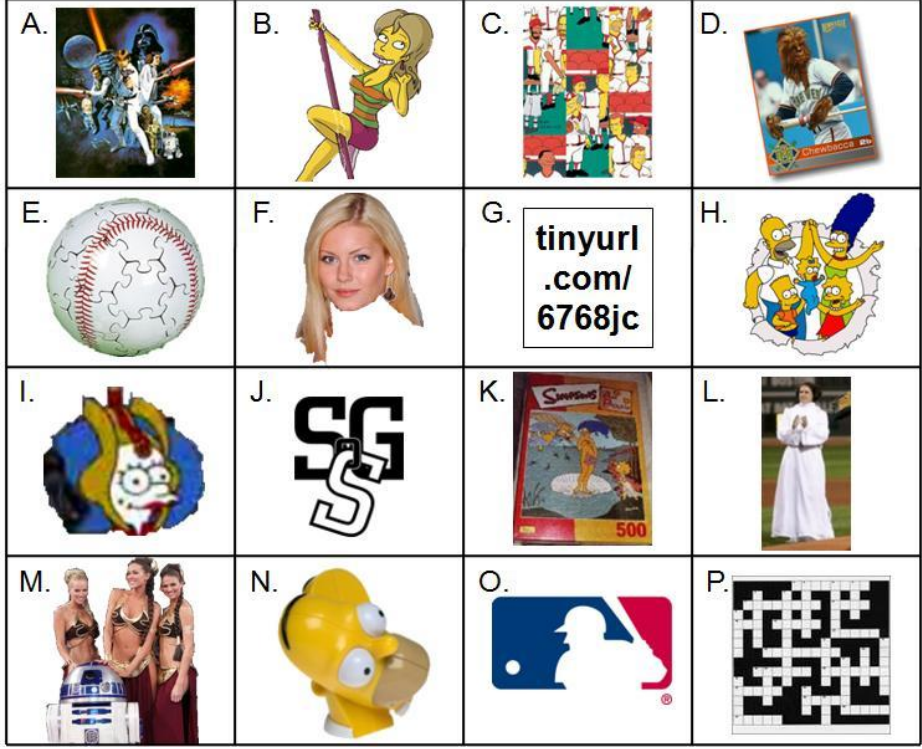


sphere-based Venn diagram "Vennice ball" $n=4$

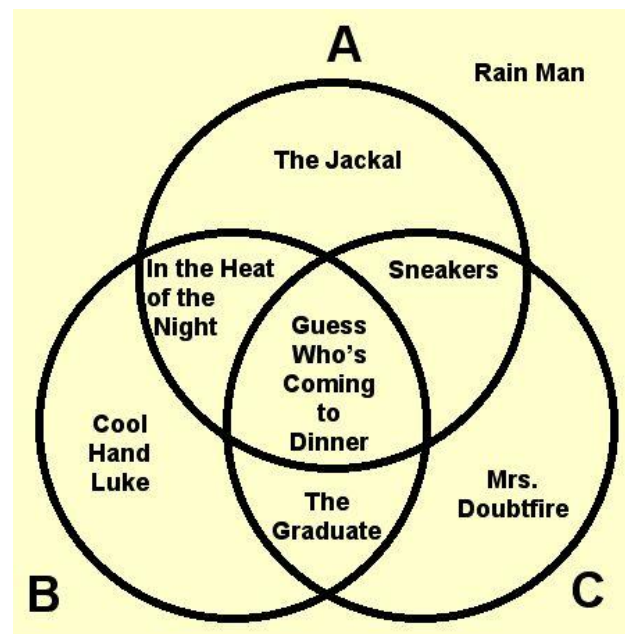
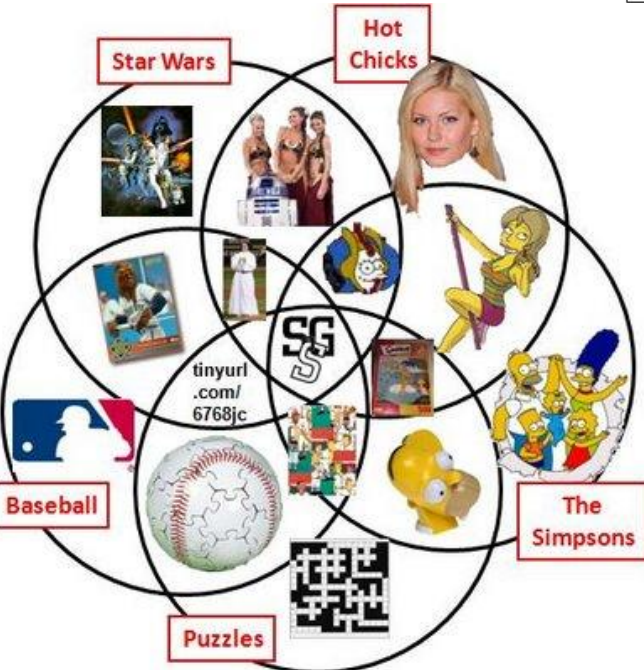
Venn diagram puzzles:



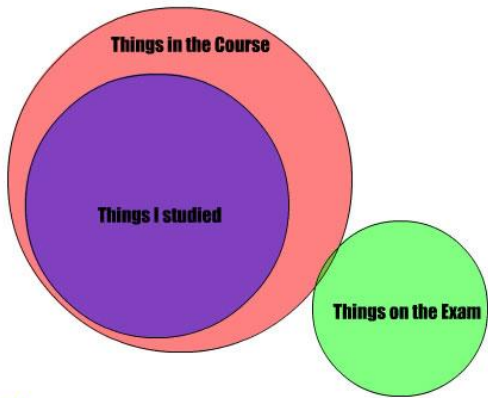
- Answer Panel:**
1. A
 2. ?
 3. ?
 4. ?
 5. ?
 6. ?
 7. ?
 8. ?
 9. ?
 10. ?
 11. ?
 12. ?
 13. ?
 14. ?
 15. ?
 16. ?
- Circle 1: ?
 Circle 2: ?
 Circle 3: ?
 Circle 4: ?
 Circle 5: ?



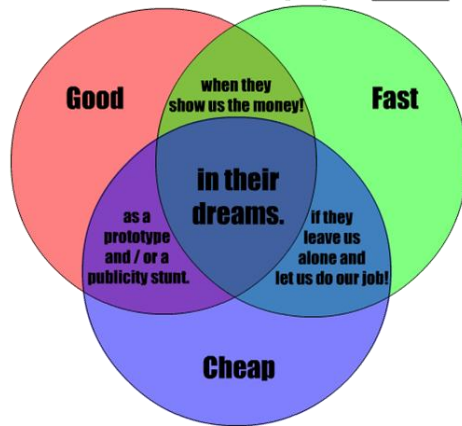
Puzzle solution:



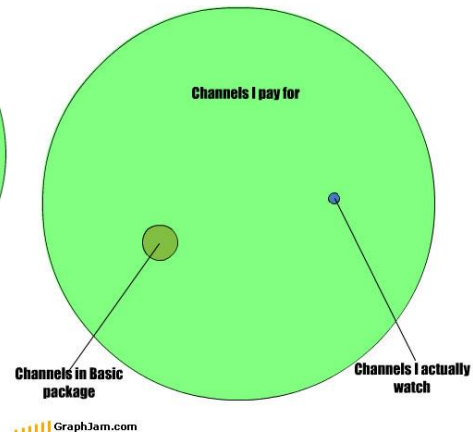
Final Exams



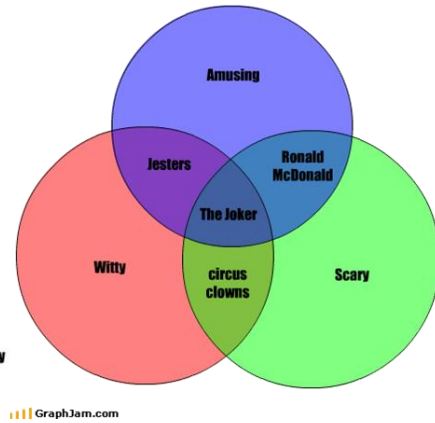
A client can have their project _____:



Cable TV



Types of clowns



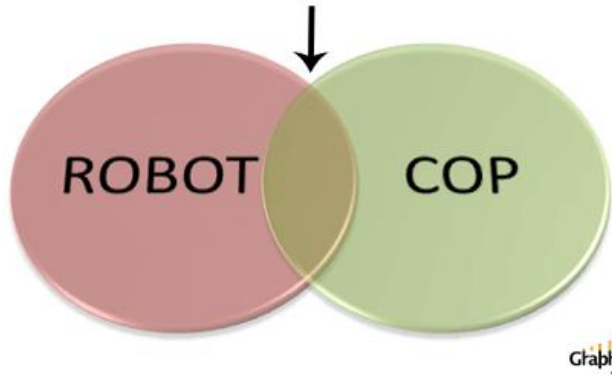
GraphJam.com

GraphJam.com

GraphJam.com

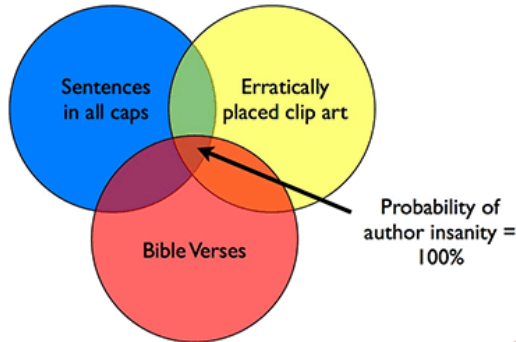
GraphJam.com

Futuristic Trends in Law Enforcement



GraphJam

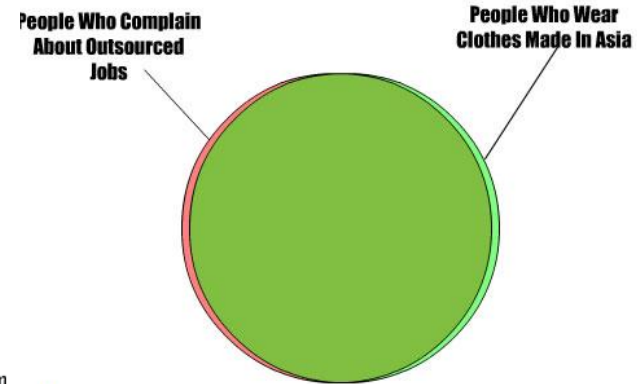
Judging Web Site Author Sanity



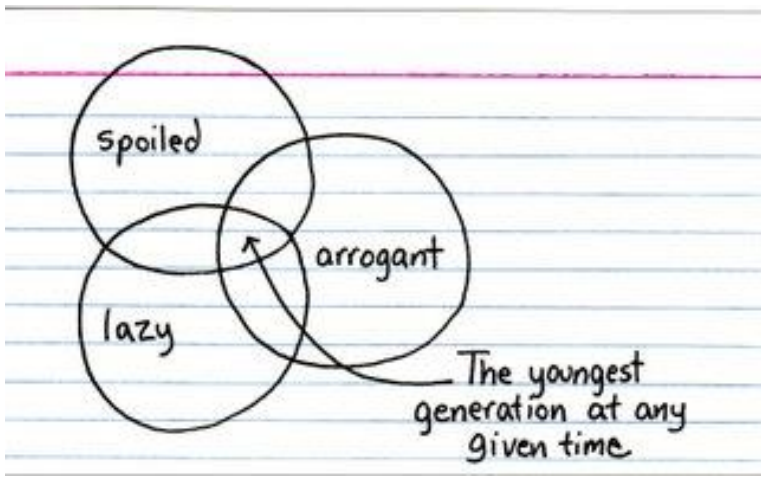
GraphJam

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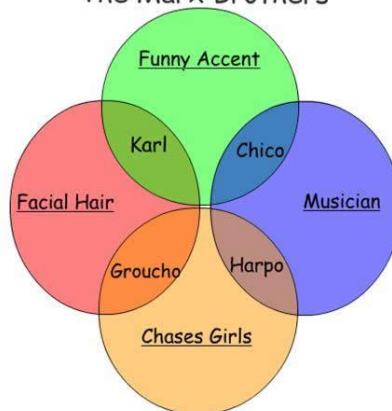
Clothing of Complainers



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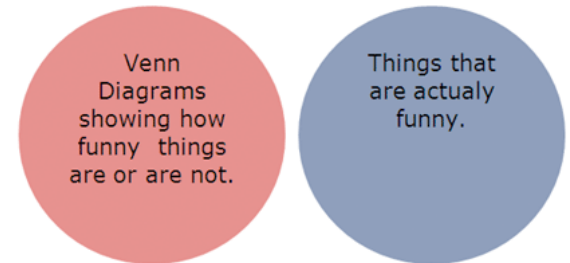


The Marx Brothers

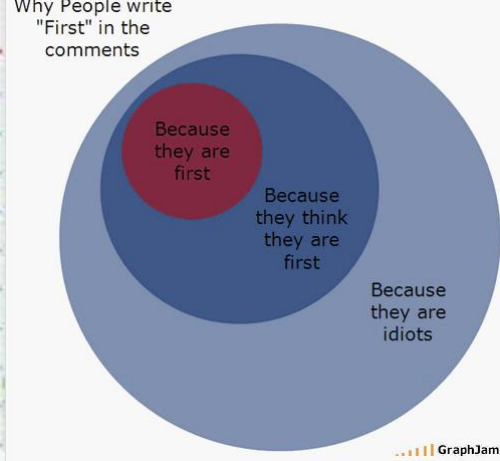
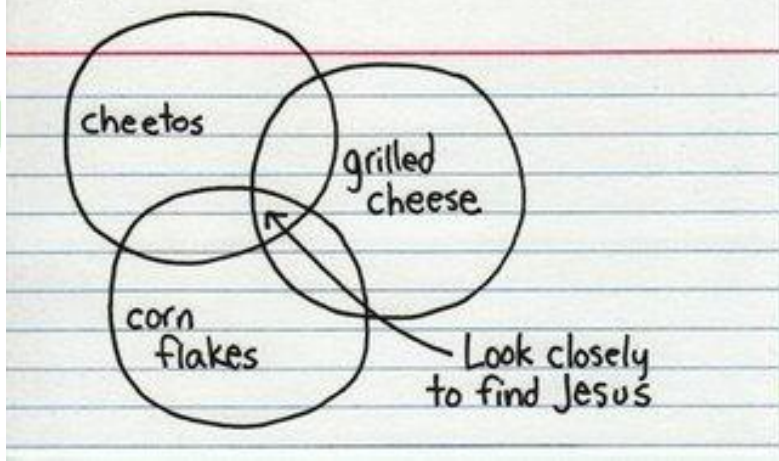
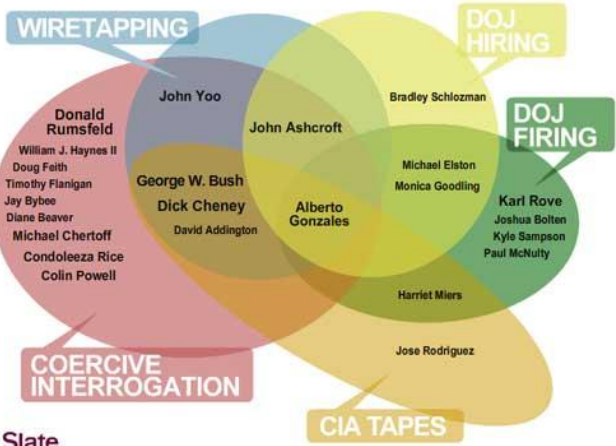


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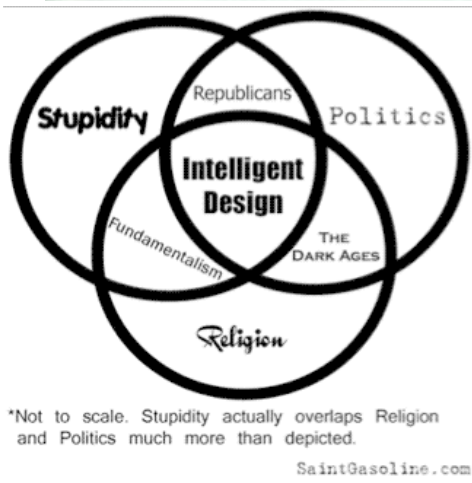
The Ironic Truth about Venn Diagrams



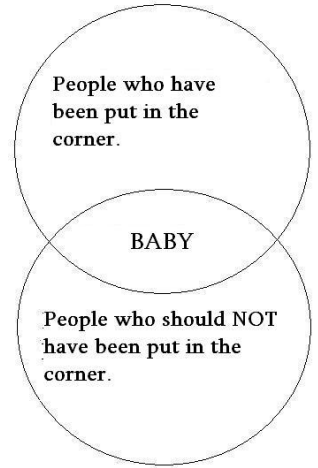
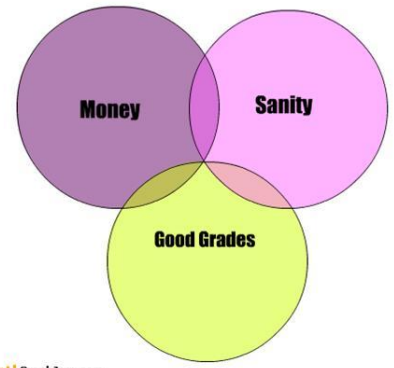
GraphJam



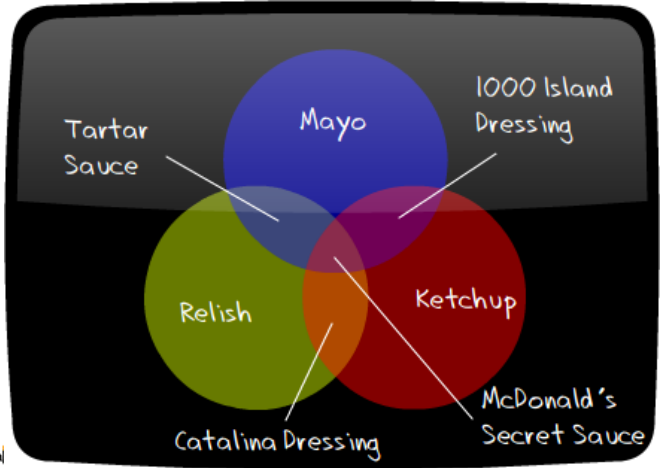
Slate

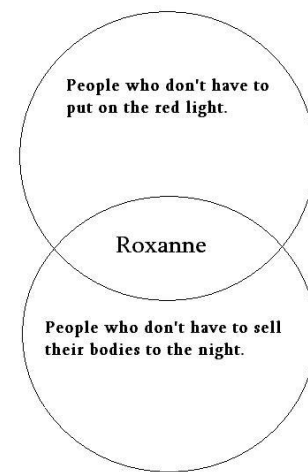
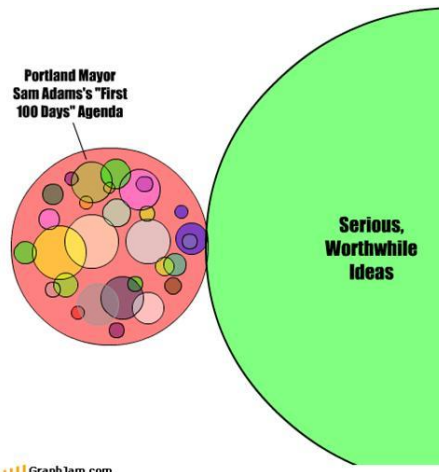
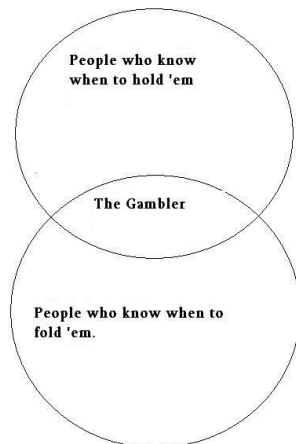
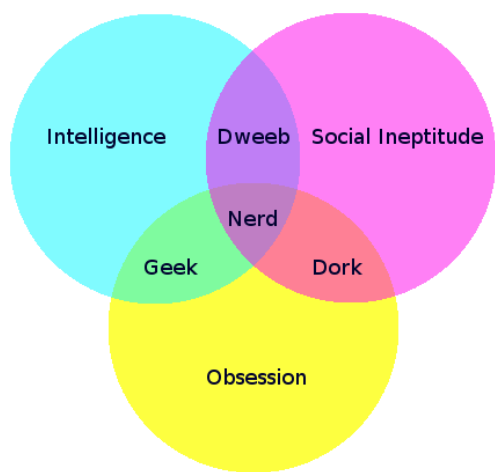


Things One Can Have While in College



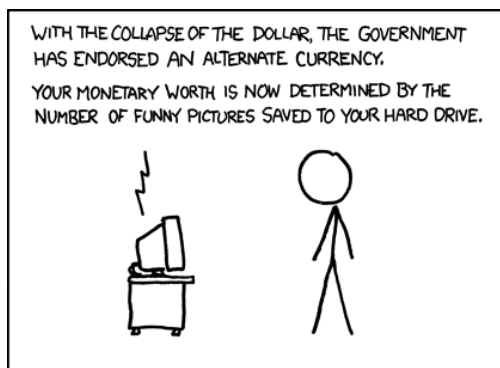
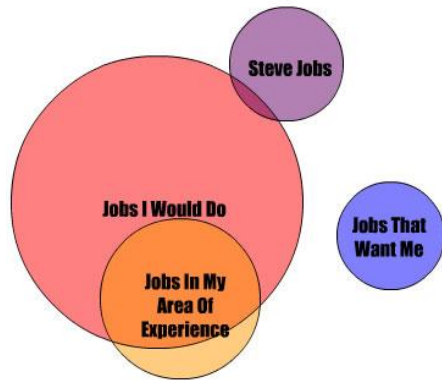
Correlation Between Visual Subgroups in Select Juveniles





GraphJam.com

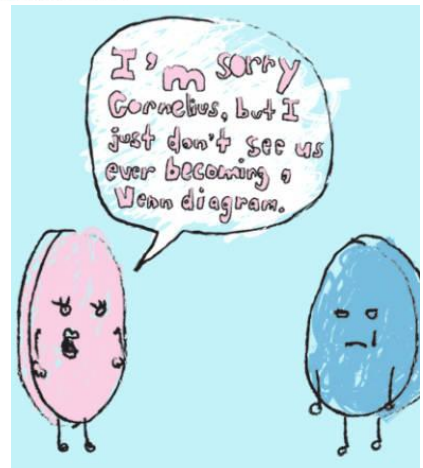
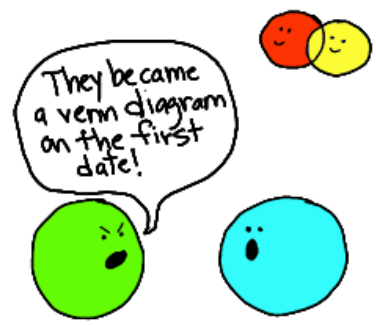
My Job Search



WITH THE COLLAPSE OF THE DOLLAR, THE GOVERNMENT HAS ENDORSED AN ALTERNATE CURRENCY. YOUR MONETARY WORTH IS NOW DETERMINED BY THE NUMBER OF FUNNY PICTURES SAVED TO YOUR HARD DRIVE.

I HAVE BEEN PREPARING FOR THIS MOMENT MY WHOLE LIFE.

GraphJam.com



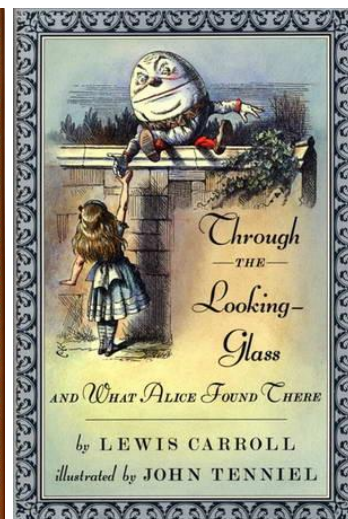
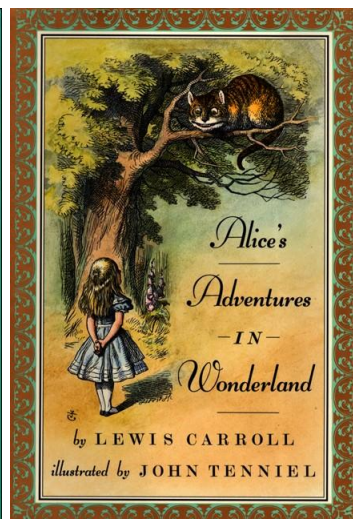
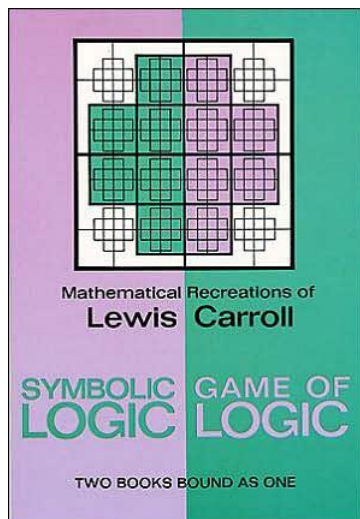
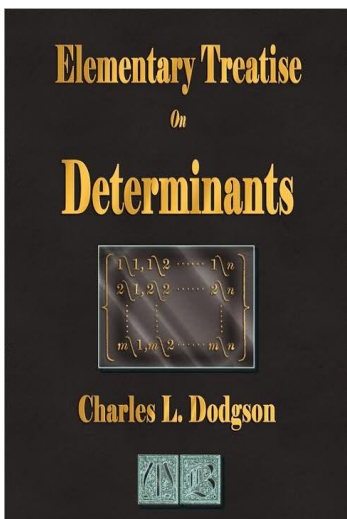
Harold had to face the painful truth. He and Daisy were never going to be a Venn diagram.



Historical Perspectives

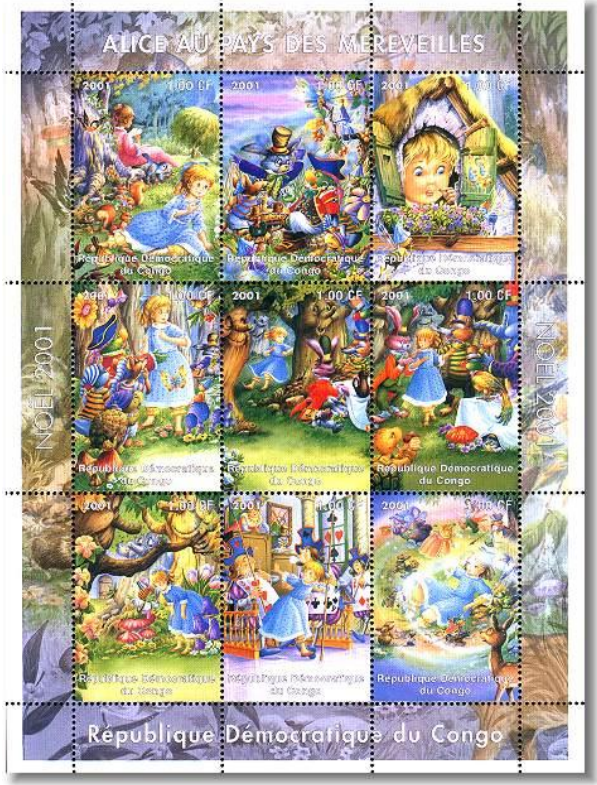
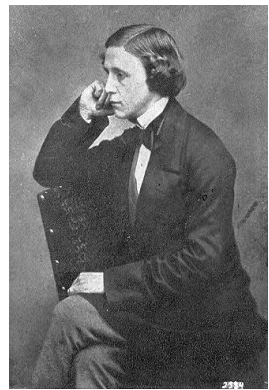
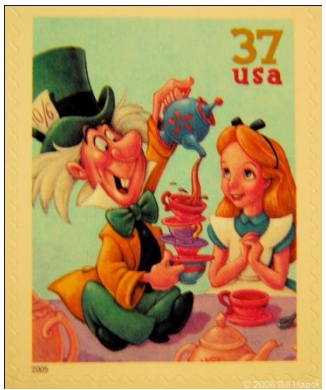
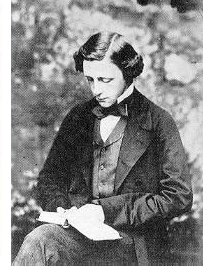
Charles Dodgson (1832-1898)

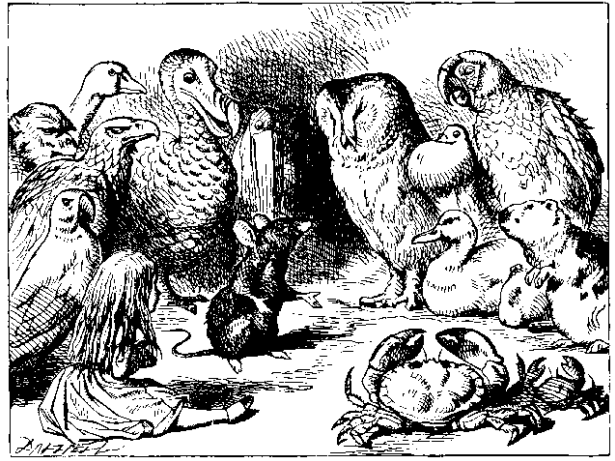
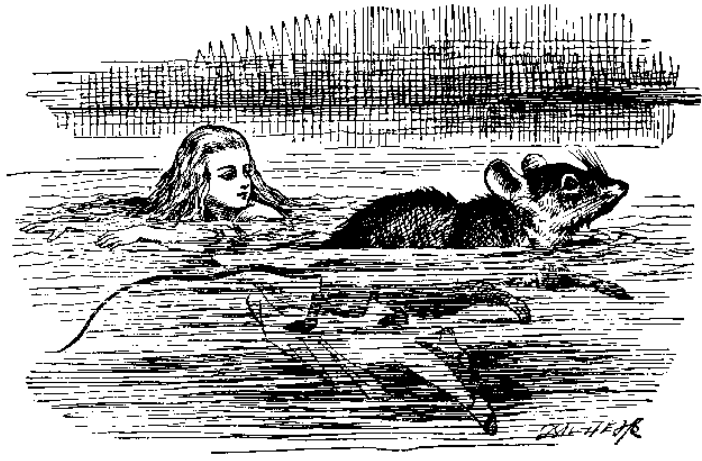
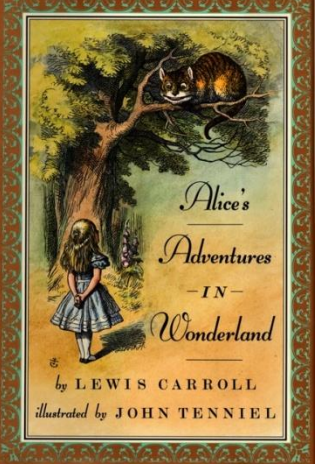
- AKA “**Lewis Carroll**”
- Mathematician, logician, author, photographer
- Wrote “**Alice in Wonderland**”, “**Jabberwocky**”, and “**Through the Looking Glass**”
- Popularized logic & syllogisms and made it fun!
- Invented “**Scrabble**” and “**word ladder**” games
- Profoundly influenced literature, art, and culture

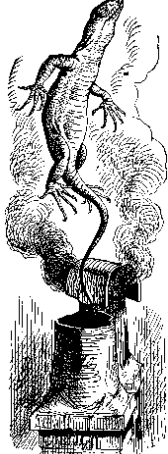


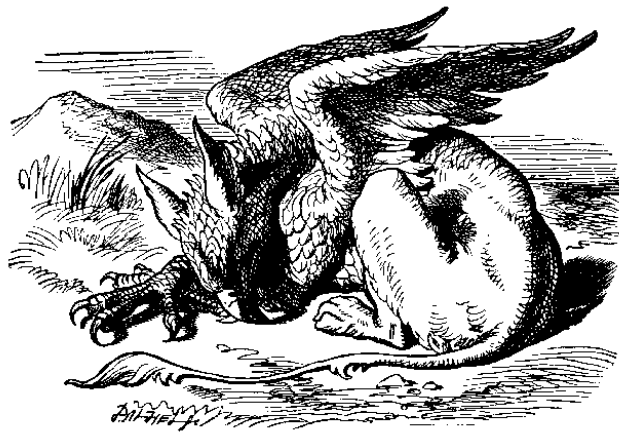
The **Disney** Classic Fairytales
in Postage Stamps

Alice in Wonderland





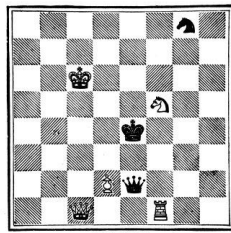
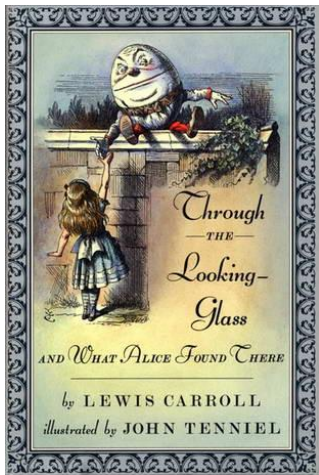






Beware
the Jabberwock
my son, the jaws
that bite, the
claws that catch,
Beware the jubjub
bird and shun
the frumious
Bandersnatch.

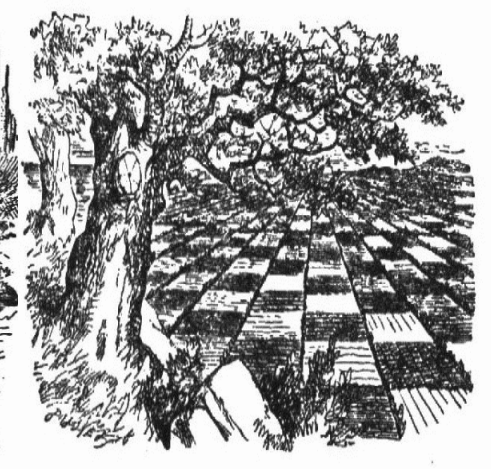
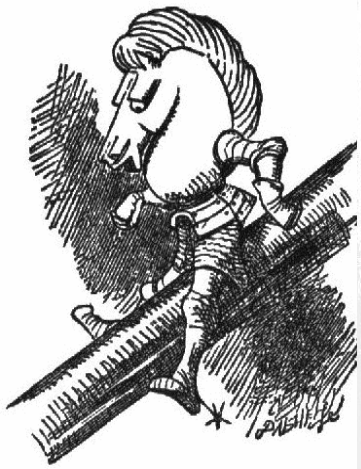
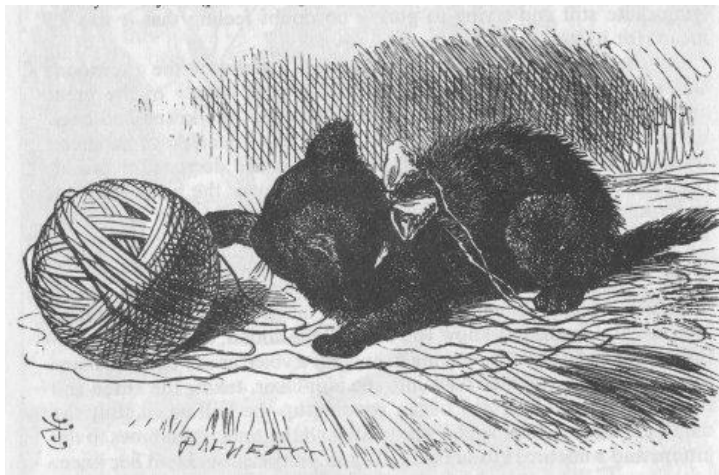


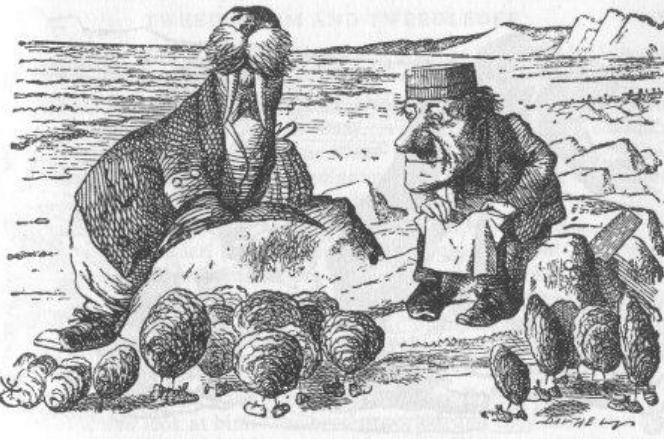
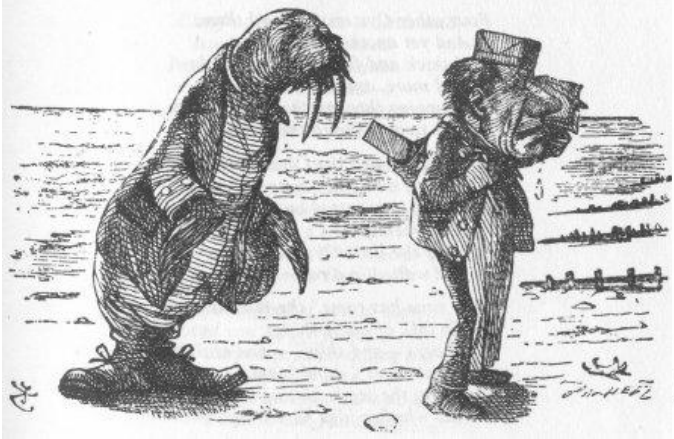
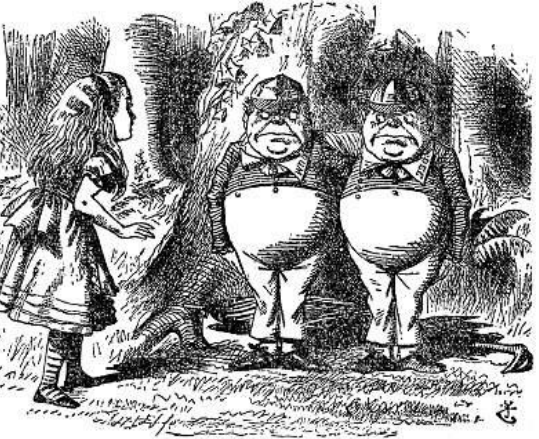
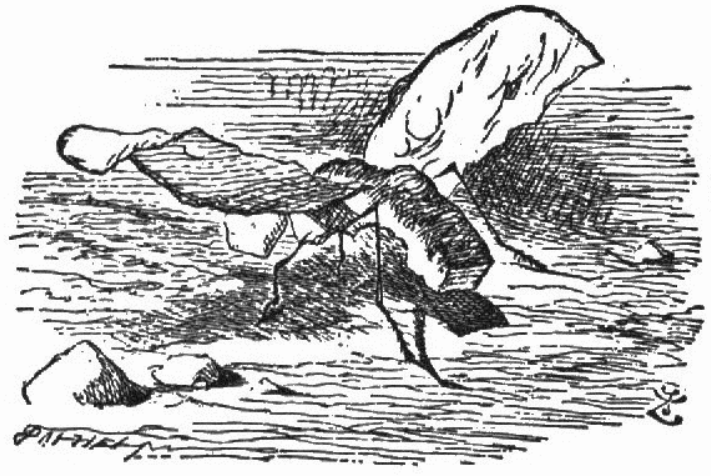
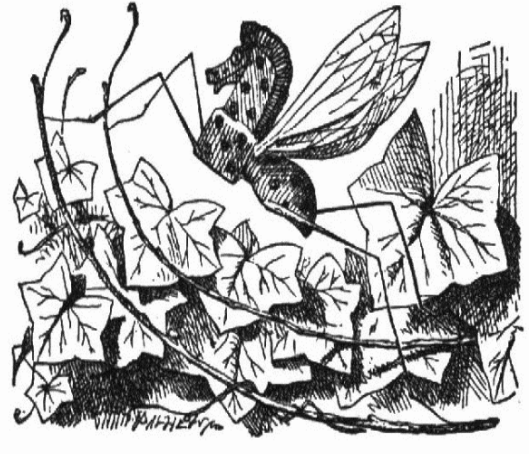
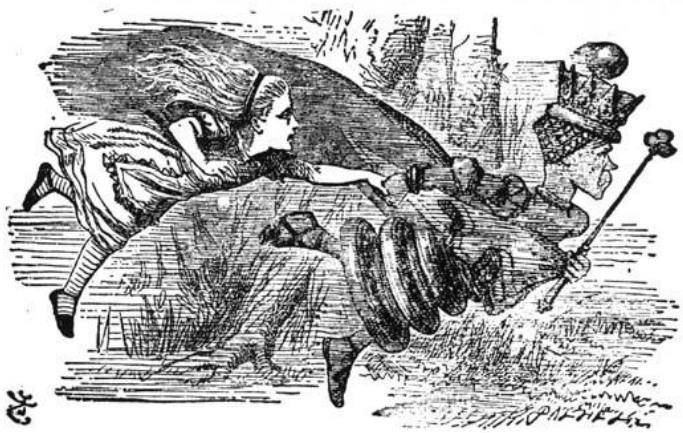


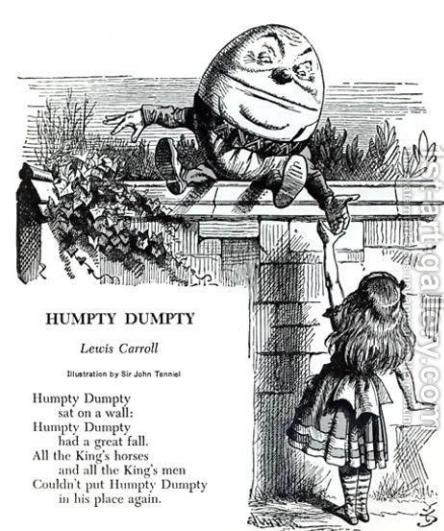
WHITE.

White Pawn (Alice) to play, and win in eleven moves.

	PAGE
1. Alice meets R. Q.	142
2. Alice through Q's 3rd (by railway).	147
3. W. Q. to Q. B's 5th (becomes sheep) and Tweedledes.	149
4. Alice meets W. Q. (with shawl).	158
5. Alice to Q's 5th (along river, sheep).	173
6. Alice to Q's 6th (Humpty Dumpty).	179
7. W. K. takes R. K1.	180
8. Alice to Q's 7th (jumps).	202
9. Alice becomes Queen.	213
10. Alice castles (feast).	223
11. Alice takes R.Q. & wins.	230







HUMPTY DUMPTY

Lewis Carroll

Illustration by Sir John Tenniel

Humpty Dumpty
sat on a wall;
Humpty Dumpty
had a great fall.
All the King's horses
and all the King's men
Couldn't put Humpty Dumpty
in his place again.







VOGUE
DIRETTORE DI VITA

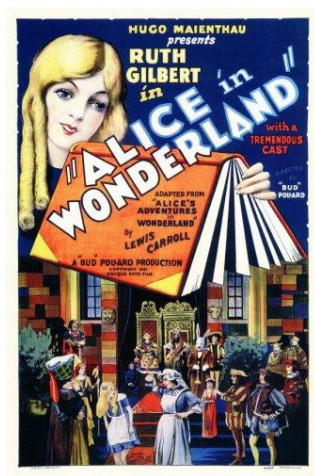
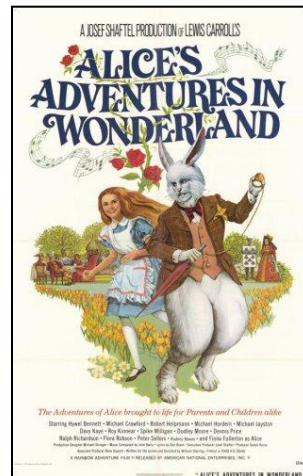
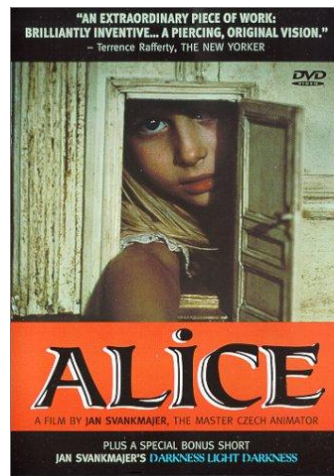
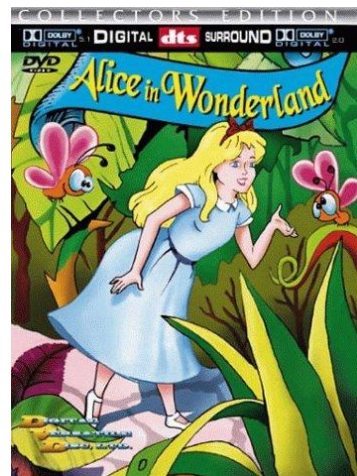
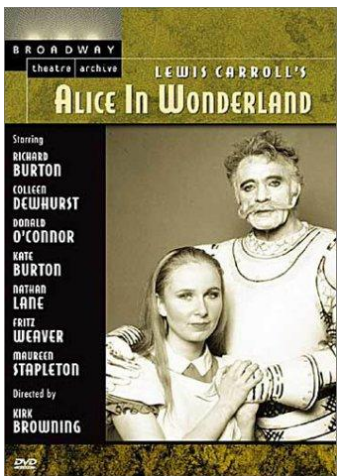
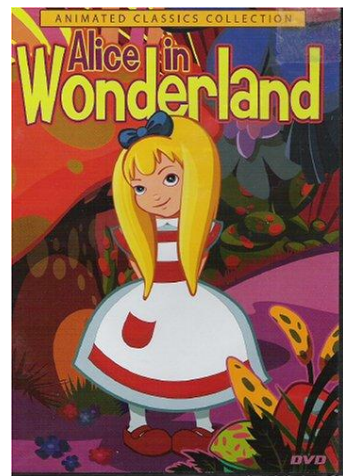
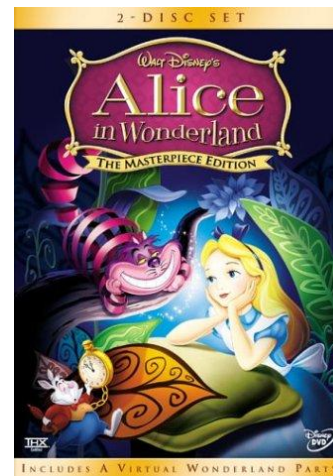
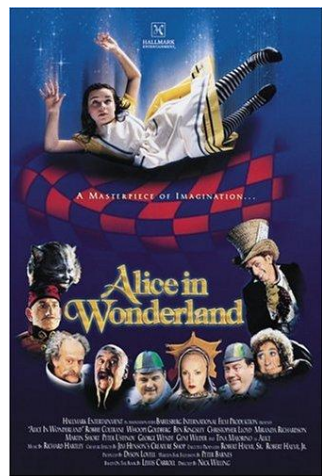
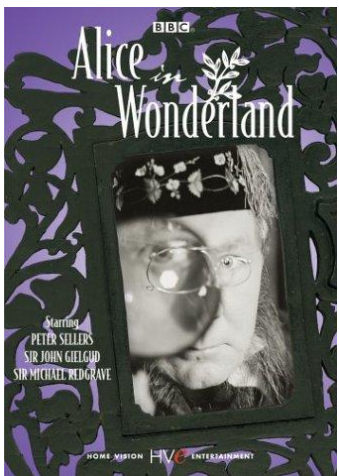
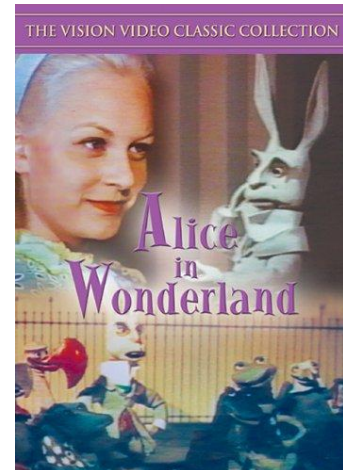
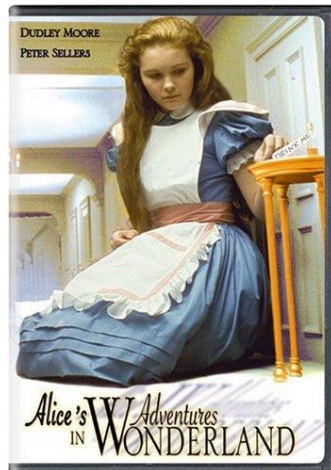
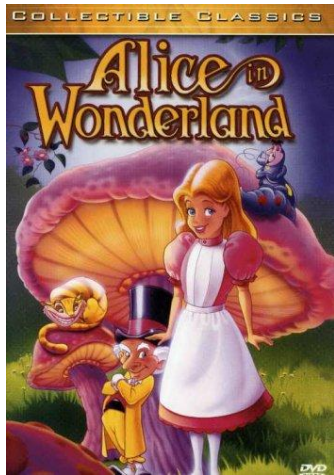
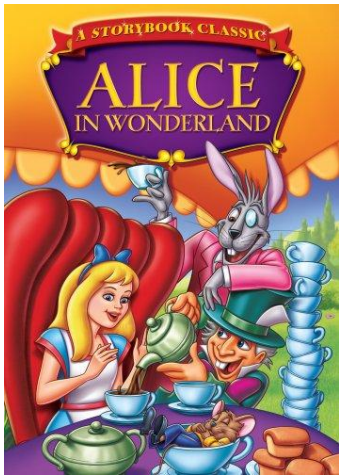


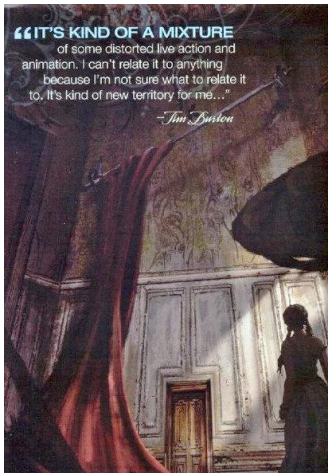
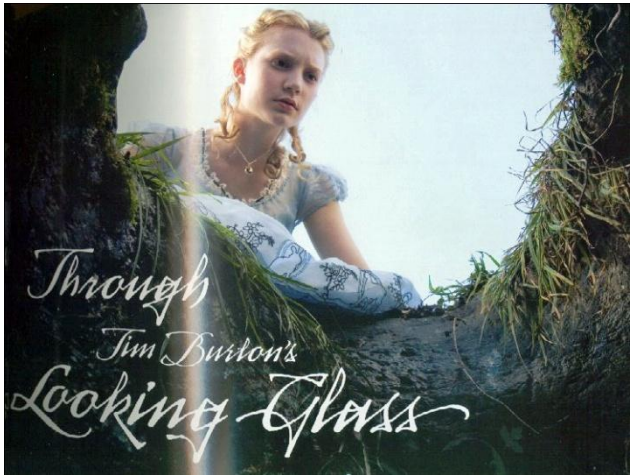
*alice in
wonderland*

*Levi's Carroll dressed her as an innocent in satin and ribbons.
Disney made her flaxen haired and saucer eyed.
In the pages of Vogue the land of merry sublimities and late-raining
rabbits shines to the again - on the world's most influential
designers dress the original little girl lost in their own visions.*

Photographed by Annie Leibowitz







Alice and the White Knight: A Lesson in Logic, Semantics, and Pointers



`You are sad,' the Knight said in an anxious tone: `let me sing you a song to comfort you.'

`Is it very long?' Alice asked, for she had heard a good deal of poetry that day.

`It's long,' said the Knight, `but it's very, *very* beautiful. Everybody that hears me sing it -- either it brings the *tears* into their eyes, or else --'

`Or else what?' said Alice, for the Knight had made a sudden pause.
logical disjunction!
law of the excluded middle!

`Or else it doesn't, you know. **The name of the song is called "Haddocks' Eyes"**.'

`Oh, that's the name of the song, is it?' Alice said, trying to feel interested.
pointer to a pointer!

`No, you don't understand,' the Knight said, looking a little vexed. `That's what the name is *called*. **The name really is "The Aged Aged Man"**.'
pointer dereferencing: meta-pointer resolved!

`Then I ought to have said "That's what the *song* is called"?' Alice corrected herself.
separation of abstractions: variable vs. pointer!

`No, you oughtn't: that's quite another thing! **The song is called "Ways and Means"**: but that's only what it's *called*, you know!'

`Well, what *is* the song, then?' said Alice, who was by this time completely bewildered.

`I was coming to that,' the Knight said. `**The song really is "A-sitting On a Gate"**': and the tune's my own invention.'
call-by-name vs. call-by-value!

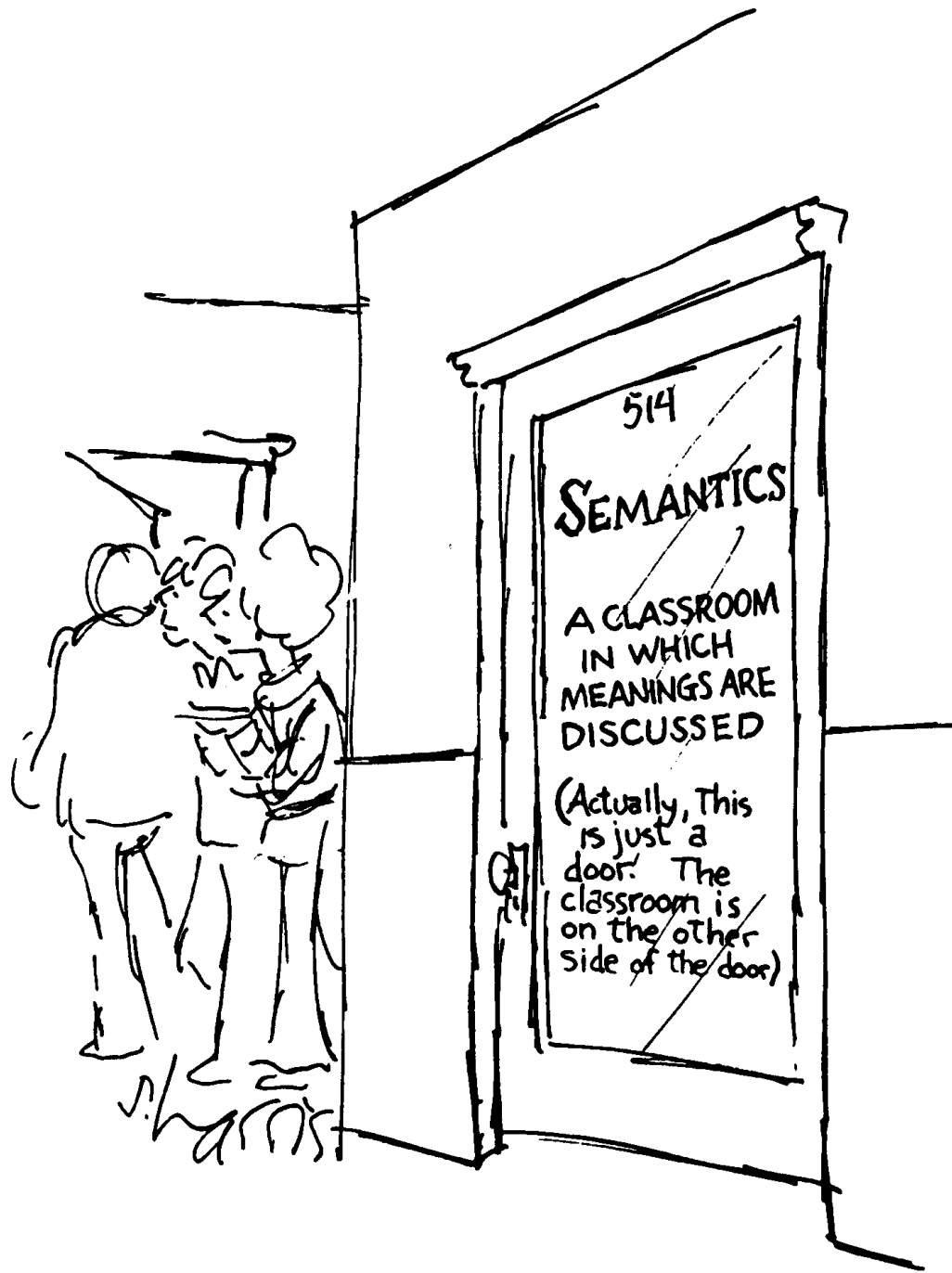
the name is called
"Haddocks' Eyes"

the name of the song is
"The Aged Aged Man"

the song is
"A-sitting On a Gate"

the song is called
"Ways and Means"





514

SEMANTICS

A CLASSROOM
IN WHICH
MEANINGS ARE
DISCUSSED

(Actually, This
is just a
door. The
classroom is
on the other
side of the door)

S. HARTS

Lewis Carroll Society of North America

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WELCOME

Welcome to The Lewis Carroll Society of North America (LCSNA) homepage. The LCSNA is a non-profit organization dedicated to furthering Carroll studies, increasing accessibility of research material, and maintaining public awareness of Carroll's contributions to society and culture. This website is one way we share information with Carroll enthusiasts around the World. If you are a Carrollian and would like to help in these endeavors, or if you simply enjoy Carroll and want to be among other people with a like interest, please consider [joining](#) the LCSNA.

For detailed information about C.L.Dodgson ("Lewis Carroll") and his creations, please access the [Lewis Carroll Homepage](#).

Spring Meeting

The 2009 Spring meeting will be held in beautiful Sante Fe, New Mexico, on May 9. Please consult the [newly updated \(as of April 24th\) meeting agenda](#) for all of the details. See you there.

Alice



An Educational Software that teaches students computer programming in a 3D environment

FREE !!

- About Alice
- Downloads
- Teaching
- Community
- Publications
- Support



 **The Alice Project announces a unique collaboration with Sun Microsystems**  [read more...](#)

Alice 3 News



Alice 3 wins Duke's Choice Award at JavaOne 2009!

[Read more...](#)

All about Alice

Alice is an innovative 3D programming environment that makes it easy to create an animation for telling a story, playing an interactive game, or a video to share on the web. Alice is a teaching tool for introductory computing. It uses 3D graphics and a drag-and-drop interface to facilitate a more engaging, less frustrating first programming experience.

[Read more...](#)

Teaching Materials

Alice is a teaching tool designed as a revolutionary approach to teaching and learning introductory programming concepts. The Alice team has developed instructional materials to support students and teachers in using this new approach. Resources include textbooks, lessons, sample syllabuses, test banks, and more. Other authors have generously joined our efforts, creating additional textbooks.

[Read more...](#)

Downloads

[Alice 2.2](#) [Alice 2.0](#)
Designed for High School and College

[Storytelling Alice](#)
Designed for Middle School

[Alice 3 beta](#)
Get a sneak peek at the future of Alice

[3D Models Gallery](#)
Additional free 3D models

Alice Blog

Check out the Alice blog! The Alice team discusses the latest in Alice development. View screencasts demonstrating new features, tips and techniques!

[Visit blog...](#)

Community Forums

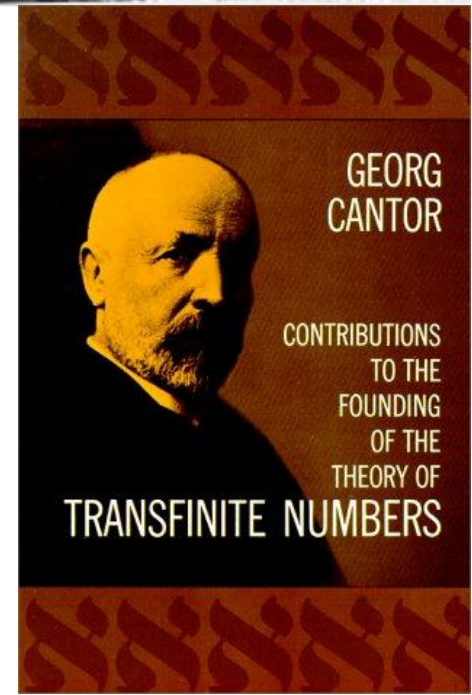
Share and gather knowledge about Alice through our community forums. Students, teachers and enthusiasts are all welcome! If you have a question or comment about Alice, post it here!

[View forums...](#)

Historical Perspectives

Georg Cantor (1845-1918)

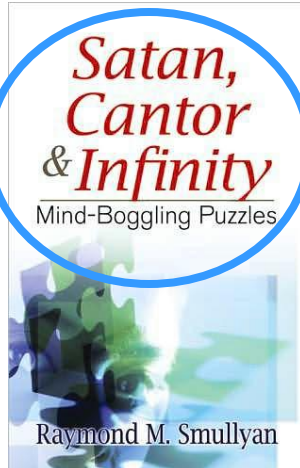
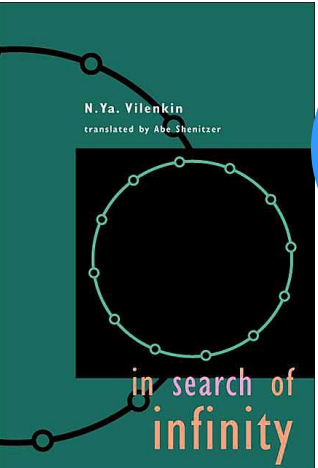
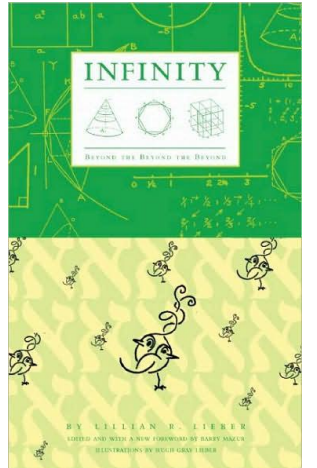
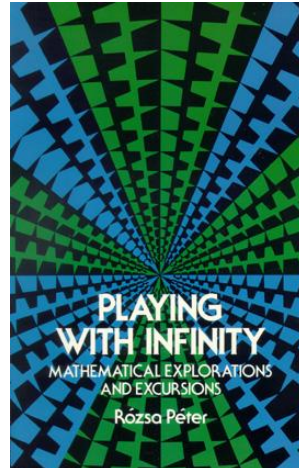
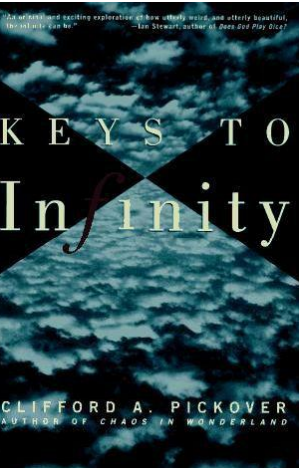
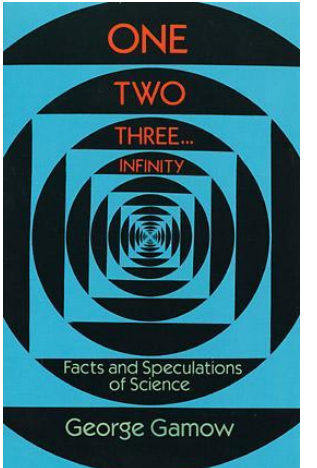
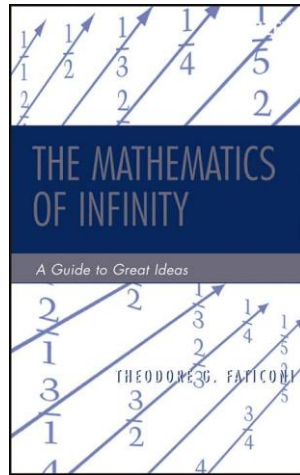
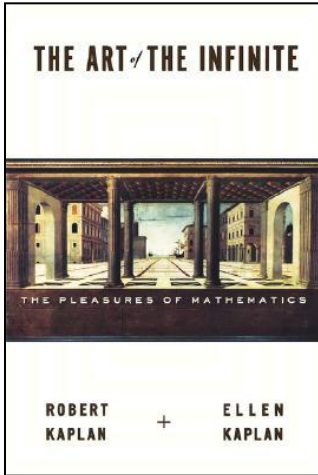
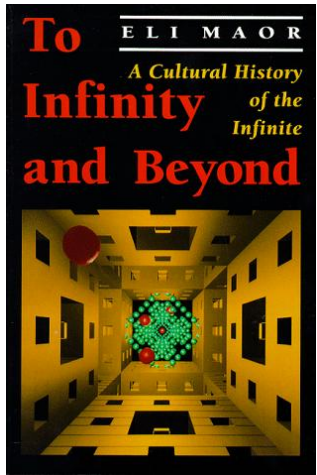
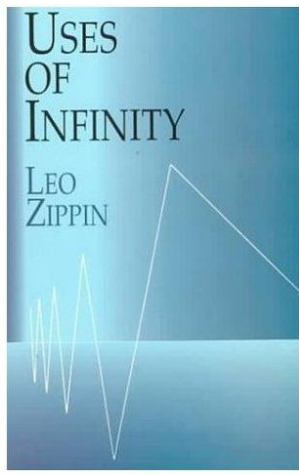
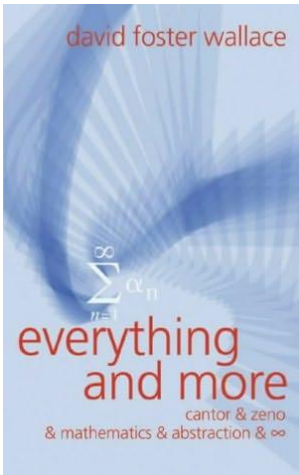
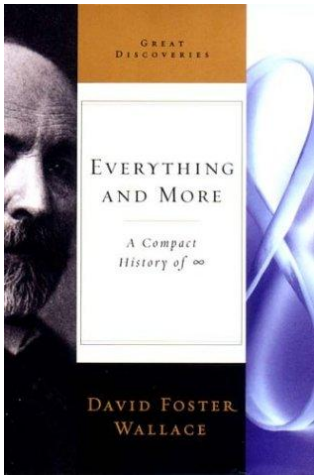
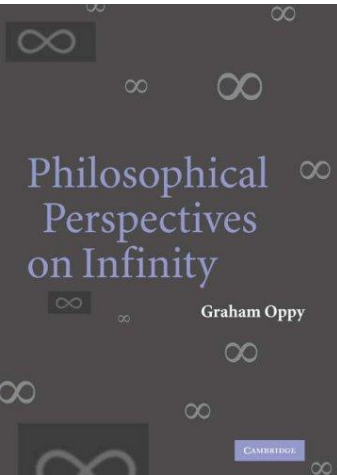
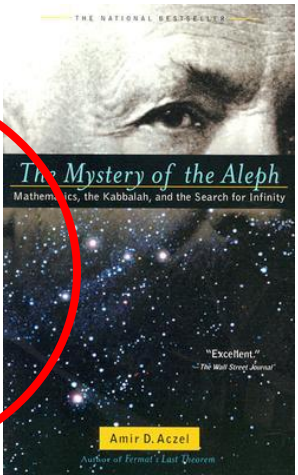
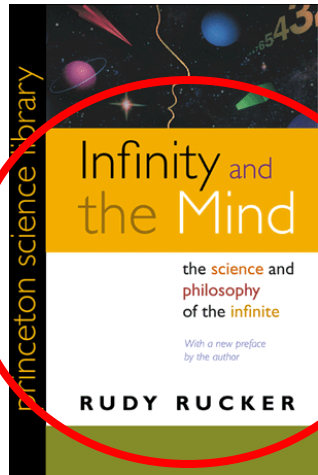
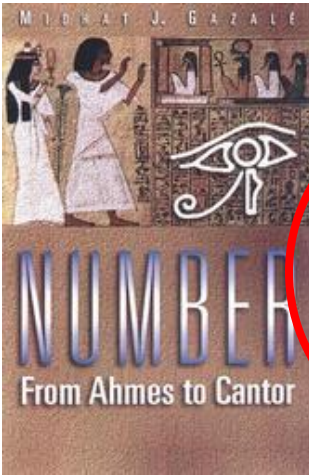
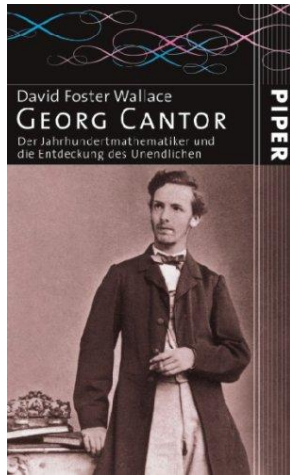
- Created modern set theory
- Invented trans-finite arithmetic (highly controversial at the time)
- Invented diagonalization argument
- First to use 1-to-1 correspondences with sets
- Proved some infinities “bigger” than others
- Showed an infinite hierarchy of infinities
- Formulated continuum hypothesis
- Cantor’s theorem, “Cantor set”, Cantor dust, Cantor cube, Cantor space, Cantor’s paradox
- Laid foundation for computer science theory
- Influenced Hilbert, Godel, Church, Turing

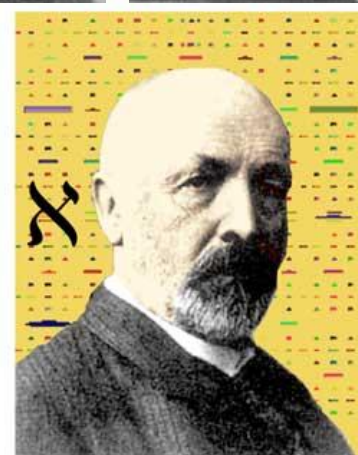
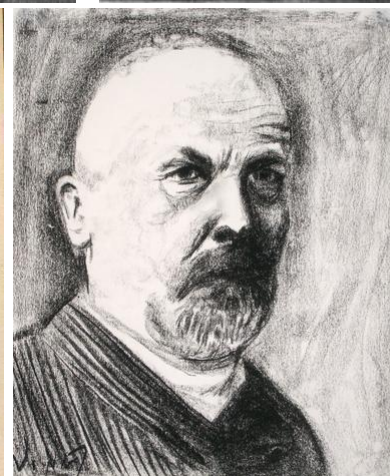


GEORG CANTOR
His Mathematics and
Philosophy of the Infinite

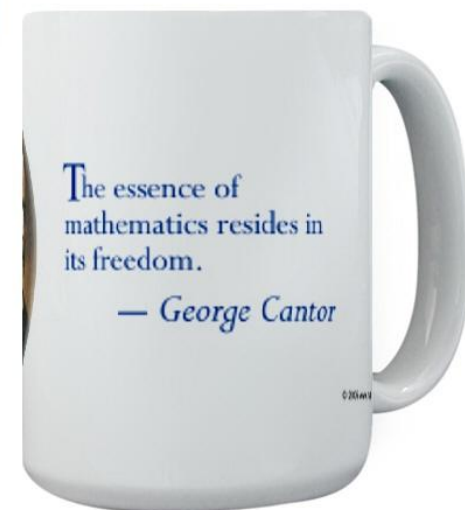


Joseph Warren Dauben



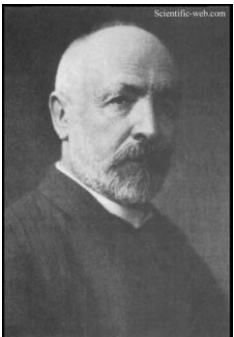


Georg Cantor
1845 - 1918



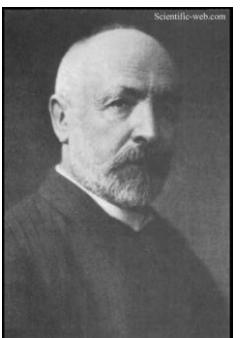
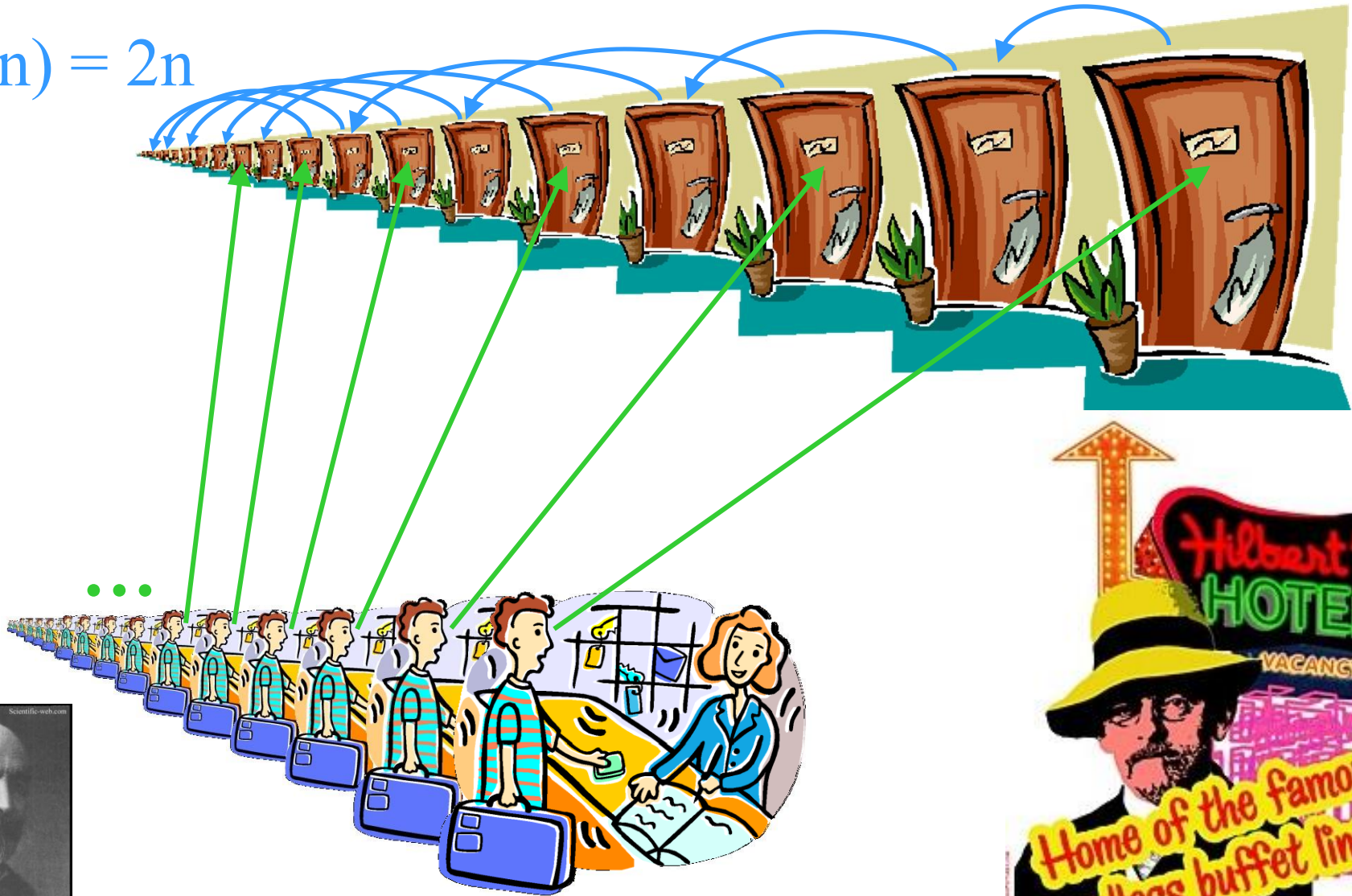
Problem: How can a **new** guest be accommodated in a **full** infinite hotel?

$$f(n) = n+1$$

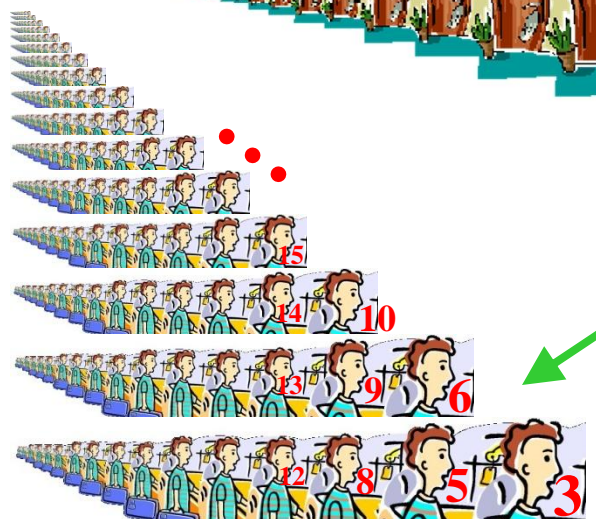


Problem: How can an **infinity** of **new** guests be accommodated in a **full** infinite hotel?

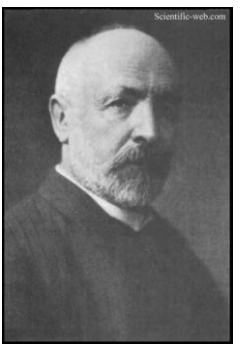
$$f(n) = 2n$$

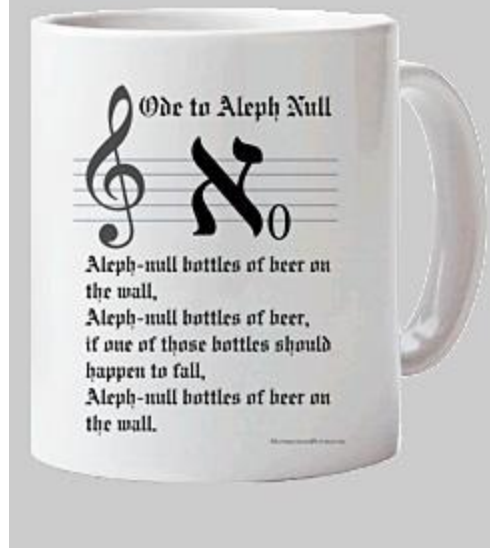
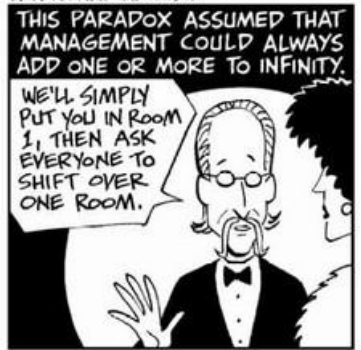


Problem: How can an infinity of infinities of new guests be accommodated in a full infinite hotel?



one-to-one
correspondence





- 8th June 2009
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HOTEL INFINITY

Amanda Boyle

average rating from 27 members ● ● ● ● ● ○

drama | 2004 | London | Switzerland/ Mandarin | 10 min

Published 28 Feb 07

What happens when an hotel of infinite rooms suddenly becomes full?

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 **PLAY NOW**
Requires [windows media player](#) or [real player](#).

synopsis

There was once a hotel in the mountains that was so popular with guests, the Manager decided to extend it. Yet still it remained full. The Manager continued to extend it, until eventually it became infinitely large. One day much to his surprise, no spare rooms could be found in his now infinite hotel. All the mathematical calculations, which he normally relied on to find the rooms, just wouldn't work...

rate this film

5 ● ● ● ● ● ○
4 ● ● ● ● ○ ○
3 ● ● ● ○ ○ ○
2 ● ● ○ ○ ○ ○
1 ● ○ ○ ○ ○ ○

Problem: Are there more rationals than natural #'s?

$\mathbb{N} \subset \mathbb{Q}$

$\mathbb{N} \neq \mathbb{Q}$

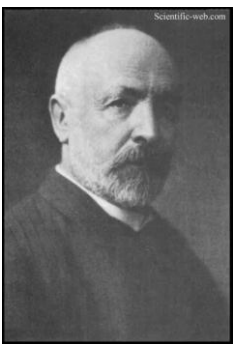
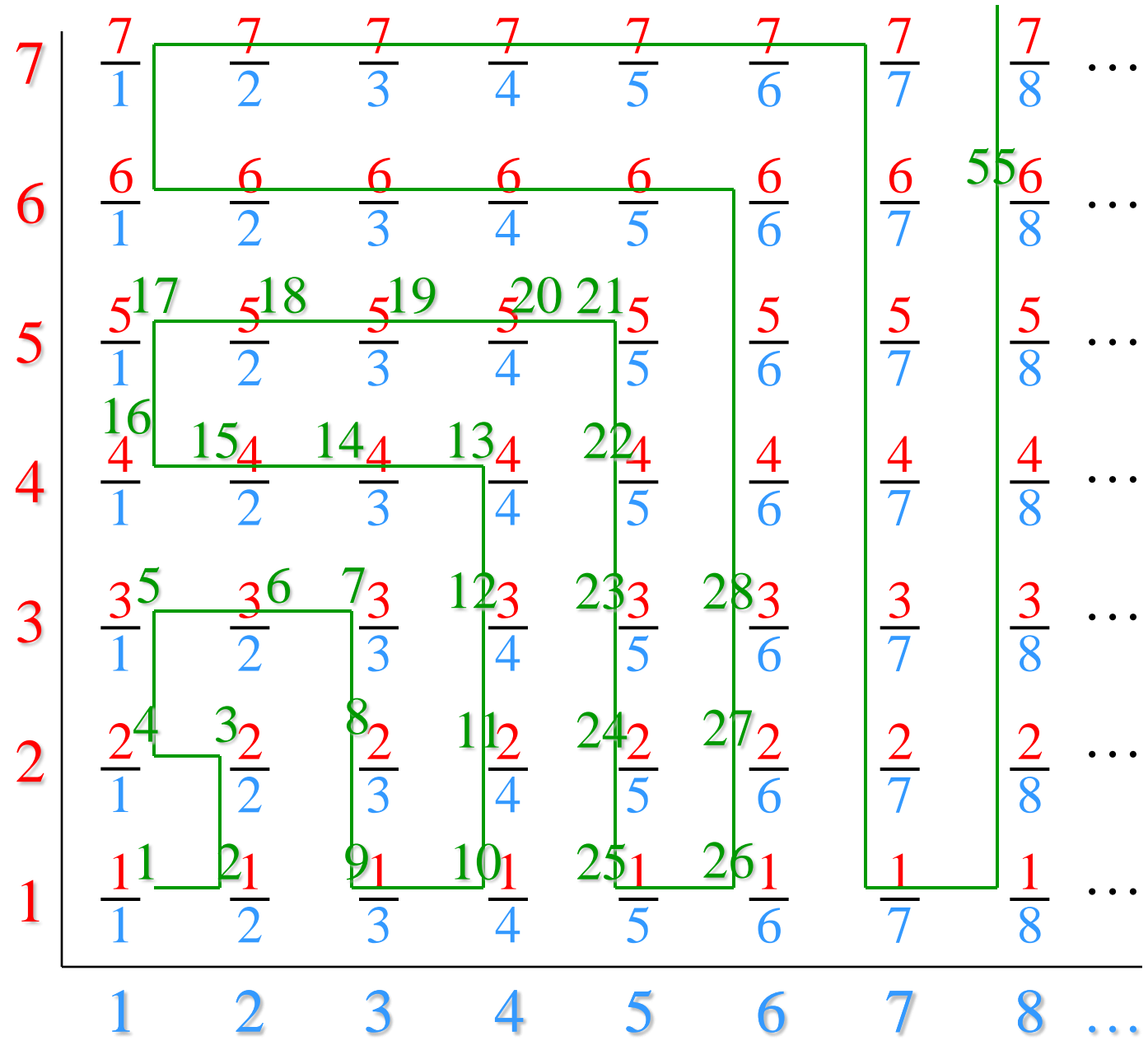
So $|\mathbb{N}| < |\mathbb{Q}|$?

Dovetailing:

Establishes 1-1 correspondence

$f: \mathbb{N} \leftrightarrow \mathbb{Q}$

$\Rightarrow |\mathbb{N}| = |\mathbb{Q}|$



Problem: Are there more rationals than natural #'s?

$\mathbb{N} \subset \mathbb{Q}$

$\mathbb{N} \neq \mathbb{Q}$

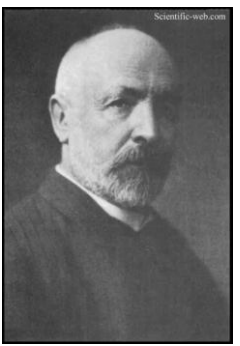
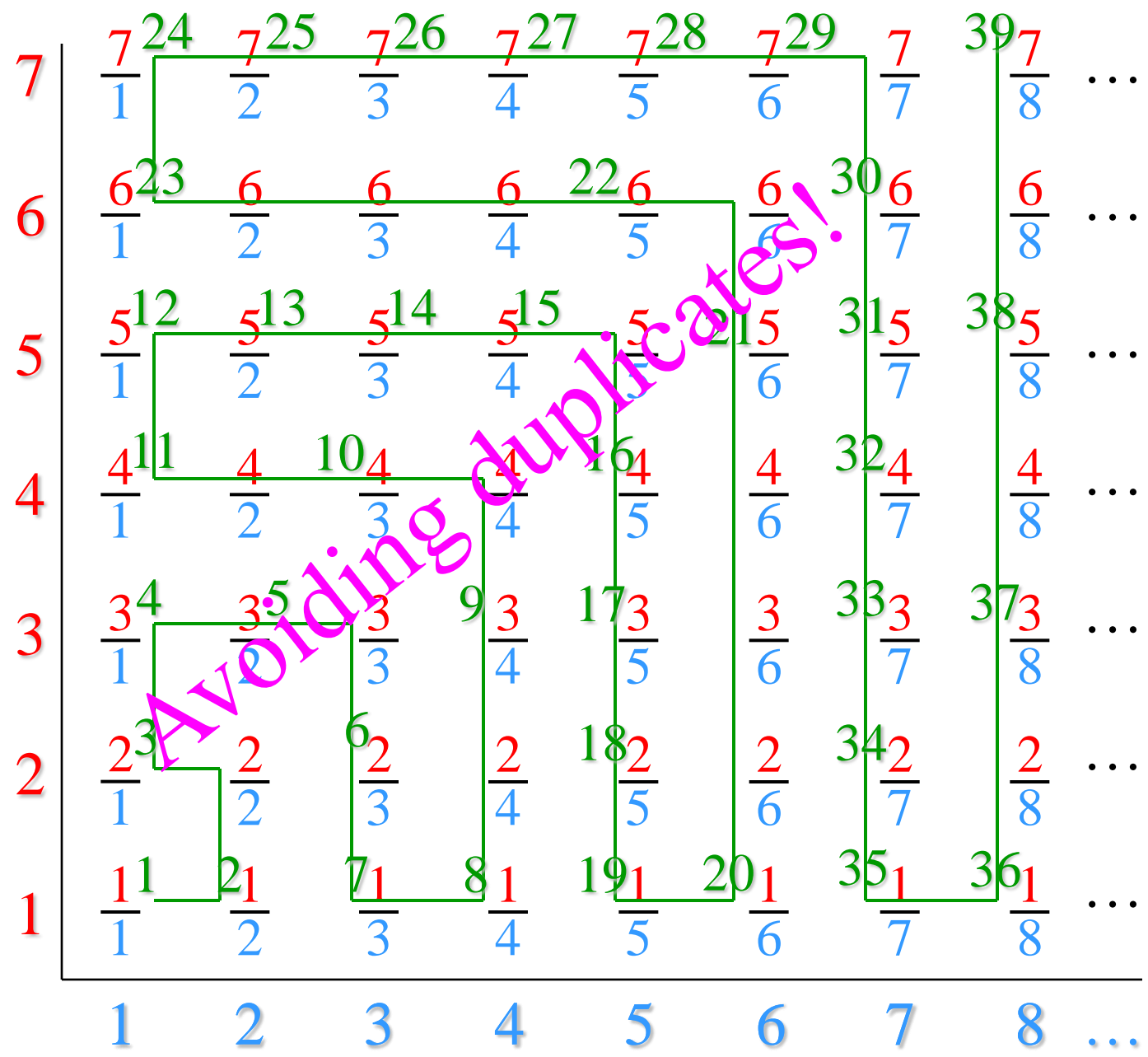
So $|\mathbb{N}| < |\mathbb{Q}|$?

Dovetailing:

Establishes 1-1 correspondence

$f: \mathbb{N} \leftrightarrow \mathbb{Q}$

$\Rightarrow |\mathbb{N}| = |\mathbb{Q}|$



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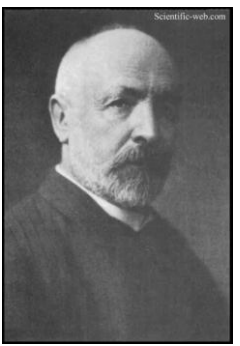
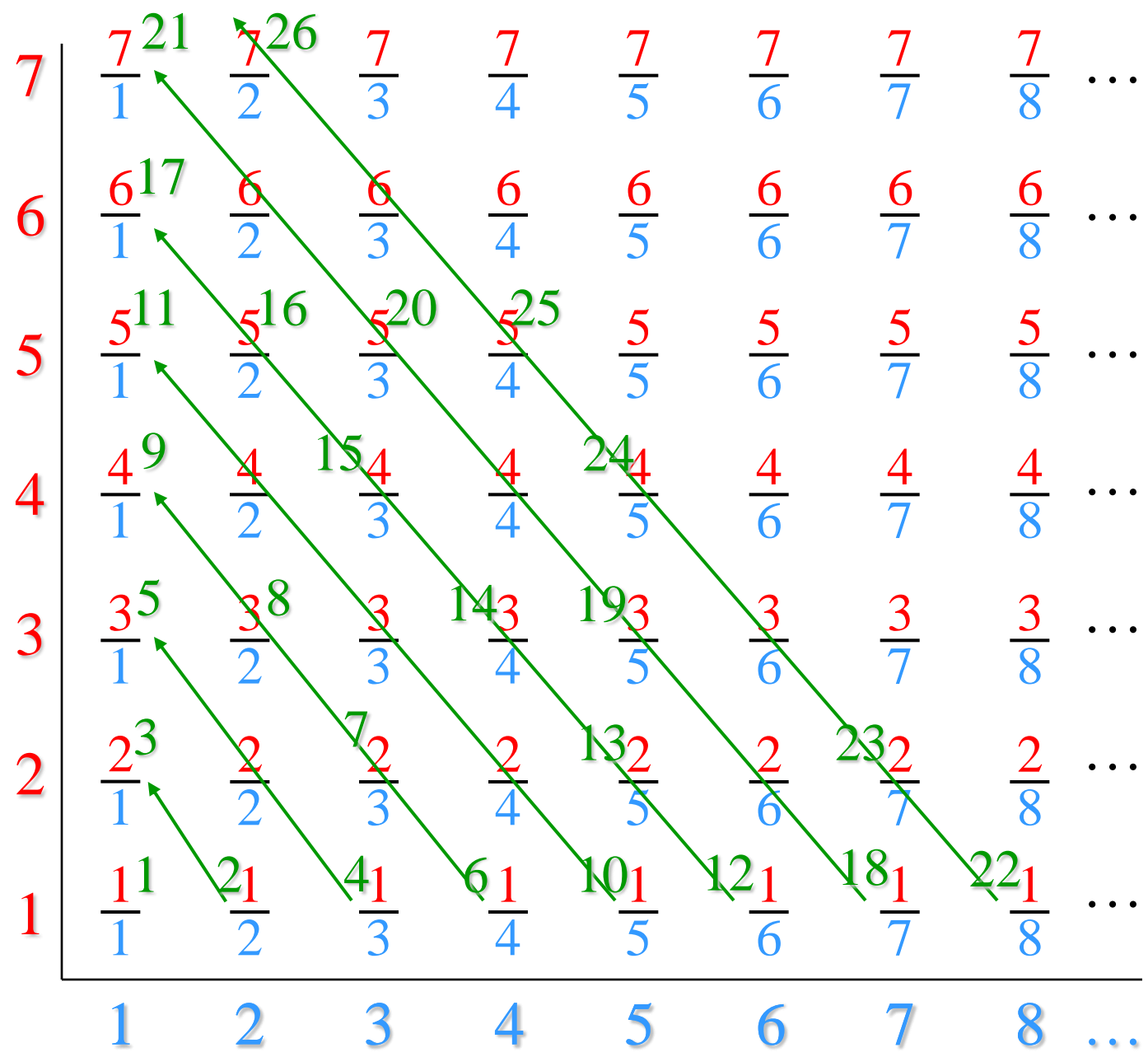
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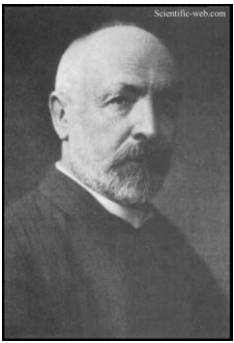
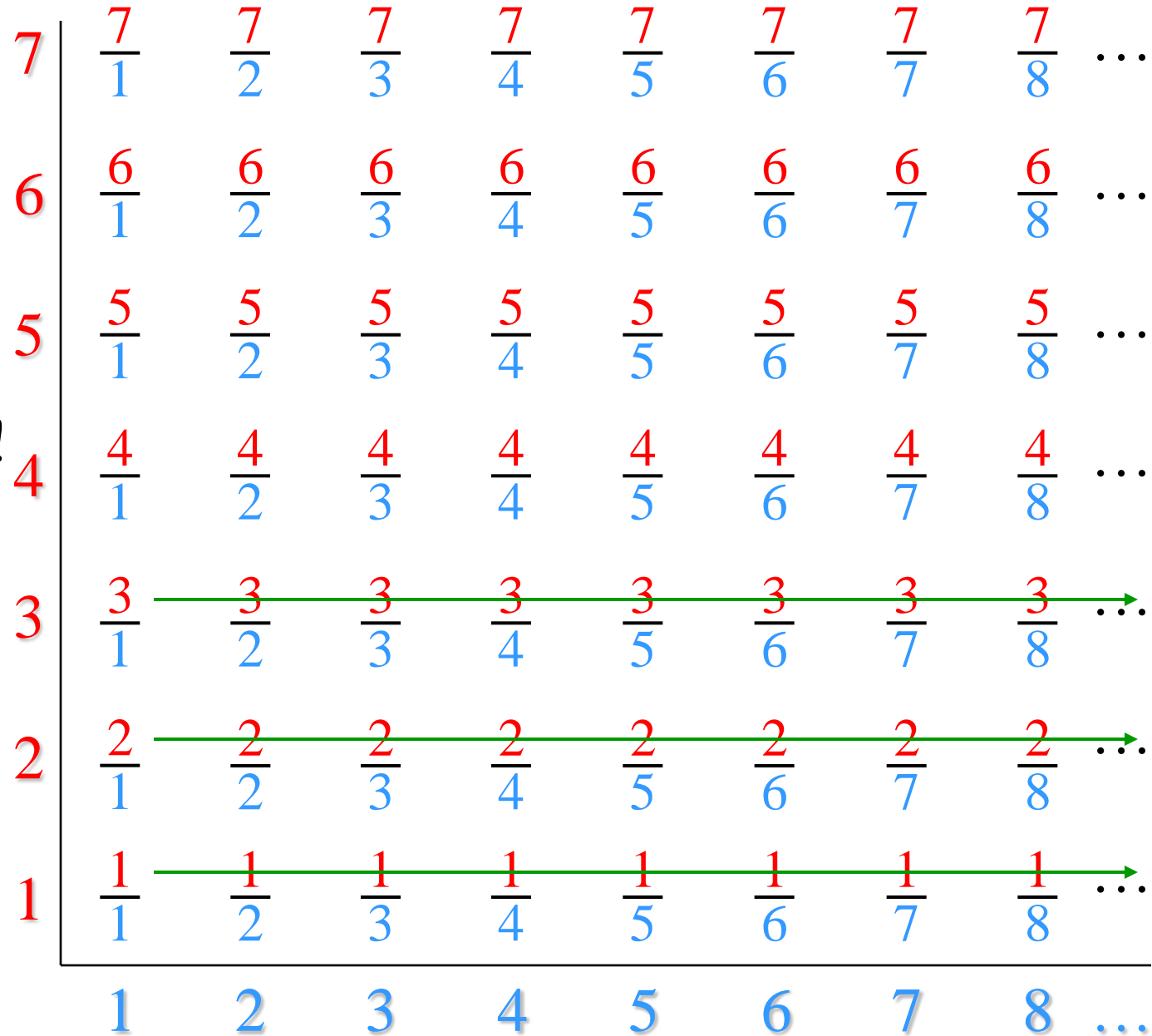


Problem: Why doesn't this "dovetailing" work?

There's no
"last" element
on the first line!

So the 2nd line
is never reached!

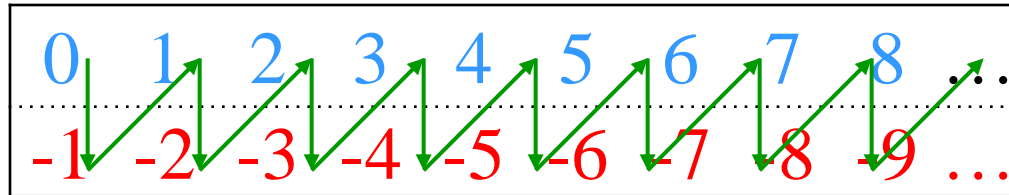
⇒ 1-1 function
is not defined!



Dovetailing Reloaded

Dovetailing: $f: \mathbb{N} \leftrightarrow \mathbb{Z}$

-4	-3	-2	-1	0	1	2	3	4
----	----	----	----	---	---	---	---	---



\mathbb{Z}

\mathbb{N} 1 2 3 4 5 6 7 8 9

To show $|\mathbb{N}| = |\mathbb{Q}|$ we can construct $f: \mathbb{N} \leftrightarrow \mathbb{Q}$ by sorting x/y by increasing key $\max(|x|, |y|)$, while avoiding duplicates:

$$\max(|x|, |y|) = 0 : \{ \}$$

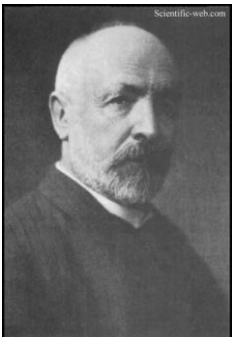
$$\max(|x|, |y|) = 1 : 0/1, 1/1$$

$$\max(|x|, |y|) = 2 : 1/2, 2/1$$

$$\max(|x|, |y|) = 3 : 1/3, 2/3, 3/1, 3/2$$

... { finite new set at each step }

- Dovetailing can have many disguises!
- So can diagonalization!



Theorem: There are more reals than rationals / integers.

Proof [Cantor]: Assume a 1-1 correspondence $f: \mathbb{N} \leftrightarrow \mathbb{R}$ i.e., there exists a table containing all of \mathbb{N} and **all** of \mathbb{R} :

\mathbb{N}	\mathbb{R}
$f(1) =$	3 . 1 4 1 5 9 2 6 5 3 ...
$f(2) =$	1 . 0 0 0 0 0 0 0 0 0 ...
$f(3) =$	2 . 7 1 8 2 8 1 8 2 8 ...
$f(4) =$	1 . 4 1 4 2 1 3 5 6 2 ...
$f(5) =$	0 . 3 3 3 3 3 3 3 3 3 ...
...	...

$X = 0.21934\dots \in \mathbb{R}$

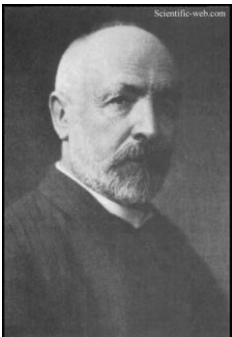
But X is missing from our table! $X \neq f(k) \forall k \in \mathbb{N}$

$\Rightarrow f$ not a 1-1 correspondence

\Rightarrow contradiction

$\Rightarrow \mathbb{R}$ is not countable!

There are more reals than rationals / integers!



Diagonalization Non-existence proof!

Problem 1: Why not just insert X into the table?

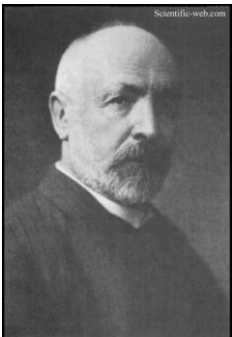
Problem 2: What if $X=0.999\dots$ but $1.000\dots$ is already in table?

\mathbb{N}	\mathbb{R}
$f(1) =$	3 . 1 4 1 5 9 2 6 5 3 ...
$f(2) =$	1 . 0 0 0 0 0 0 0 0 ...
$f(3) =$	2 . 7 1 8 2 8 1 8 2 8 ...
$f(4) =$	1 . 4 1 4 2 1 3 5 6 2 ...
$f(5) =$	0 . 3 3 3 3 3 3 3 3 ...
...	...

$X = 0.21934\dots \in \mathbb{R}$

Diagonalization Non-existence proof!

- Table with X inserted will have X' **still missing!**
Inserting X (or any number of X's) will not help!
- To enforce **unique table values**, we can avoid using 9's and 0's in X.





Celebrity Cruises X a true departure

**WELCOME
TO
INFINITY**

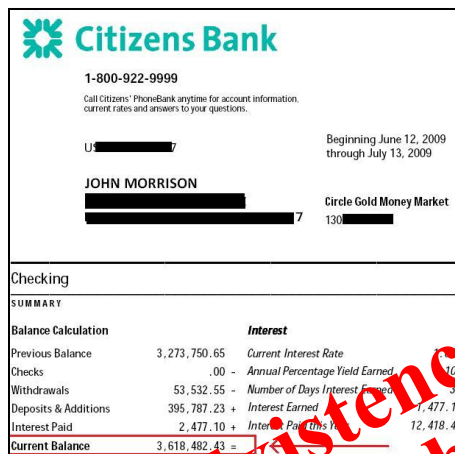
Infinity

Non-Existence Proofs

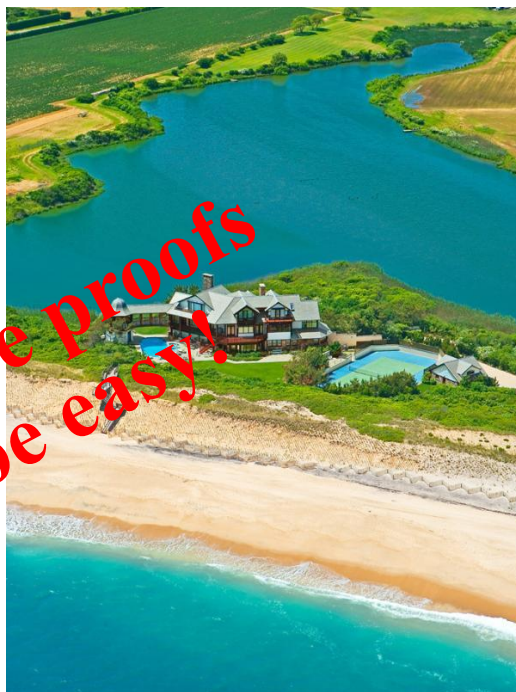
- Must cover **all possible** (usually infinite) scenarios!
- Examples / counter-examples are not convincing!
- **Not “symmetric”** to existence proofs!

Ex: proof that you **are** a millionaire:

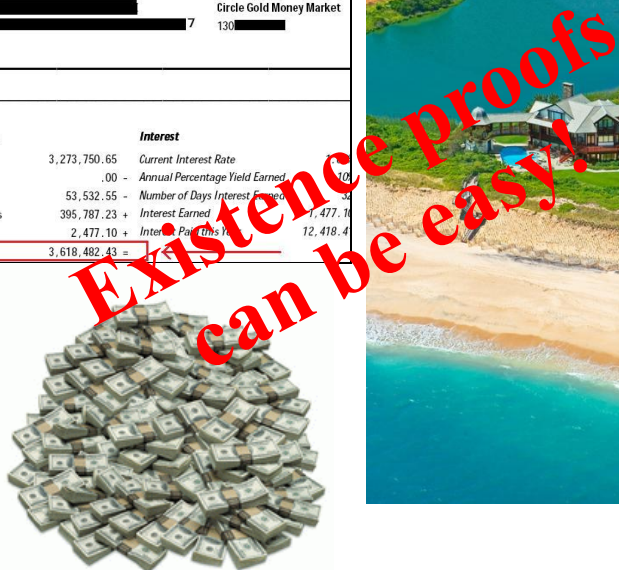
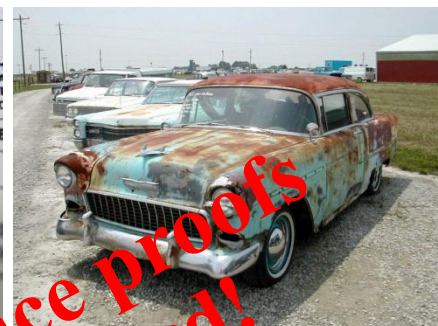
“Proof” that you **are not** a millionaire ?



Balance Calculation		Interest	
Previous Balance	3,273,750.65	Current Interest Rate	
Checks	.00	Annual Percentage Yield Earned	10.34
Withdrawals	53,532.55	Number of Days Interest Earned	30
Deposits & Additions	395,787.23	Interest Earned	1,477.14
Interest Paid	2,477.10	Interest Paid (this year)	12,418.44
Current Balance	3,618,482.43		



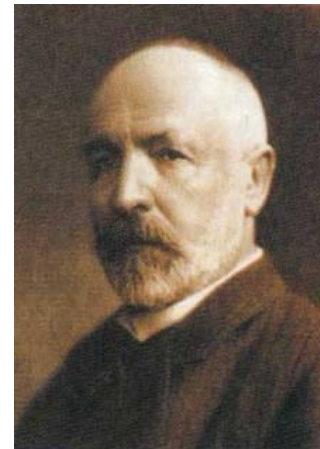
DESCRIPTON	AMOUNT
TOTAL DEPOSITS	2,710.00
YTD INTEREST PAID	



Existence proofs can be easy!

Non-existence proofs are often hard!

$P \neq NP$



Cantor set:

Start with **unit segment**

- Remove (open) **middle third**
- **Repeat recursively** on all remaining segments
- Cantor set is all the **remaining points**



Total **length** removed: $1/3 + 2/9 + 4/27 + 8/81 + \dots = 1$

Cantor set **does not contain any intervals**

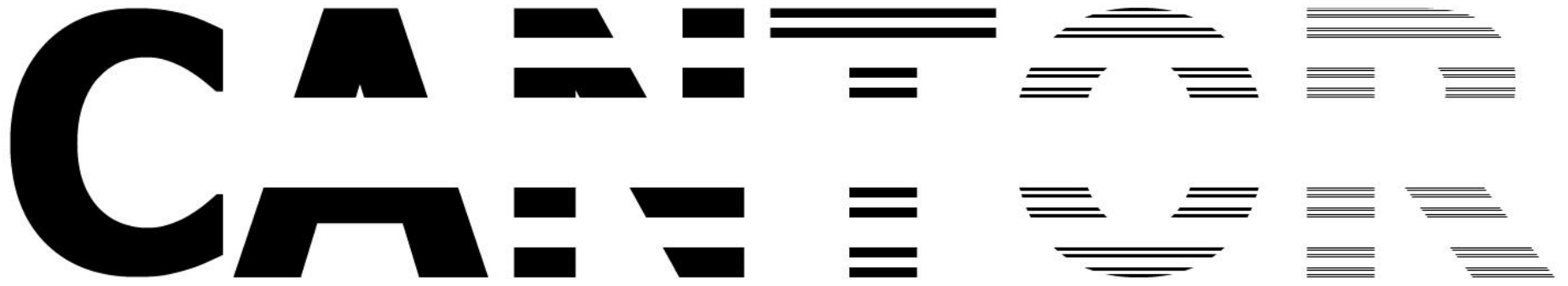
Cantor set is **not empty** (since, e.g. interval endpoints remain)

An **uncountable number of non-endpoints** remain as well (e.g., $1/4$)

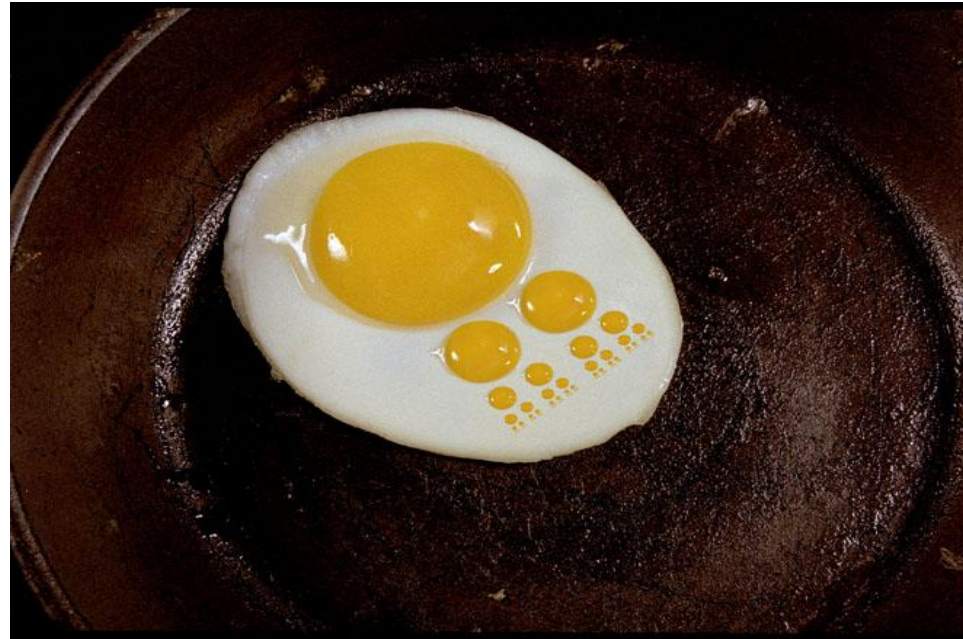
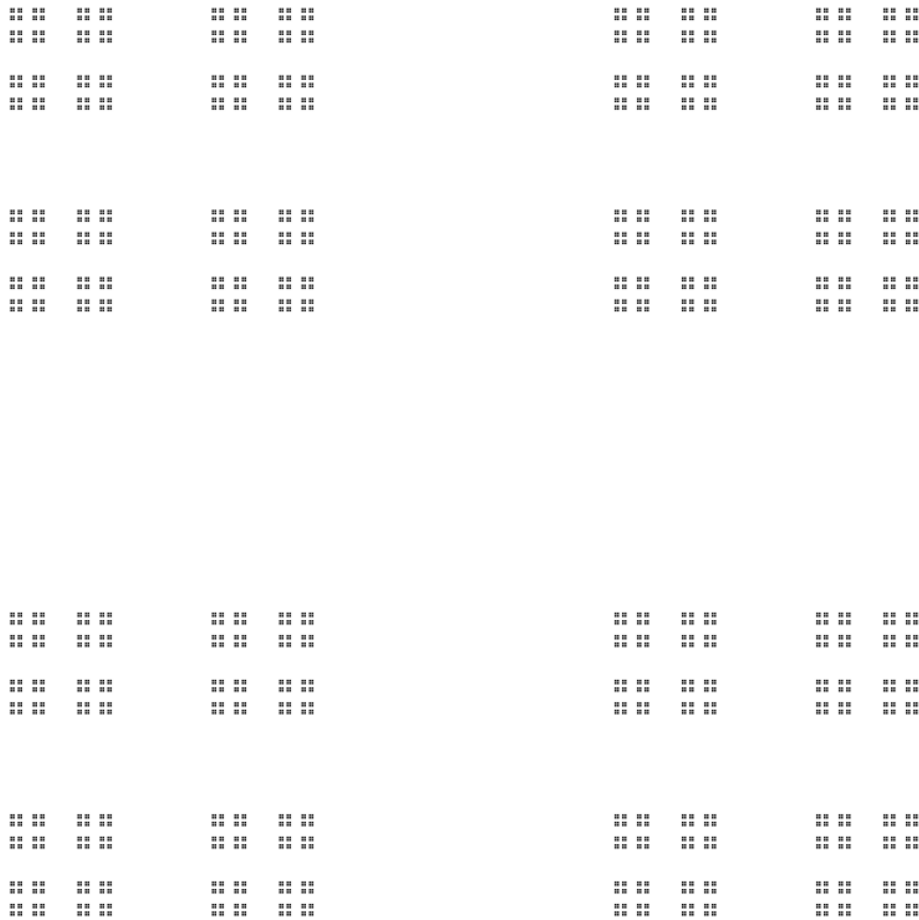
Cantor set is **totally disconnected** (no nontrivial connected subsets)

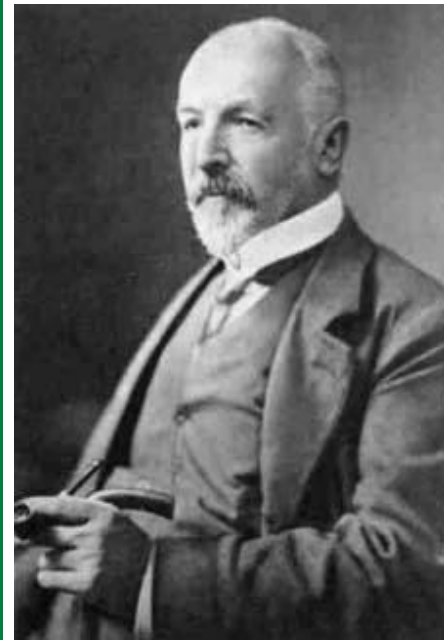
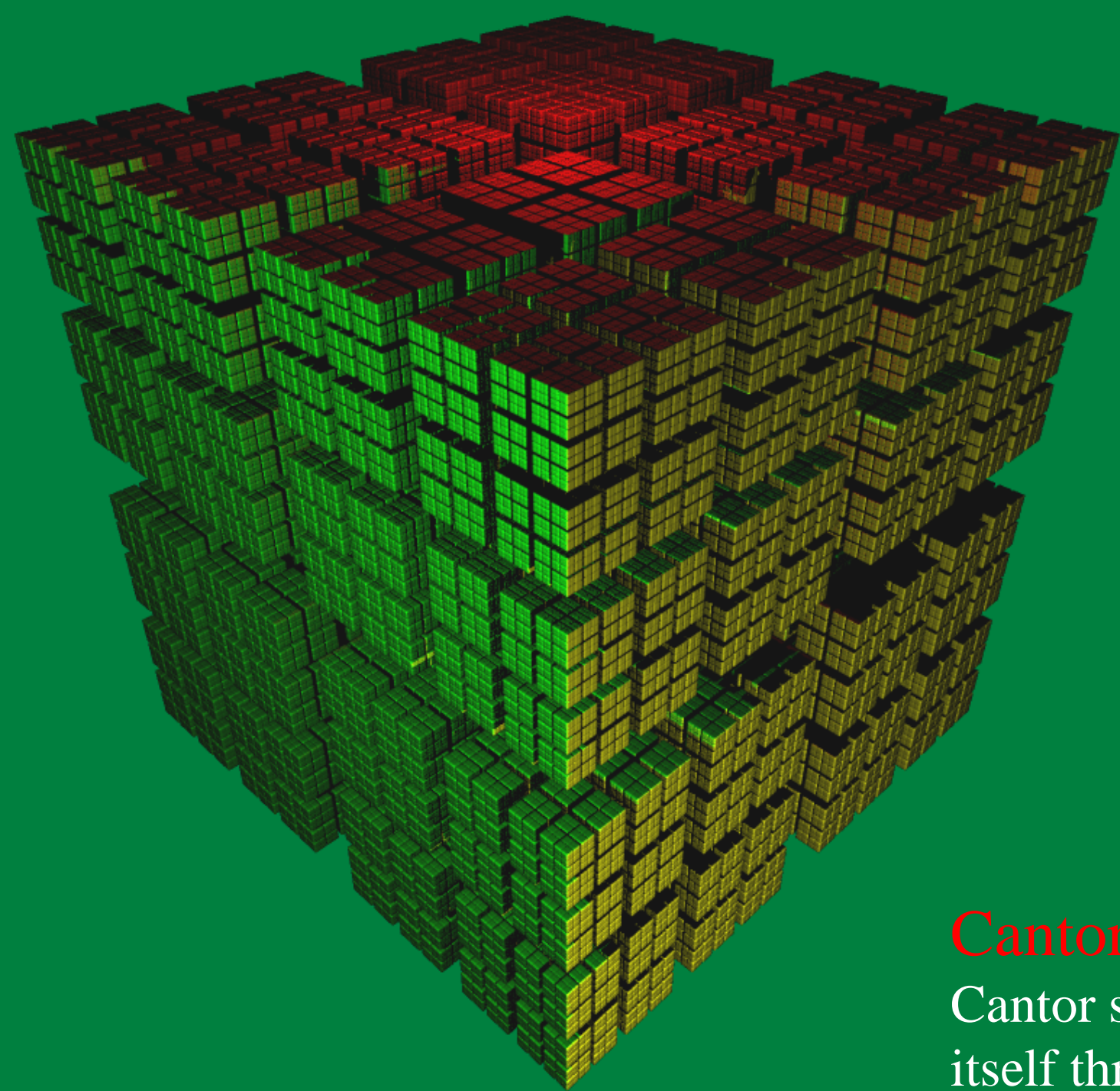
Cantor set is **self-similar** with Hausdorff dimension of $\log_3 2 = 1.585$

Cantor set is a **closed**, totally bounded, **compact**, complete metric space, with **uncountable** cardinality and lebesgue **measure zero**



Cantor dust (2D generalization): Cantor set crossed with itself





Cantor cube (3D):
Cantor set crossed with
itself three times

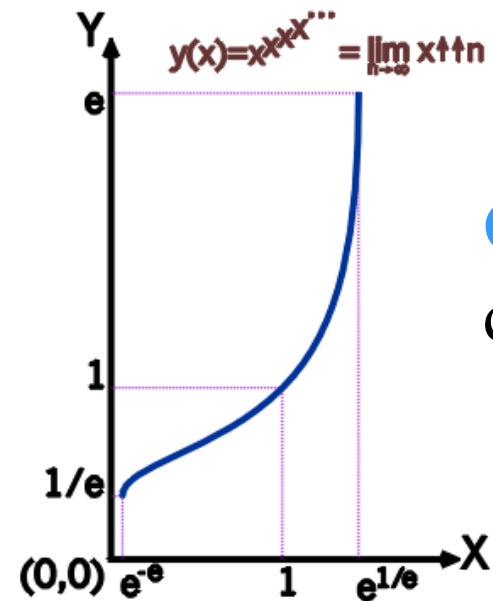
Problem: Solve the following equation for X:

$$X^{X^{X^{X^{\dots}}}} = 2$$

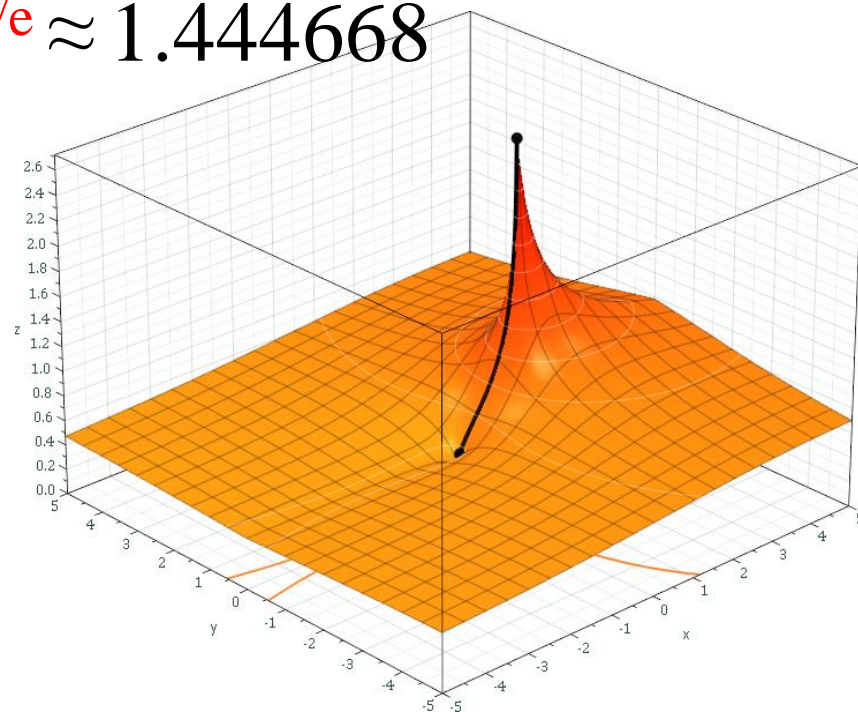
where the stack of exponentiated x's extends forever.

This “power tower” **converges** for:

$$0.065988 \approx e^{-e} < X < e^{1/e} \approx 1.444668$$



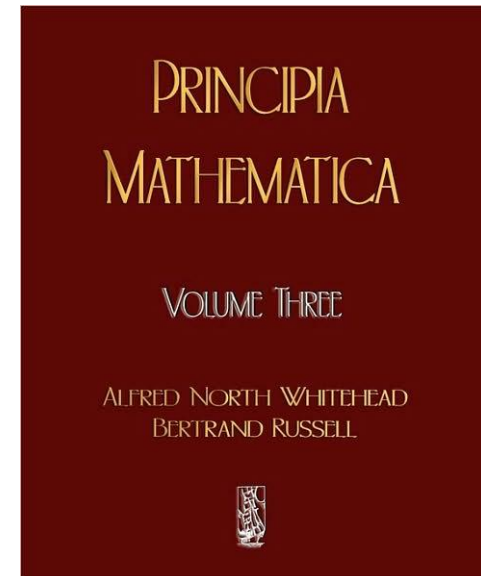
Generalization to
complex numbers:



Historical Perspectives


Bertrand Russell (1872-1970)

- Philosopher, logician, mathematician, historian, social reformist, and pacifist
- Co-authored “[Principia Mathematica](#)” (1910)
- [Axiomatized mathematics](#) and set theory
- Co-founded [analytic philosophy](#)
- Originated [Russell’s Paradox](#)
- [Activist: humanitarianism](#), pacifism, education, free trade, nuclear disarmament, birth control gender & racial equality, gay rights
- Profoundly [transformed math & philosophy](#), mentored Wittgenstein, influenced Godel
- Laid [foundation](#) for computer science theory
- Won Nobel Prize in literature (1950)




The Problems of Philosophy

... Bertrand Russell, a Welsh atheist, wore many hats including philosopher, historical, logician, mathematician, and social reformer. In 1950 he won a Nobel Prize in literature for his humanitarianism and freedom of thought. In this book Russell attempts to give an easily accessible look at problems in philosophy.




Bertrand Russell

Russell




An Inquiry into Meaning and Truth

The Analysis of the Mind



Bertrand Russell


Bertrand Russell



INTRODUCTION TO MATHEMATICAL PHILOSOPHY

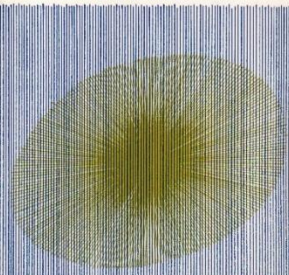
BERTRAND RUSSELL

MY PHILOSOPHICAL DEVELOPMENT



RELIGION AND SCIENCE

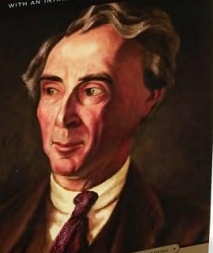
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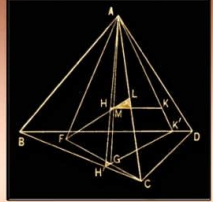
OUR KNOWLEDGE OF THE EXTERNAL WORLD

BERTRAND RUSSELL

WITH AN INTRODUCTION BY AMIT KAGAR




THE FOUNDATIONS OF GEOMETRY



BERTRAND A. W. RUSSELL

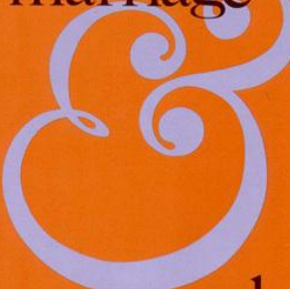
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AUTHORITY AND THE INDIVIDUAL

Bertrand Russell

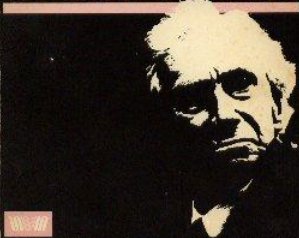
marriage



morals

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THE CONQUEST OF HAPPINESS




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Freedom versus Organization

1814-1914

THE PATTERN OF POLITICAL CHANGES IN 19TH CENTURY EUROPEAN HISTORY.



THE ART OF PHILOSOPHIZING AND OTHER ESSAYS

BERTRAND RUSSELL

The A B C Of Atoms


Bertrand Russell

Bertrand RUSSELL'S

Dictionary of MIND MATTER and MORALS

UNWIN BOOKS


Principles of Social Reconstruction



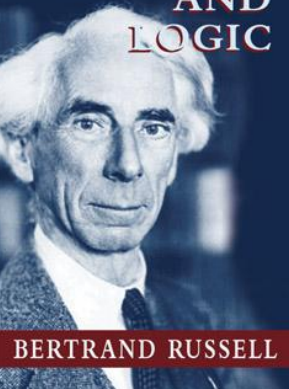
BERTRAND RUSSELL

Bertrand Russell

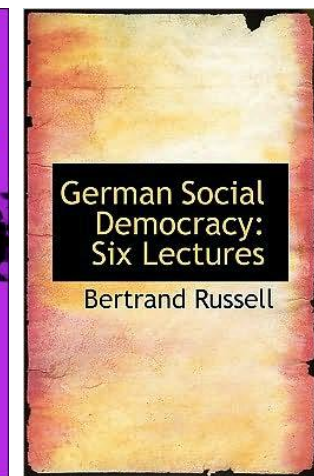
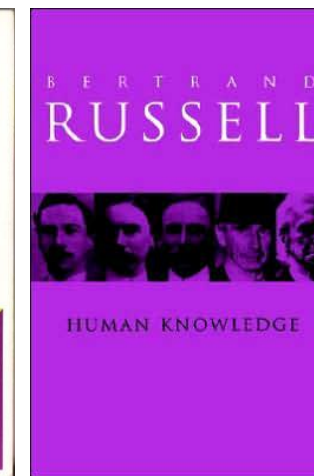
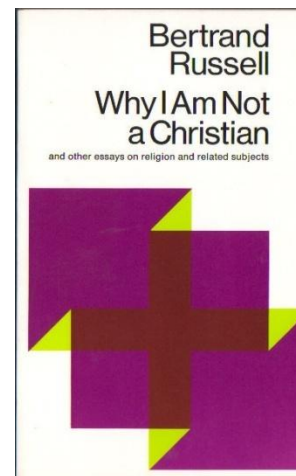
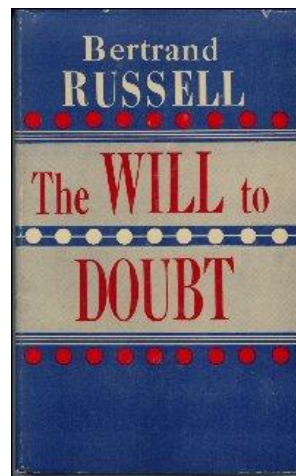
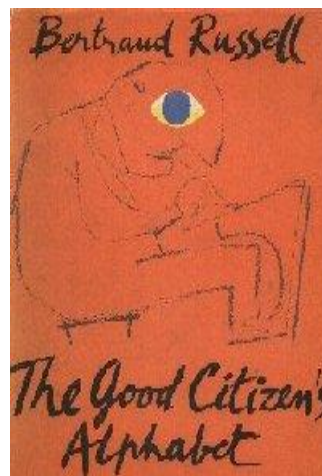
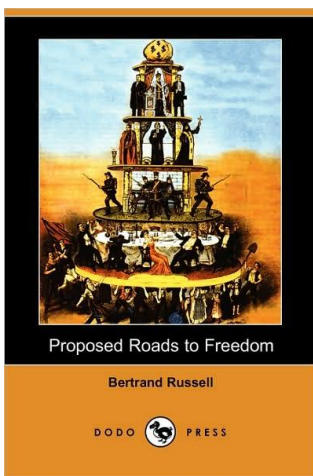
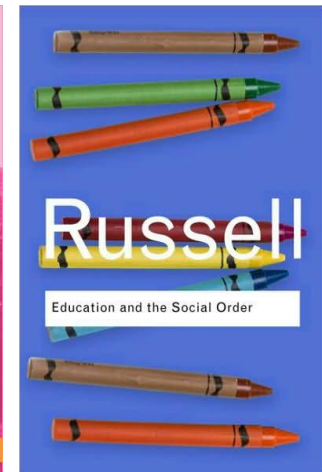
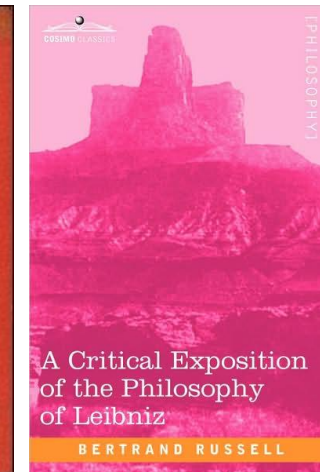
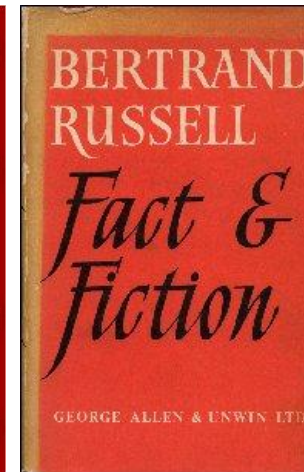
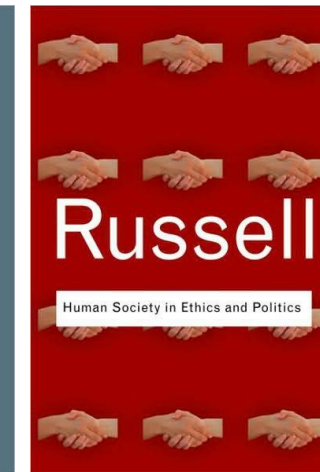
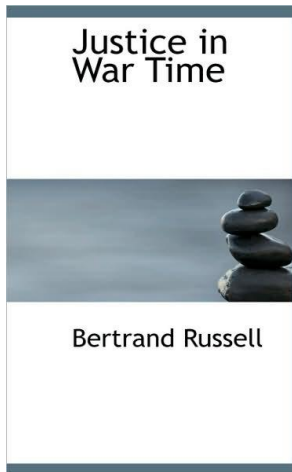
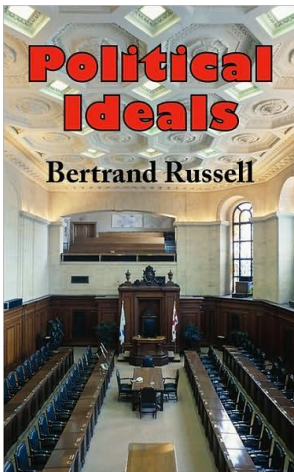
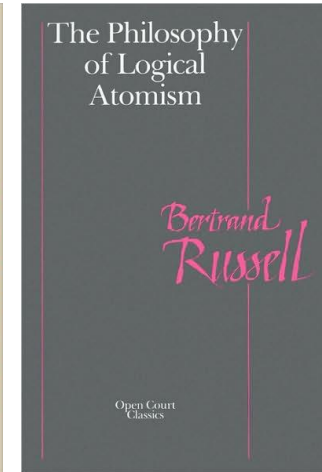
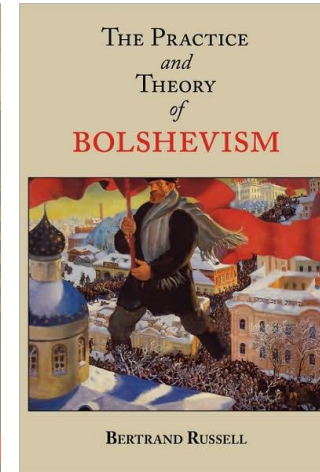
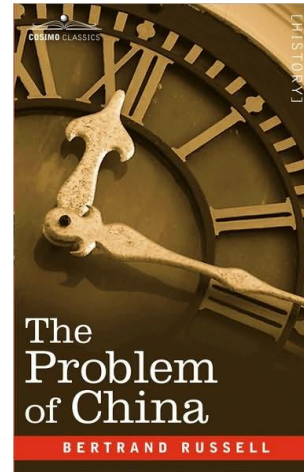
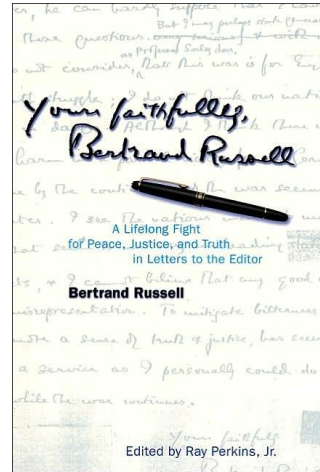
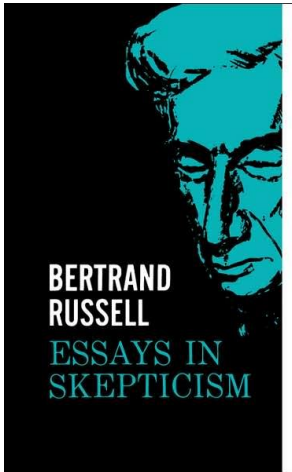
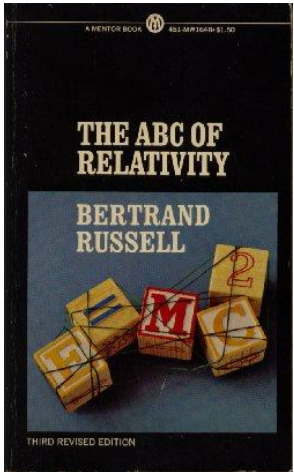
ion On Edu
Education
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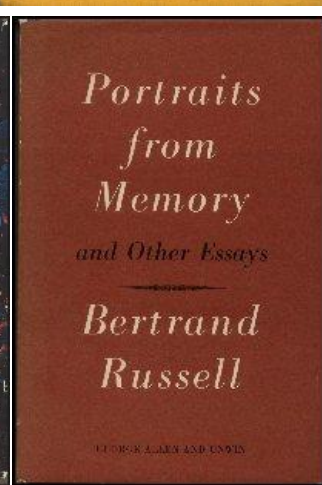
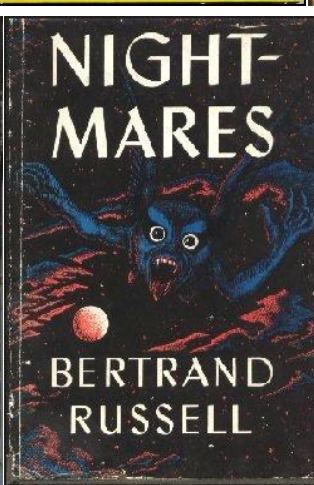
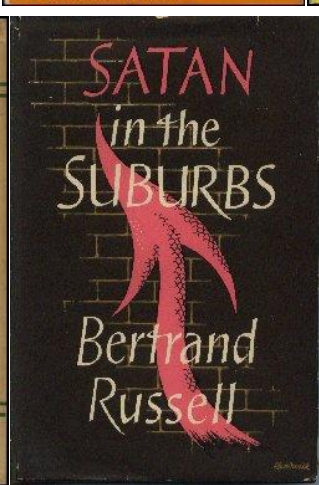
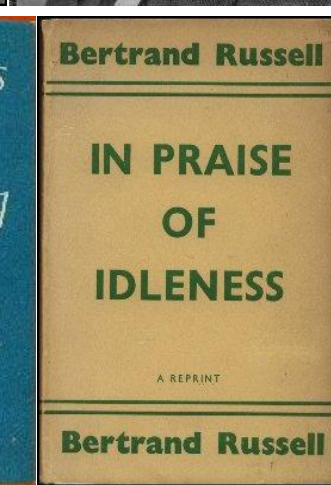
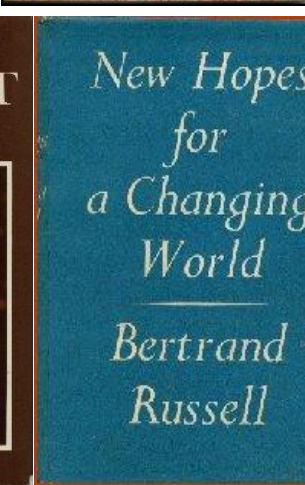
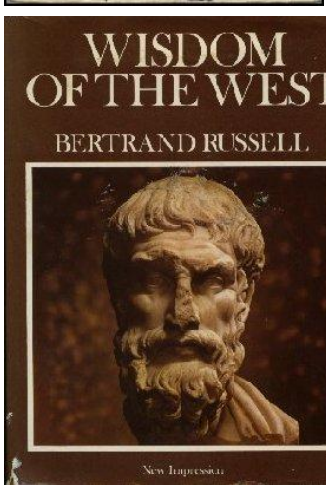
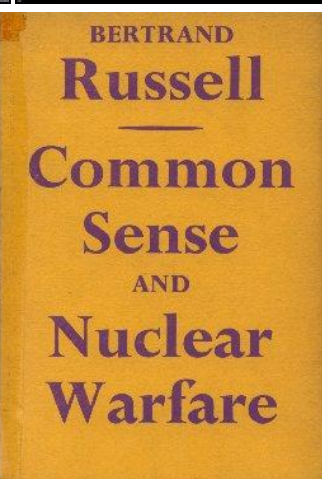
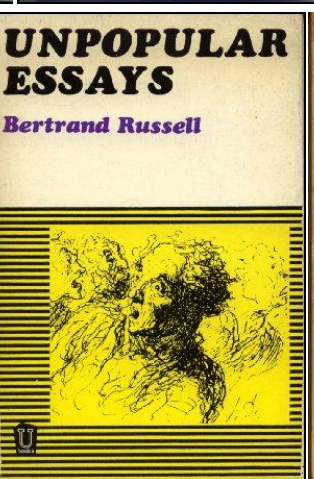
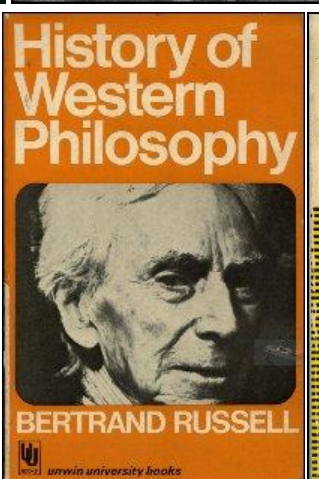
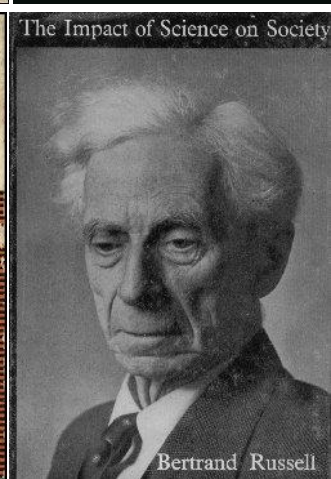
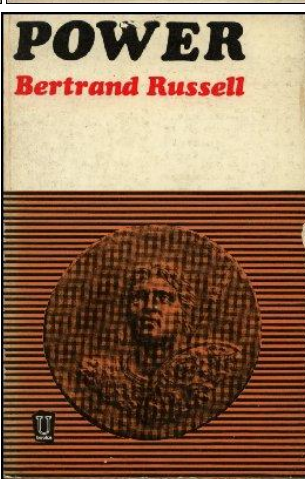
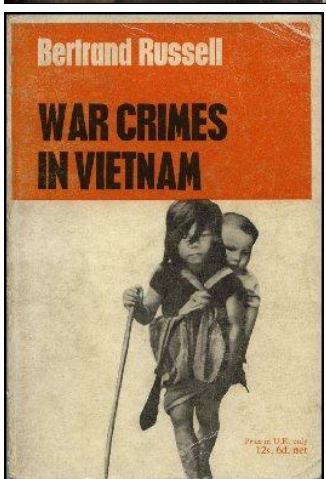
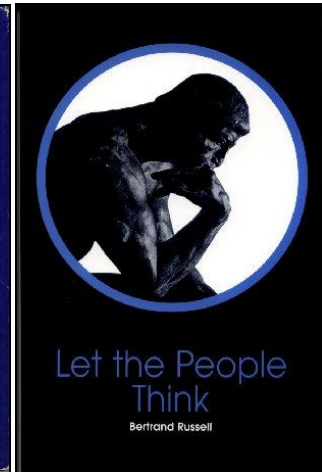
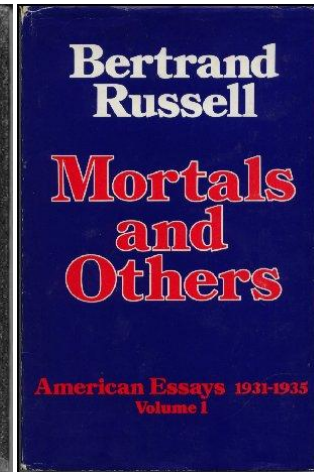
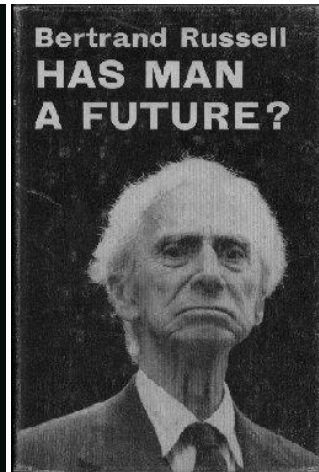
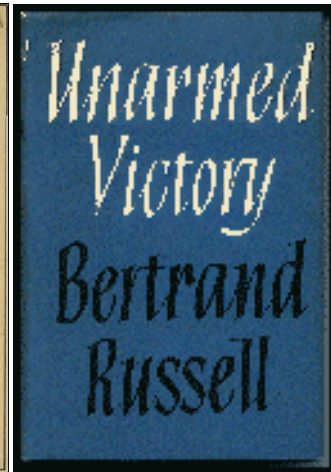
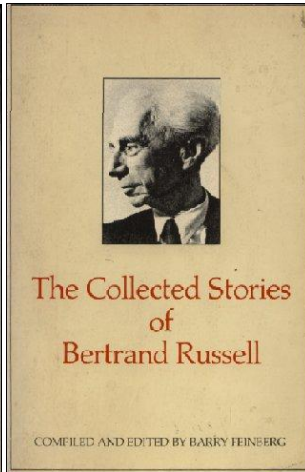
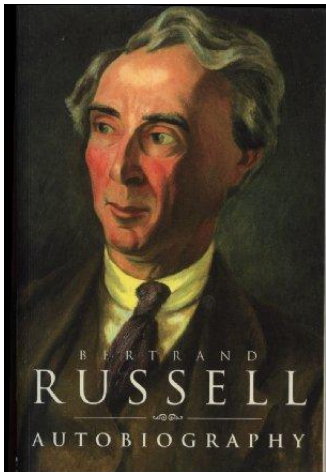


MYSTICISM AND LOGIC



BERTRAND RUSSELL





*5442. $\vdash :: \alpha \in 2. \supset :: \beta \subset \alpha. \mathfrak{A}! \beta. \beta \neq \alpha. \equiv . \beta \in \iota' \alpha$

Dem.

$\vdash . *544. \supset \vdash :: \alpha = \iota' x \cup \iota' y. \supset ::$

$\beta \subset \alpha. \mathfrak{A}! \beta. \equiv : \beta = \Lambda. \vee . \beta = \iota' x. \vee . \beta = \iota' y. \vee . \beta = \alpha : \mathfrak{A}! \beta :$

[*24:53:56.*51:161] $\equiv : \beta = \iota' x. \vee . \beta = \iota' y. \vee . \beta = \alpha$ (1)

$\vdash . *54:25. \text{Transp.} *52:22. \supset \vdash : x \neq y. \supset . \iota' x \cup \iota' y \neq \iota' x. \iota' x \cup \iota' y \neq \iota' y :$

[*13:12] $\supset \vdash : \alpha = \iota' x \cup \iota' y. x \neq y. \supset . \alpha \neq \iota' x. \alpha \neq \iota' y$ (2)

$\vdash . (1). (2). \supset \vdash :: \alpha = \iota' x \cup \iota' y. x \neq y. \supset ::$

$\beta \subset \alpha. \mathfrak{A}! \beta. \beta \neq \alpha. \equiv : \beta = \iota' x. \vee . \beta = \iota' y :$

[*51:235] $\equiv : (\mathfrak{A}z). z \in \alpha. \beta = \iota' z :$

[*37:6] $\equiv : \beta \in \iota' \alpha$ (3)

$\vdash . (3). *11:11:35. *54:101. \supset \vdash . \text{Prop}$

*5443. $\vdash :: \alpha, \beta \in 1. \supset : \alpha \cap \beta = \Lambda. \equiv . \alpha \cup \beta \in 2$

Dem.

$\vdash . *54:26. \supset \vdash :: \alpha = \iota' x. \beta = \iota' y. \supset : \alpha \cup \beta \in 2. \equiv . x \neq y.$

[*51:231] $\equiv . \iota' x \cap \iota' y = \Lambda.$

[*13:12] $\equiv . \alpha \cap \beta = \Lambda$ (1)

$\vdash . (1). *11:11:35. \supset$

$\vdash :: (\mathfrak{A}x, y). \alpha = \iota' x. \beta = \iota' y. \supset : \alpha \cup \beta \in 2. \equiv . \alpha \cap \beta = \Lambda$ (2)

$\vdash . (2). *11:54. *52:1. \supset \vdash . \text{Prop}$

From this proposition it will follow, when arithmetical addition has been defined, that $1 + 1 = 2$.

*5444. $\vdash :: z, w \in \iota' x \cup \iota' y. \supset_{z,w} . \phi(z, w) : \equiv . \phi(x, x) . \phi(x, y) . \phi(y, x) . \phi(y, y)$

Dem.

$\vdash . *51:234. *11:62. \supset \vdash :: z, w \in \iota' x \cup \iota' y. \supset_{z,w} . \phi(z, w) : \equiv :$

$z \in \iota' x \cup \iota' y. \supset_z . \phi(z, x) . \phi(z, y) :$

[*51:234.*10:29] $\equiv : \phi(x, x) . \phi(x, y) . \phi(y, x) . \phi(y, y) : \supset \vdash . \text{Prop}$

*54441. $\vdash :: z, w \in \iota' x \cup \iota' y. z \neq w. \supset_{z,w} . \phi(z, w) : \equiv :: x = y : \vee : \phi(x, y) . \phi(y, x)$

Dem.

$\vdash . *5:6. \supset \vdash :: z, w \in \iota' x \cup \iota' y. z \neq w. \supset_{z,w} . \phi(z, w) : \equiv ::$

$z, w \in \iota' x \cup \iota' y. \supset_{z,w} : z = w. \vee . \phi(z, w) :$

[*54:44] $\equiv : x = x. \vee . \phi(x, x) : x = y. \vee . \phi(x, y) :$

$y = x. \vee . \phi(y, x) : y = y. \vee . \phi(y, y) :$

[*13:15] $\equiv : x = y. \vee . \phi(x, y) : y = x. \vee . \phi(y, x) :$

[*13:16.*4:41] $\equiv : x = y. \vee . \phi(x, y) . \phi(y, x)$

This proposition is used in *163:42, in the theory of relations of mutually exclusive relations.

*110632. $\vdash : \mu \in \text{NC}. \supset . \mu +_c 1 = \hat{\xi} \{ (\mathfrak{A}y) . y \in \xi . \xi - \iota' y \in \text{sm}'' \mu \}$

Dem.

$\vdash . *110:631. *51:211:22. \supset$

$\vdash : \text{Hp}. \supset . \mu +_c 1 = \hat{\xi} \{ (\mathfrak{A}y, y) . \gamma \in \text{sm}'' \mu . y \in \xi . \gamma = \xi - \iota' y \}$

[*13:195] $= \hat{\xi} \{ (\mathfrak{A}y) . y \in \xi . \xi - \iota' y \in \text{sm}'' \mu \} : \supset \vdash . \text{Prop}$

*11064. $\vdash . 0 +_c 0 = 0$ [*110:62]

*110641. $\vdash . 1 +_c 0 = 0 +_c 1 = 1$ [*110:51:61. *101:2]

*110642. $\vdash . 2 +_c 0 = 0 +_c 2 = 2$ [*110:51:61. *101:31]

*110643. $\vdash . 1 +_c 1 = 2$

Dem.

$\vdash . *110:632. *101:21:28. \supset$

$\vdash . 1 +_c 1 = \hat{\xi} \{ (\mathfrak{A}y) . y \in \xi . \xi - \iota' y \in 1 \}$

[*54:3] $= 2. \supset \vdash . \text{Prop}$

The above proposition is occasionally useful. It is used at least three times, in *113:66 and *120:123:472.

*110:7:71 are required for proving *110:72, and *110:72 is used in *117:3, which is a fundamental proposition in the theory of greater and less.

*1107. $\vdash : \beta \subset \alpha. \supset . (\mathfrak{A}\mu) . \mu \in \text{NC}. \text{Nc}' \alpha = \text{Nc}' \beta +_c \mu$

Dem.

$\vdash . *24:411:21. \supset \vdash : \text{Hp}. \supset . \alpha = \beta \cup (\alpha - \beta) . \beta \cap (\alpha - \beta) = \Lambda.$

[*110:32] $\supset . \text{Nc}' \alpha = \text{Nc}' \beta +_c \text{Nc}' (\alpha - \beta) : \supset \vdash . \text{Prop}$

*11071. $\vdash : (\mathfrak{A}\mu) . \text{Nc}' \alpha = \text{Nc}' \beta +_c \mu. \supset . (\mathfrak{A}\delta) . \delta \text{ sm } \beta . \delta \subset \alpha$

Dem.

$\vdash . *100:3. *110:4. \supset$

$\vdash : \text{Nc}' \alpha = \text{Nc}' \beta +_c \mu. \supset . \mu \in \text{NC} - \iota' \Lambda$ (1)

$\vdash . *110:3. \supset \vdash : \text{Nc}' \alpha = \text{Nc}' \beta +_c \text{Nc}' \gamma. \equiv . \text{Nc}' \alpha = \text{Nc}' (\beta + \gamma).$

[*100:3:31] $\supset . \alpha \text{ sm } (\beta + \gamma).$

[*73:1] $\supset . (\mathfrak{A}R) . R \in 1 \rightarrow 1 . D' R = \alpha . \mathfrak{A}' R = \downarrow \Lambda \gamma'' \iota'' \beta \cup \Lambda \beta \downarrow \iota'' \gamma'' .$

[*37:15] $\supset . (\mathfrak{A}R) . R \in 1 \rightarrow 1 . \downarrow \Lambda \gamma'' \iota'' \beta \subset \mathfrak{A}' R . R'' \downarrow \Lambda \gamma'' \iota'' \beta \subset \alpha .$

[*110:12.*73:22] $\supset . (\mathfrak{A}\delta) . \delta \subset \alpha . \delta \text{ sm } \beta$ (2)

$\vdash . (1). (2). \supset \vdash . \text{Prop}$

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Theorem pm54.43 4699

Description: Theorem *54.43 of [WhiteheadRussell] p. 360. "From this proposition it will follow, when arithmetical addition has been defined, that $1+1=2$." See http://en.wikipedia.org/wiki/Principia_Mathematica#Quotations. This theorem states that two sets of cardinality 1 are disjoint iff their union has cardinality 2.

Whitehead and Russell define 1 as the collection of all sets with cardinality 1 (i.e. all singletons; see [card1](#) 4965), so that their $A \in 1$ means, in our notation, $A \in \{x \mid (\text{card}'x) = 1_o\}$ i.e. $(\text{card}'A) = 1_o$ (by [elab](#) 1939) i.e. $A \approx 1_o$ (by [carden](#) 4963 and [cardm](#) 4954). We do not have several of their earlier lemmas available (which would otherwise be unused by our different approach to arithmetic), so our proof is longer. (It is also longer because we must show every detail.)

Theorem [pml10.643](#) 5057 shows the derivation of $1+1=2$ for cardinal numbers from this theorem.

Assertion

Ref	Expression
pm54.43	$\vdash ((A \approx 1_o \wedge B \approx 1_o) \rightarrow ((A \cap B) = \emptyset \leftrightarrow (A \cup B) \approx 2_o))$

Proof of Theorem pm54.43

Step	Hyp	Ref	Expression
1		lon 4262	$\vdash 1_o \in On$
2	1	onurri 3006	$\vdash \neg 1_o \in 1_o$
3		disjsn 2493	$\vdash ((1_o \cap \{1_o\}) = \emptyset \leftrightarrow \neg 1_o \in 1_o)$
4	2, 3	mpbir 188	$\vdash (1_o \cap \{1_o\}) = \emptyset$
5		unen 4563	$\vdash (((A \approx 1_o \wedge B \approx \{1_o\}) \wedge ((A \cap B) = \emptyset \wedge (1_o \cap \{1_o\}) = \emptyset)) \rightarrow (A \cup B) \approx (1_o \cup \{1_o\}))$
6	4, 5	mpanr2 713	$\vdash (((A \approx 1_o \wedge B \approx \{1_o\}) \wedge (A \cap B) = \emptyset) \rightarrow (A \cup B) \approx (1_o \cup \{1_o\}))$
7	6	ex 371	$\vdash ((A \approx 1_o \wedge B \approx \{1_o\}) \rightarrow ((A \cap B) = \emptyset \rightarrow (A \cup B) \approx (1_o \cup \{1_o\})))$
8	1	elisseti 1860	$\vdash 1_o \in V$
9	8	ensnl 4553	$\vdash \{1_o\} \approx 1_o$
10	8, 9	ensymi 4542	$\vdash 1_o \approx \{1_o\}$
11		entr 4543	$\vdash ((B \approx 1_o \wedge 1_o \approx \{1_o\}) \rightarrow B \approx \{1_o\})$
12	10, 11	mpan2 699	$\vdash (B \approx 1_o \rightarrow B \approx \{1_o\})$
13	7, 12	sylan2 453	$\vdash ((A \approx 1_o \wedge B \approx 1_o) \rightarrow ((A \cap B) = \emptyset \rightarrow (A \cup B) \approx (1_o \cup \{1_o\})))$
14		df-2o 4258	$\vdash 2_o = \text{suc } 1_o$
15		df-suc 2971	$\vdash \text{suc } 1_o = (1_o \cup \{1_o\})$
16	14, 15	eqtri 1534	$\vdash 2_o = (1_o \cup \{1_o\})$
17	16	breq2i 2691	$\vdash ((A \cup B) \approx 2_o \leftrightarrow (A \cup B) \approx (1_o \cup \{1_o\}))$
18	13, 17	syl6ibr 211	$\vdash ((A \approx 1_o \wedge B \approx 1_o) \rightarrow ((A \cap B) = \emptyset \rightarrow (A \cup B) \approx 2_o))$

19		sneq 2466 15 $\vdash (x = y \rightarrow \{x\} = \{y\})$
20	19	uneq2d 2232 14 $\vdash (x = y \rightarrow (\{x\} \cup \{x\}) = (\{x\} \cup \{y\}))$
21		unidm 2223 14 $\vdash (\{x\} \cup \{x\}) = \{x\}$
22	20, 21	syl5reqr 1561 13 $\vdash (x = y \rightarrow (\{x\} \cup \{y\}) = \{x\})$
23		visset 1855 15 $\vdash x \in V$
24	23	ensn1 4553 14 $\vdash \{x\} \approx 1_o$
25		l1dom2 4656 14 $\vdash 1_o \prec 2_o$
26		ensdomtr 4600 14 $\vdash (((\{x\} \approx 1_o \wedge 1_o \prec 2_o) \rightarrow \{x\} \prec 2_o)$
27	24, 25, 26	mp2an 700 13 $\vdash \{x\} \prec 2_o$
28	22, 27	syl6eqbr 2716 12 $\vdash (x = y \rightarrow (\{x\} \cup \{y\}) \prec 2_o)$
29		sdomnen 4516 12 $\vdash (((\{x\} \cup \{y\}) \prec 2_o \rightarrow \neg (\{x\} \cup \{y\}) \approx 2_o)$
30	28, 29	syl 10 11 $\vdash (x = y \rightarrow \neg (\{x\} \cup \{y\}) \approx 2_o)$
31	30	necon2ai 1650 10 $\vdash (((\{x\} \cup \{y\}) \approx 2_o \rightarrow x \neq y)$
32		disjsn2 2494 10 $\vdash (x \neq y \rightarrow (\{x\} \cap \{y\}) = \emptyset)$
33	31, 32	syl 10 9 $\vdash (((\{x\} \cup \{y\}) \approx 2_o \rightarrow (\{x\} \cap \{y\}) = \emptyset)$
34	33	alim 8 8 $\vdash ((A = \{x\} \wedge B = \{y\}) \rightarrow (((\{x\} \cup \{y\}) \approx 2_o \rightarrow (\{x\} \cap \{y\}) = \emptyset))$
35		uneq12 2227 9 $\vdash ((A = \{x\} \wedge B = \{y\}) \rightarrow (A \cup B) = (\{x\} \cup \{y\}))$
36	35	breql2 2693 8 $\vdash ((A = \{x\} \wedge B = \{y\}) \rightarrow ((A \cup B) \approx 2_o \leftrightarrow (\{x\} \cup \{y\}) \approx 2_o)$
37		ineq12 2260 9 $\vdash ((A = \{x\} \wedge B = \{y\}) \rightarrow (A \cap B) = (\{x\} \cap \{y\}))$
38	37	equeql1 1522 8 $\vdash ((A = \{x\} \wedge B = \{y\}) \rightarrow ((A \cap B) = \emptyset \leftrightarrow (\{x\} \cap \{y\}) = \emptyset)$
39	34, 36, 38	3imtr4d 545 7 $\vdash ((A = \{x\} \wedge B = \{y\}) \rightarrow ((A \cup B) \approx 2_o \rightarrow (A \cap B) = \emptyset)$
40	39	ex 371 6 $\vdash (A = \{x\} \rightarrow (B = \{y\} \rightarrow ((A \cup B) \approx 2_o \rightarrow (A \cap B) = \emptyset)))$
41	40	19.23adv 1248 5 $\vdash (A = \{x\} \rightarrow (\exists y B = \{y\} \rightarrow ((A \cup B) \approx 2_o \rightarrow (A \cap B) = \emptyset)))$
42	41	19.23aiv 1330 4 $\vdash (\exists x A = \{x\} \rightarrow (\exists y B = \{y\} \rightarrow ((A \cup B) \approx 2_o \rightarrow (A \cap B) = \emptyset)))$
43	42	imp 348 3 $\vdash ((\exists x A = \{x\} \wedge \exists y B = \{y\}) \rightarrow ((A \cup B) \approx 2_o \rightarrow (A \cap B) = \emptyset)$
44		enl 4555 3 $\vdash (A \approx 1_o \leftrightarrow \exists x A = \{x\})$
45		enl 4555 3 $\vdash (B \approx 1_o \leftrightarrow \exists y B = \{y\})$
46	43, 44, 45	syl2anb 457 2 $\vdash ((A \approx 1_o \wedge B \approx 1_o) \rightarrow ((A \cup B) \approx 2_o \rightarrow (A \cap B) = \emptyset))$
47	18, 46	impbid 518 1 $\vdash ((A \approx 1_o \wedge B \approx 1_o) \rightarrow ((A \cap B) = \emptyset \leftrightarrow (A \cup B) \approx 2_o))$

Colors of variables: wff set class

Syntax hints: \neg wn 2 \rightarrow wi 3 \leftrightarrow wb 144 \wedge wa 221 $=$ wceq 989 \in wcel 991 \exists wex 1013 \neq wne 1624 \cup cun 2093 \cap cin 2094 \emptyset c0 2328 $\{$ csn 2458 *class class class* wbr 2683 \cap con0 2965 *suc* csuc 2967
 1_o clo 4252 2_o c2o 4253 \approx cen 4493 \prec csdm 4495

This theorem is referenced by: [pml10.643](#) 5057 [unpde2eg2](#) 10809

This theorem was proved from axioms: [ax-1](#) 4 [ax-2](#) 5 [ax-3](#) 6 [ax-mp](#) 7 [ax-7](#) 965 [ax-gen](#) 996 [ax-8](#) 997 [ax-9](#) 998 [ax-10](#) 999 [ax-11](#) 1000 [ax-12](#) 1001 [ax-13](#) 1002 [ax-14](#) 1003 [ax-17](#) 1004 [ax-4](#) 1006 [ax-5o](#) 1008 [ax-6o](#) 1011
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Theorem pm110.643 ⁵⁰⁵⁷

Description: $1+1=2$ for cardinal number addition. Theorem *110.643 of *Principia Mathematica*, vol. II, p. 86, which adds the remark, "The above proposition is occasionally useful." Unlike us, Whitehead and Russell define cardinal addition on collections of all sets equinumerous to 1 and 2 (which for us are proper classes unless we restrict them as in [karden](#) 4856), but after applying definitions, our theorem is equivalent. See also the comment for [pm54.43](#) 4699. The comment for [cdavali](#) 5054 explains why we use \approx instead of $=$.

Assertion

Ref	Expression
pm110.643	$1 \vdash (1_o +_c 1_o) \approx 2_o$

Proof of Theorem pm110.643

Step	Hyp	Ref	Expression
1		lon 4262	$1 \vdash 1_o \in On$
2	1	elisseti 1860	$1 \vdash 1_o \in V$
3	2, 2	cdavali 5054	$2 \vdash (1_o +_c 1_o) = ((1_o \times \{\emptyset\}) \cup (1_o \times \{1_o\}))$
4		xp01disj 4267	$3 \vdash ((1_o \times \{\emptyset\}) \cap (1_o \times \{1_o\})) = \emptyset$
5		0ex 2777	$5 \vdash \emptyset \in V$
6	2, 5	xpsnen 4564	$4 \vdash (1_o \times \{\emptyset\}) \approx 1_o$
7	2, 2	xpsnen 4564	$4 \vdash (1_o \times \{1_o\}) \approx 1_o$
8		pm54.43 4699	$4 \vdash (((1_o \times \{\emptyset\}) \approx 1_o \wedge (1_o \times \{1_o\}) \approx 1_o) \rightarrow (((1_o \times \{\emptyset\}) \cap (1_o \times \{1_o\})) = \emptyset \leftrightarrow ((1_o \times \{\emptyset\}) \cup (1_o \times \{1_o\})) \approx 2_o)$
9	6, 7, 8	mp2an 700	$3 \vdash (((1_o \times \{\emptyset\}) \cap (1_o \times \{1_o\})) = \emptyset \leftrightarrow ((1_o \times \{\emptyset\}) \cup (1_o \times \{1_o\})) \approx 2_o)$
10	4, 9	mpbi 187	$2 \vdash ((1_o \times \{\emptyset\}) \cup (1_o \times \{1_o\})) \approx 2_o$
11	3, 10	eqbrtri 2698	$1 \vdash (1_o +_c 1_o) \approx 2_o$

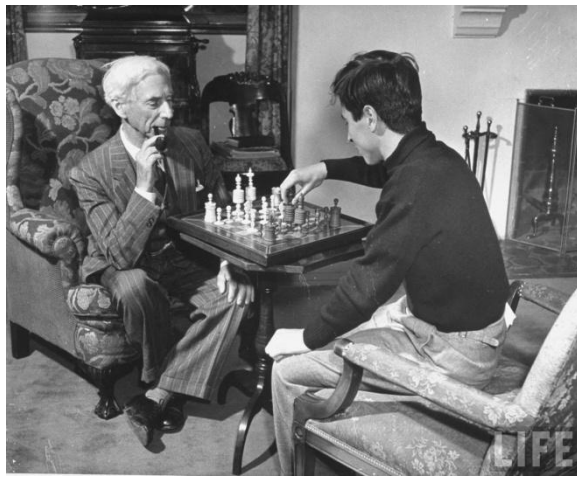
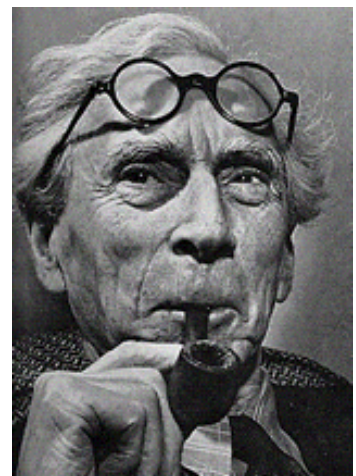
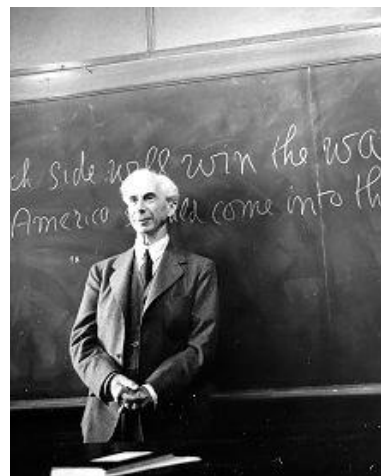
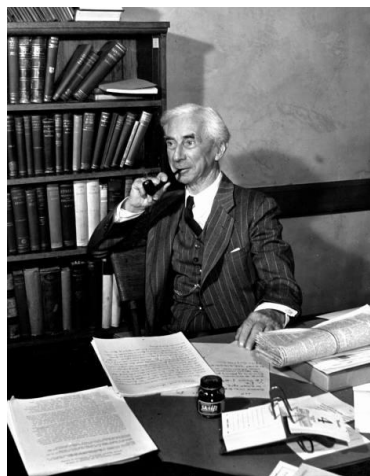
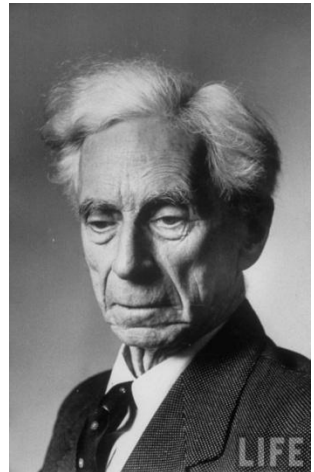
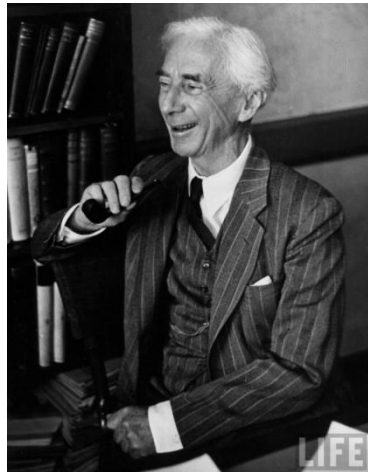
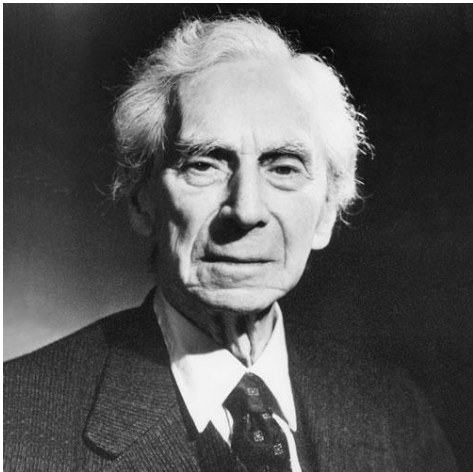
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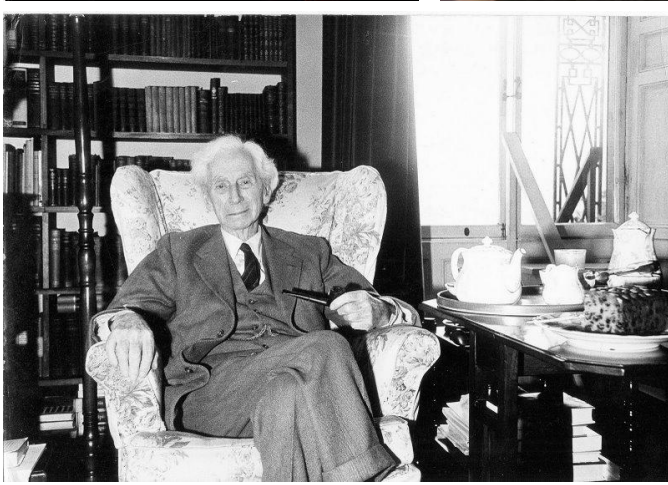
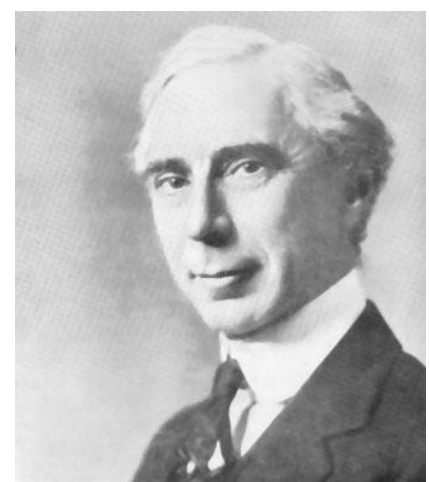
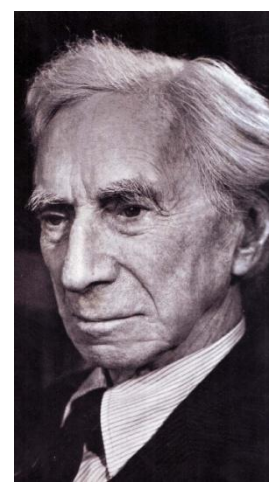
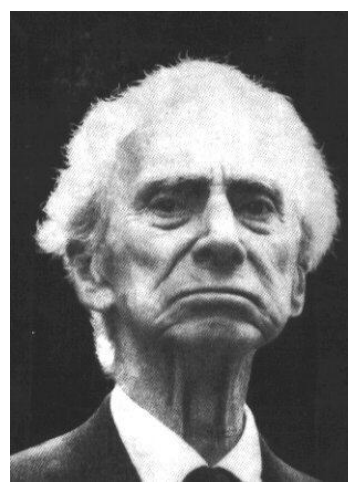
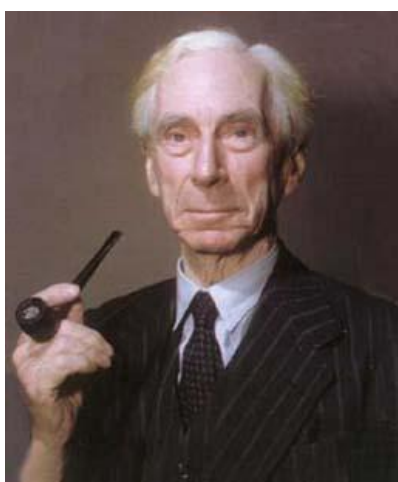
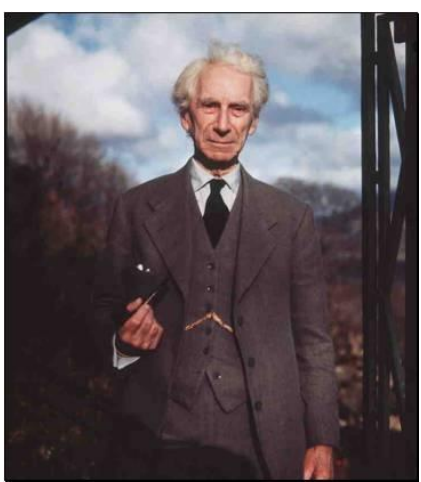
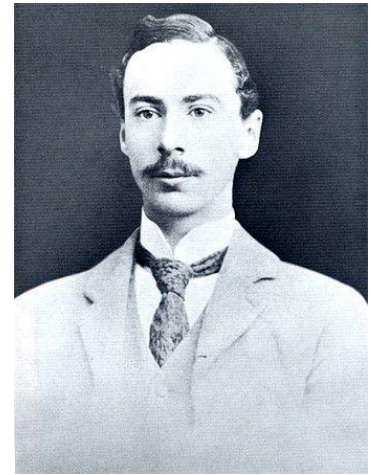
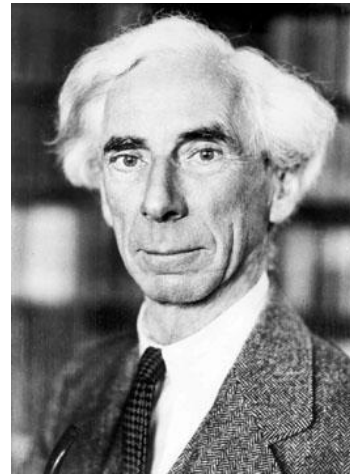
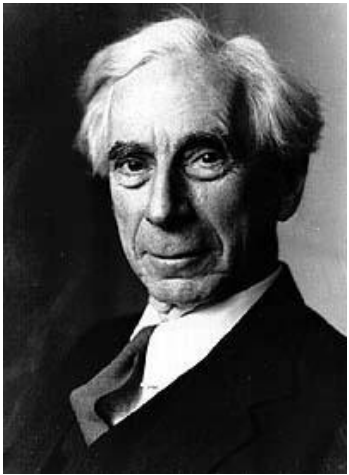
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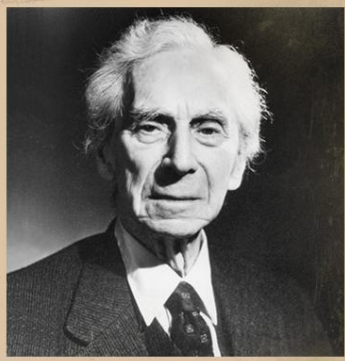
This theorem was proved from axioms: [ax-1](#) 4 [ax-2](#) 5 [ax-3](#) 6 [ax-mp](#) 7 [ax-7](#) 995 [ax-gen](#) 996 [ax-8](#) 997 [ax-9](#) 998 [ax-10](#) 999 [ax-11](#) 1000 [ax-12](#) 1001 [ax-13](#) 1002 [ax-14](#) 1003 [ax-17](#) 1004 [ax-4](#) 1006 [ax-5o](#) 1008 [ax-6o](#) 1011 [ax-9o](#) 1156 [ax-10o](#) 1174 [ax-16](#) 1244 [ax-11o](#) 1252 [ax-ext](#) 1496 [ax-rep](#) 2759 [ax-sep](#) 2769 [ax-nul](#) 2776 [ax-pow](#) 2809 [ax-pr](#) 2844 [ax-un](#) 3079

This theorem depends on definitions: [df-bi](#) 145 [df-or](#) 222 [df-an](#) 223 [df-3or](#) 779 [df-3an](#) 780 [df-ex](#) 1014 [df-sb](#) 1206 [df-eu](#) 1417 [df-mo](#) 1418 [df-clab](#) 1502 [df-cleq](#) 1507 [df-clel](#) 1510 [df-ne](#) 1626 [df-ral](#) 1691 [df-rex](#) 1692 [df-reu](#) 1693 [df-rab](#) 1694 [df-v](#) 1854 [df-sbc](#) 1983 [df-csb](#) 2048 [df-dif](#) 2097 [df-un](#) 2098 [df-in](#) 2099 [df-ss](#) 2101 [df-nul](#) 2329 [df-pw](#) 2451 [df-sn](#) 2461 [df-pr](#) 2462 [df-tp](#) 2464 [df-op](#) 2465 [df-uni](#) 2561 [df-int](#) 2592 [df-br](#) 2684 [df-opab](#) 2732 [df-tr](#) 2746 [df-eprel](#) 2899 [df-id](#) 2902 [df-po](#) 2907 [df-so](#) 2919 [df-fi](#) 2937 [df-we](#) 2952 [df-ord](#) 2968 [df-on](#) 2969 [df-suc](#) 2971 [df-xp](#) 3255 [df-rel](#) 3256 [df-cnv](#) 3257 [df-co](#) 3258 [df-dm](#) 3259 [df-m](#) 3260 [df-res](#) 3261 [df-ima](#) 3262 [df-fun](#) 3263 [df-fi](#) 3264 [df-f](#) 3265 [df-fl](#) 3266 [df-fo](#) 3267 [df-flo](#) 3268 [df-fv](#) 3269 [df-opr](#) 4011 [df-oprab](#) 4012 [df-lo](#) 4257 [df-2o](#) 4258 [df-er](#) 4389 [df-en](#) 4497 [df-dom](#) 4498 [df-sdom](#) 4499 [df-cda](#) 5052

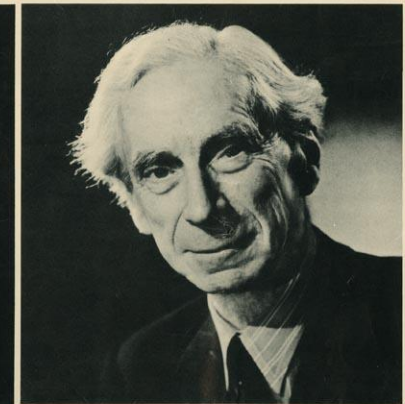
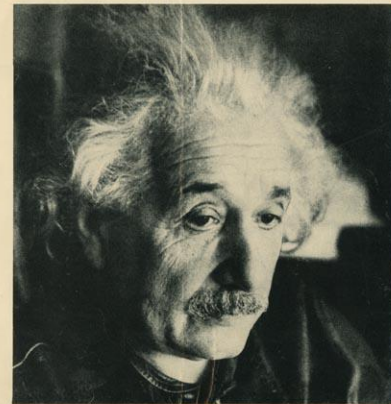
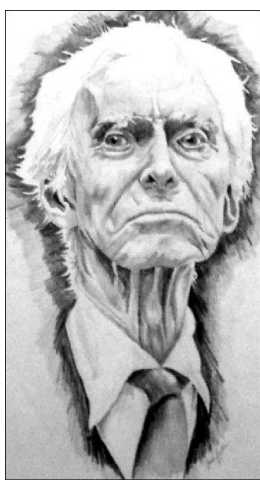
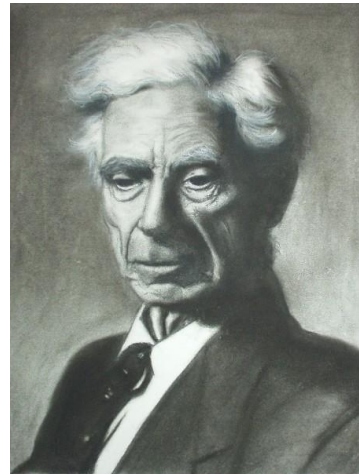
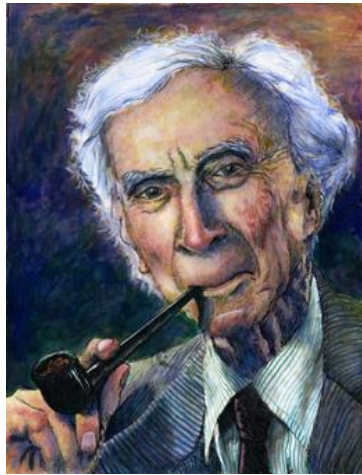
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Bertrand Russell

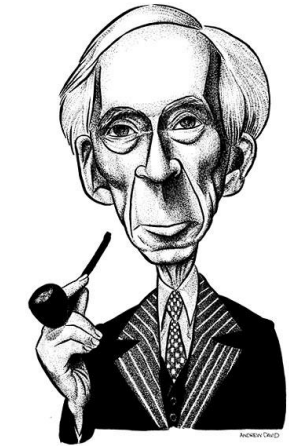
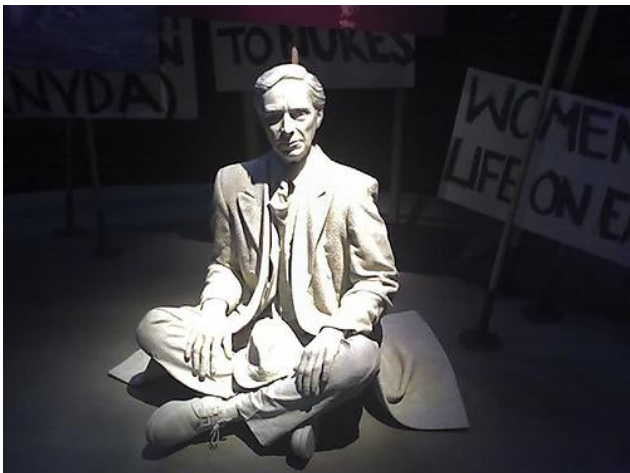


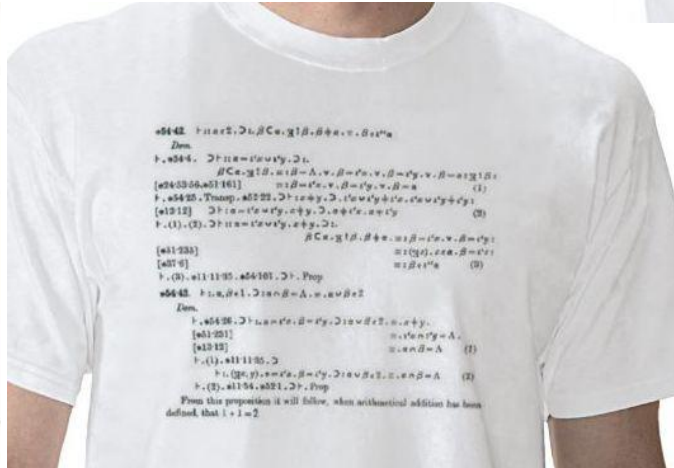
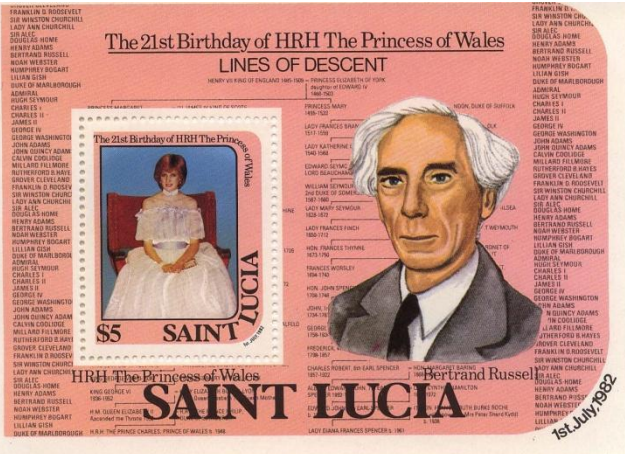
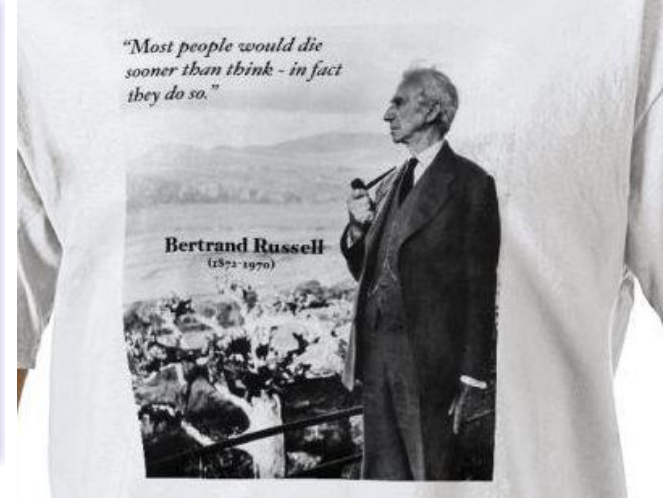
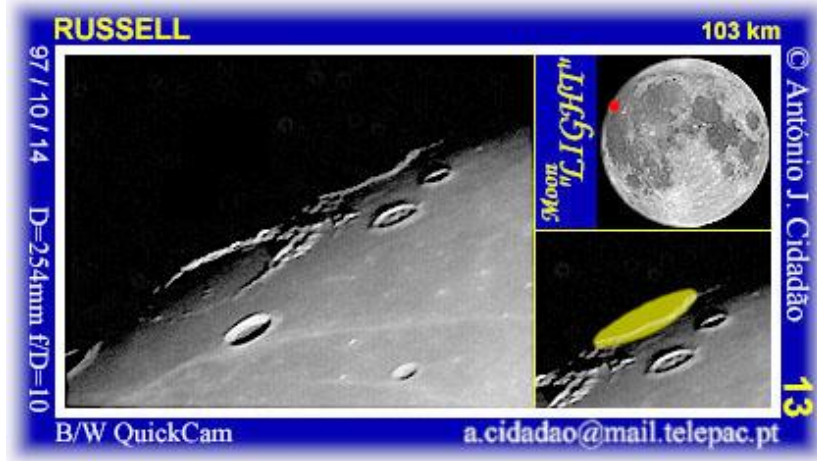
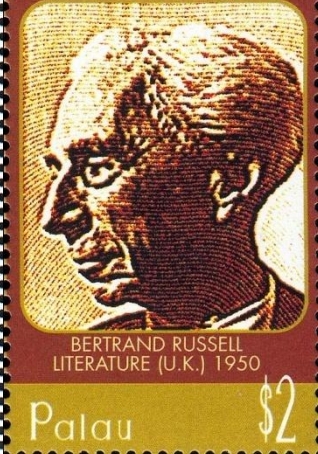
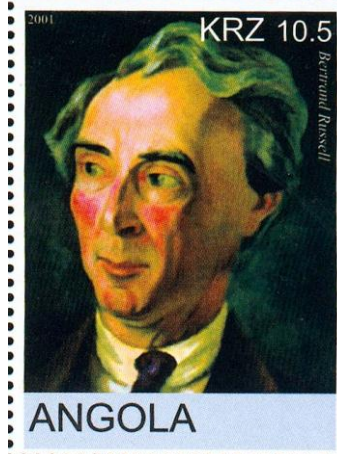
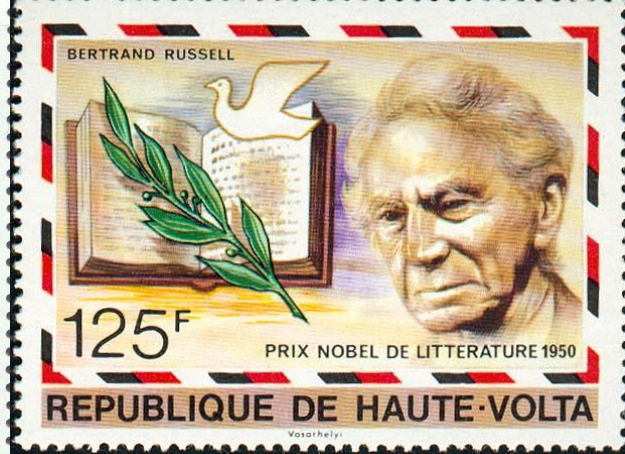
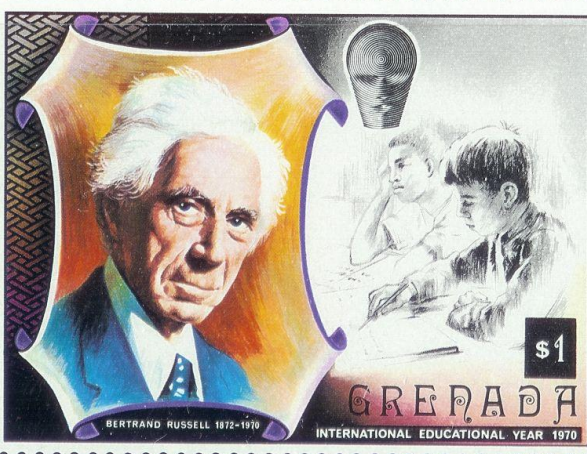
Albert Einstein *Bertrand Russell*

**NOTICE
TO THE WORLD**

*... renounce war or perish!
... world peace or universal death!*

AUDIO MASTERWORKS LPA 1225





"Most people would die sooner than think; in fact, they do so."
 - Bertrand Russell (1872-1970)

Russell's paradox was invented by Russell in 1901 to show that naïve set theory is self-contradictory:

Define: set of all sets that **do not contain themselves**

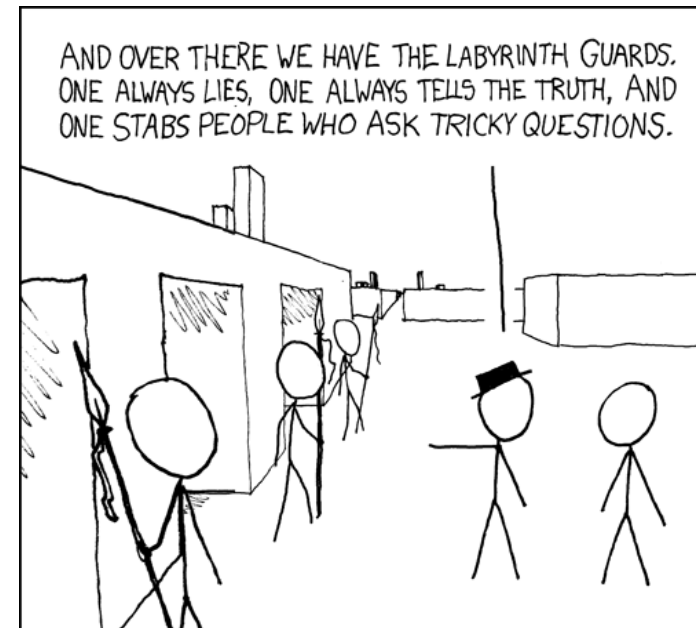
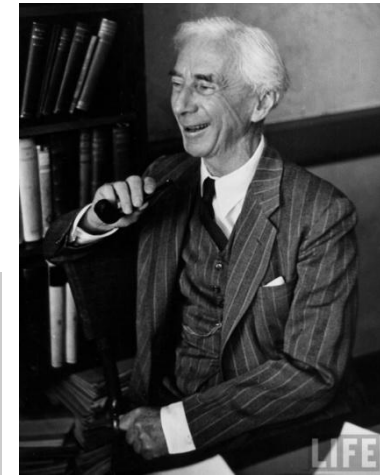
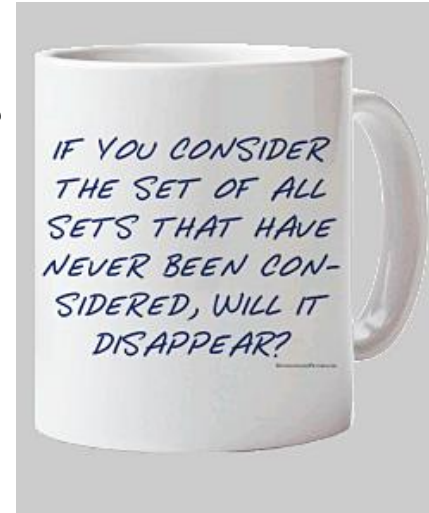
$$S = \{ T \mid T \notin T \}$$

Q: does S contain itself as an element?

$$S \notin S \Leftrightarrow S \in S \quad \text{contradiction!}$$

Similar **paradoxes**:

- “**A barber who shaves exactly those who do not shave themselves.**”
- “This sentence is false.”
- “I am lying.”
- “Is the answer to this question ‘no’?”
- “The smallest positive integer not describable in twenty words or less.”

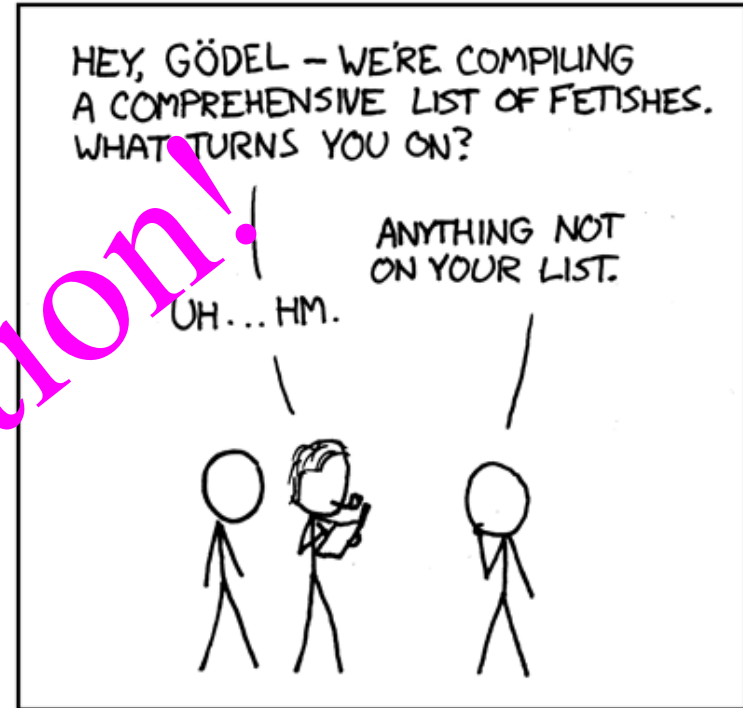




Star Trek, 1967, "I, Mudd" episode
Captain James Kirk and Harry Mudd use a logical paradox to cause hostile android "Norman" to crash

AUTHOR KATHARINE GATES RECENTLY ATTEMPTED TO MAKE A CHART OF ALL SEXUAL FETISHES.

LITTLE DID SHE KNOW THAT RUSSELL AND WHITEHEAD HAD ALREADY FAILED AT THIS SAME TASK.



Diagonalization!

