

# Algorithms

## Problem Set 3

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Please make all algorithms as efficient as you can, and state their time and space complexities.

1-46. Solve the following problems from the [Cormen, Third Edition, 2009] algorithms textbook:

- p. 456: 17.1-1, 17.1-2, 17.1-3
- p. 458-459: 17.2-1, 17.2-3
- p. 462-463: 17.3-2, 17.3-4, 17.3-6, 17.3-7
- p. 471: 17.4-1
- p. 473: 17-2
- p. 593: 22.1-5, 22.1-6, 22.1-7, 22.1-8,
- p. 602: 22.2-4, 22.2-7, 22.2-8, 22.2-9,
- p. 612: 22.3-12, 22.3-13
- p. 614-615: 22.4-2, 22.4-3, 22.4-4, 22.4-5
- p. 621: 22.5-7
- p. 623: 22-3, 22-4
- p. 629-630: 23.1-1, 23.1-5, 23.1-6, 23.1-7, 23.1-11
- p. 637-640: 23.2-4, 23.3-5, 23.2-6, 23.2-7, 23-1, 23-3,
- p. 655: 24.1-6
- p. 658: 24.2-4
- p. 663-664: 24.3-6, 24.3-8, 24.3-9
- p. 678-679: 24-2, 24-3,

- 47. Give an **optimal** algorithm to determine the diameter of a weighted tree (i.e., longest path between any two nodes). Prove the optimality of your algorithm. What is your time complexity?
- 48. Give an algorithm to determine whether any two of  $N$  segments on a line intersect.
- 49. Give an algorithm to find the  $N$ -way intersection of  $N$  rectangles with sides parallel to the axis.
- 50. Give an algorithm to compute the diameter of a given pointset (i.e., the distance between the farthest pair of points).
- 51. Give a linear-time algorithm to compute the area of a given convex polygon.
- 52. Give a linear-time algorithm to compute the area of a given non-convex polygon.
- 53. What is the worst-case cost of a traveling salesman tour over  $n$  points located arbitrarily in the unit square? Prove it. (Express the cost asymptotically as  $O(f(n))$  for some function  $f$ .)
- 54. What is the expected (average) cost of a minimum spanning tree over  $n$  points uniformly distributed in the unit square? (Express the cost asymptotically as  $O(f(n))$  for some function  $f$ .)

55. Given  $N$  **red** points and  $N$  **blue** points arbitrarily placed in the plane, we seek to match these  $2N$  points into  $N$  pairs, so that no two of the line segments defined by these pairs intersect. (Assume that no three of the  $2N$  distinct points are collinear.) Prove or disprove: there always exists such a non-self-intersecting matching for any arrangement of  $N$  red and  $N$  blue points arbitrarily located in the plane. Devise an efficient algorithm that given an arbitrary set of  $2N$  red & blue points in the plane, finds such a matching, when one exists. Generalize all this to dimension  $D$ .