This assignment asks you to prepare written answers to questions on garbage collection and exceptions. Each of the questions has a short answer. You may discuss this assignment with other students and work on the problems together. However, your write-up should be your own individual work.

*Please print your name and email address on your homework!* We need this information so that we can give you credit for the assignment and so that we can return it to you.

   
   (a) Assume that the garbage collector is run only when the user program runs out of memory, i.e. when a call to `new` cannot be satisfied. Is one of these two GC algorithms ‘faster’ than the other? Which algorithm needs to be run more frequently?

   (b) Does either algorithm use strictly more memory than the other?

   Python uses reference counting for its garbage collector. It uses a special ‘cycle detector’ to clean up cyclical data structures periodically.

   (c) Are reference cycles common in everyday data structures?

   (d) Briefly describe how one might implement a cycle detector. When can a cycle be cleaned?

2. Imagine that we have an updated version of Cool that supports `try`, `catch` and `throw` (as in the ‘Exceptions and Error Handling’ lecture notes). Suppose further that we want to add a new construct to the language: `protect e`, which works as follows:
• At compile time, the type checker verifies that the expression \( e \) does not have any uncaught exceptions. In other words, \( e \) can contain `throw` clauses, but they must all be caught using a `try / catch` construct. If \( e \) does not meet this requirement, the compiler should reject the program.

• At runtime, `protect` \( e \) simply evaluates to \( e \). Another way to look at this is that the `protect` keyword is ignored at runtime.

In order to check `protect` expressions at compile time, we need to extend our typing judgments to track a boolean \( E \) indicating whether an exception can be thrown out of the expression:

\[
O, M, C \vdash e : T, E
\]

Note that \( O \), \( M \), and \( C \) are unchanged. For example, the extended rule for `+` would be:

\[
\frac{O, M, C \vdash e_1 : Int, E_1 \quad O, M, C \vdash e_2 : Int, E_2}{O, M, C \vdash e_1 + e_2 : Int, E_1 \lor E_2} \quad \text{[Plus]}
\]

(a) Assume that a `try / catch` block catches all exceptions (i.e. the exception type being caught is always `Object`).

Give the new type rules for `protect`, `try / catch`, and `throw`. For example, the rule for `try / catch` should look similar to this:

\[
\cdots
\]

\[
O, M, C \vdash \text{try } e_0 \text{ catch } x : \text{Object} \Rightarrow e_1 : T_0 \sqcup T_1, E'
\]

(b) Suppose you wanted a conservative type system that was sound and safe (but might reject too many programs). How might you handle dynamic dispatch under this scheme? Think of the `throws` clause in Java.