History of Programming Languages

Functional Programming (1/2)
Cunning Plan

- History Lesson
- Functional Programming
  - OCaml
  - Types
  - Pattern Matching
  - Higher-Order Functions
- Basic Syntax
- Data Structures
- Higher-Order Functions
  - Fold
Gone In Sixty Seconds

- **Imperative**: change state, assignments
- **Structured**: if/block/routine control flow
- **Object-Oriented**: message passing (= dynamic dispatch), inheritance
- **Functional**: functions are first-class citizens that can be passed around or called recursively. We can avoid changing state by passing copies.
Why Study History?

• Those who cannot remember George Santayana are condemned to misquote him.
  - *Supernatural*, 1999
Why Study History?

- Progress, far from consisting in change, depends on retentiveness. **Those who cannot remember the past are condemned to repeat it.**

- Through meticulous analysis of history I will find a way to make the people worship me. By studying the conquerors of days gone by, I’ll discover the mistakes that made them go awry.
  - The Brain, *A Meticulous Analysis of History*, P031
Modern Era

- 1972 - C
- 1983 - Ada
- 1983 - C++
- 1987 - Perl
- 1990 - Python
- 1991 - Java
- 1993 - Ruby
- 1996 - OCaml
- 2000 - C#

Systems programming, ASM
US DOD, static type safety
classes, default args, STL
dynamic scripting language
interp OO + readability
portable OO lang (for iTV)
Perl + Smalltalk
ML + C++
“simple” Java + delegates

I invented the term Object-Oriented, and I did not have C++ in mind.
- Alan Kay
Time Travel

• Back to an earlier time when the US was worried about a Communist “perfect attack”
• In Soviet Russia, noun verbs you! (-1 Redundant)
The Land Before Time

- It was a time very different now ...

- **Senator Joseph McCarthy 1950**
  - “I have here in my hand a list of 205 — a list of names …”

- **John McCarthy 1958**
  - **LISP** = List Processing Language
  - basic datatype is the List, programs themselves are lists, can self-modify, dynamic allocation, garbage collection (!), functional
There are only two kinds of programming languages: those people always [complain] about and those nobody uses.
- Bjarne Stroustrup

I fear the new OO systems may suffer the fate of LISP, in that they can do many things, but the complexity of the class hierarchies may cause them to collapse under their own weight.
- Bill Joy

Computer language design is just like a stroll in the park. *Jurassic Park*, that is.
- Larry Wall

Oh what a tangled web we weave,
When first we practise to deceive!
- Sir Walter Scott, 1771-1832
Functional Programming

- You know OO and Structured Imperative
- **Functional Programming**
  - Computation = evaluating (math) functions
  - Avoid “global state” and “mutable data”
  - Get stuff done = apply (higher-order) functions
  - Avoid sequential commands
- Important Features
  - Higher-order, first-class functions
  - Closures and recursion
  - Lists and list processing
State

• The **state** of a program is all of the current variable and heap values
• **Imperative** programs destructively modify existing state

```
add_elem(SET, y)

SET {x}
```
State

• The state of a program is all of the current variable and heap values
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\[ SET \{x, y\} \]
State

• The **state** of a program is all of the current variable and heap values

• **Imperative** programs destructively modify existing state

• **Functional** programs yield new similar states over time

\[
\text{SET}_1 \{x\} \\
\text{SET}_2 = \text{add}_\text{elem}(\text{SET}_1, y)
\]
State

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\[ \text{SET}_2 = \text{add}\_\text{elem}(\text{SET}_1, y) \]
Basic OCaml

• Let's Start With C

double avg(int x, int y) {
    double z = (double)(x + y);
    z = z / 2;
    printf("Answer is %g\n", z);
    return z;
}

Basic OCaml

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define avg (x:int) (y:int) : float = begin
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let avg (x:int) (y:int) : float = begin
    let z = float_of_int (x + y) in
    end
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    let z = z /. 2.0 in
    printf "Answer is %g\n" z ;
    z
end
The Tuple (or Pair)

let x = (22, 58) in (* tuple creation *)
...
let y, z = x in (* tuple field extraction *)
printf “first element is %d\n” y ; ...

let add_points p1 p2 =
let x1, y1 = p1 in
let x2, y2 = p2 in
(x1 + x2, y1 + y2)
List Syntax in OCaml

- Empty List: `[ ]`
- Singleton: `[ element ]`
- Longer List: `[ e1 ; e2 ; e3 ]`
- Cons: `x :: [y;z] = [x;y;z]`
- Append: `[w;x]@[y;z] = [w;x;y;z]`
- List.length, List.filter, List.fold, List.map …
- More on these later!
- Every element in list **must have same type**
Functional Example

• Simple Functional Set (built out of lists)
  
  ```ml
  let rec add_elem (s, e) = 
    if s = [] then [e]
    else if List.hd s  = e then s
    else List.hd s :: add_elem(List.tl s, e)
  ```

• Pattern-Matching Functional (same effect)
  
  ```ml
  let rec add_elem (s,e) = match s with
    | [] -> [e]
    | hd :: tl when e = hd -> s
    | hd :: tl -> hd :: add_elem(tl, e)
  ```
Imperative Code

• More cases to handle

```c
List* add_elem(List *s, item e) {
    if (s == NULL)
        return list(e, NULL);
    else if (s->hd == e)
        return s;
    else if (s->tl == NULL) {
        s->tl = list(e, NULL); return s;
    } else
        return add_elem(s->tl, e);
}
```

I have stopped reading Stephen King novels. Now I just read C code instead.
- Richard O’Keefe
Q: Advertising (785 / 842)

• Identify the company associated with two of the following four advertising slogans or symbols.
  - "Fill it to the rim."
  - "I bet you can't eat just one."
  - "Snap, Crackle, Pop"
  - "The San Francisco Treat"
Correcting English Prose

1. It might be different if they had something in common, but they were on such polar ends of the pole.

4. Lizzy drank in the sight of him like a thirst craven man consumes water.

116. Wanda, seeing that, no, Kurt was not just going to leave her bags or things there, like she thought most boys would do, that he was, in fact, in laments terms, being a gentleman.

244. The pupils, who where only in high school grade, cheered for their known idles, inspired by their god-like techniques.

352. "Good evening my league." He picked her up by the wrist. "I think that you and I have some talking to do, actually I have a preposition"

Next time, you pick. Choose three numbers between 1 and 454, inclusive.
Functional-Style Advantages

• Tractable program semantics
  - Procedures are functions
  - Formulate and prove assertions about code
  - More readable

• Referential transparency
  - Replace any expression by its value without changing the result

• No side-effects
  - Fewer errors
Functional-Style Disadvantages

- Efficiency
  - Copying takes time

- Compiler implementation
  - Frequent memory allocation

- Unfamiliar (to you!)
  - New programming style

- Not appropriate for every program
  - Operating systems, etc.

<table>
<thead>
<tr>
<th>Language</th>
<th>Speed</th>
<th>Space</th>
</tr>
</thead>
<tbody>
<tr>
<td>C (gcc)</td>
<td>1.0</td>
<td>1.1</td>
</tr>
<tr>
<td>C++ (g++)</td>
<td>1.0</td>
<td>1.6</td>
</tr>
<tr>
<td>OCaml</td>
<td>1.5</td>
<td>2.9</td>
</tr>
<tr>
<td>Java (JDK -server)</td>
<td>1.7</td>
<td>9.1</td>
</tr>
<tr>
<td>Lisp</td>
<td>1.7</td>
<td>11</td>
</tr>
<tr>
<td>C# (mono)</td>
<td>2.4</td>
<td>5.6</td>
</tr>
<tr>
<td>Python</td>
<td>6.5</td>
<td>3.9</td>
</tr>
<tr>
<td>Ruby</td>
<td>16</td>
<td>5.0</td>
</tr>
</tbody>
</table>

17 small benchmarks
There are many ways of trying to understand programs. People often rely too much on one way, which is called “debugging” and consists of running a partly-understood program to see if it does what you expected. Another way, which ML advocates, is to install some means of understanding in the very programs themselves.

- Robin Milner, 1997
Type System

• Type Inference
  
  let rec add_elem (s,e) = match s with
  
  | [] -> [e]  
  | hd :: tl when e = hd -> s
  | hd :: tl -> hd :: add_elem(tl, e)
  
  - val add_elem : α list * α -> α list
  - “α list” means “List<T>” or “List<α>”

• ML infers types
  
  - Inconsistent or incomplete type is an error

• Optional type declarations  (exp : type)
  
  - Clarify ambiguous cases, documentation
Pattern Matching

- Simplifies Code (eliminates ifs, accessors)

```ocaml
type btree = (* binary tree of strings *)
  | Node of btree * string * btree
  | Leaf of string

let rec height tree = match tree with
  | Leaf _ -> 1
  | Node(x,_,y) -> 1 + max (height x) (height y)

let rec mem tree elt = match tree with
  | Leaf str -> str = elt
  | Node(x,str,y) -> str = elt ||
    mem x elt || mem y elt
```
Pattern Matching Mistakes

• What if I forget a case?
  
  - let rec is_odd x = match x with
    - | 0 -> false
    - | 2 -> false
    - | x when x > 2 -> is_odd (x-2)

  - Warning P: this pattern-matching is not exhaustive.

  - Here is an example of a value that is not matched: 1
Polymorphism

• Functions and type inference are polymorphic
  - Operate on more than one type
  
  ```
  let rec length x = match x with
  | [] -> 0
  | hd :: tl -> 1 + length tl
  ```

  - val length : \( \alpha \) list \( \to \) int

  length \([1;2;3]\) = 3
  length \[“algol”; ”smalltalk”; ”ml”\] = 3
  length \([1 ; “algol” ]\) = ?

\(\alpha\) means “any one type”
Higher-Order Functions

- Function are first-class values
  - Can be used whenever a value is expected
  - Notably, can be passed around
  - Closure captures the environment

```ocaml
let rec map f lst = match lst with
  | []   -> []
  | hd :: tl -> f hd :: map f tl
- val map : (α -> β) -> α list -> β list
```

```ocaml
let offset = 10 in
let myfun x = x + offset in
- val myfun : int -> int
map myfun [1;8;22] = [11;18;32]
```

- Extremely powerful programming technique
  - General iterators
  - Implement abstraction
The Story of Fold

• We’ve seen length and map

• We can also imagine ...
  - sum \([1; 5; 8]\) = 14
  - product \([1; 5; 8]\) = 40
  - and \([\text{true; true; false}]\) = false
  - or \([\text{true; true; false}]\) = true
  - filter \((\text{fun } x \rightarrow x>4)\) \([1; 5; 8]\) = \([5; 8]\)
  - reverse \([1; 5; 8]\) = \([8; 5; 1]\)
  - mem 5 \([1; 5; 8]\) = true

• Can we build all of these?
The House That Fold Built

• The **fold** operator comes from Recursion Theory (Kleene, 1952)

```ocaml
let rec fold f acc lst = match lst with
  | [] -> acc
  | hd :: tl -> fold f (f acc hd) tl
val fold : (α -> β -> α) -> α -> β list -> α
```

• Imagine we’re summing a list \( f = \text{addition} \):

\[
\begin{array}{c}
9 \rightarrow 2 \rightarrow 7 \rightarrow 4 \rightarrow 5 \rightarrow (\ldots) \\
\end{array}
\begin{array}{c}
acc \quad lst
\end{array}
\]

\[
\begin{array}{c}
11 \quad 7 \rightarrow 4 \rightarrow 5 \rightarrow (\ldots) \\
\end{array}
\begin{array}{c}
\text{f}
\end{array}
\]

\[
\begin{array}{c}
18 \rightarrow 4 \rightarrow 5 \rightarrow (\ldots) \\
\end{array}
\begin{array}{c}
\text{acc}
\end{array}
\]

\[
\begin{array}{c}
27 \\
\end{array}
\begin{array}{c}
\text{lst}
\end{array}
\]
It’s Lego Time

• Let’s build things out of Fold!
  - **length** lst =  **fold** (fun acc elt -> acc + 1) 0 lst
  - **sum** lst =  **fold** (fun acc elt -> acc + elt) 0 lst
  - **product** lst= **fold** (fun acc elt -> acc * elt) 1 lst
  - **and** lst =  **fold** (fun acc elt -> acc & elt) true lst

• How would we do **or**?
• How would we do **reverse**?
Tougher Legos

• Examples:
  - `reverse` `lst = fold (fun acc e -> acc @ [e]) [] lst`
    • Note typing: `(acc : α list) (e : α)`
  - `filter` `keep_it lst = fold (fun acc elt -> if keep_it elt then elt :: acc else acc) [] lst`
  - `mem` `wantedlst = fold (fun acc elt -> acc || wanted = elt) false lst`
    • Note typing: `(acc : bool) (e : α)`

• How do we do `map`?
  - Recall: `map (fun x -> x +10) [1;2] = [11;12]`
  - Let’s write it on the board ...
Map From Fold

- let map myfun lst =
-   fold (fun acc elt -> (myfun elt) :: acc) [] lst

  - Types: (myfun : α -> β)
  - Types: (lst : α list)
  - Types: (acc : β list)
  - Types: (elt : α)

- How do we do sort?
  - (sort : (α * α -> bool) -> α list -> α list)

Do nothing which is of no use.
- Miyamoto Musashi, 1584-1645
Sorting Examples

• langs = [ “fortran”; “algol”; “c” ]
• courses = [ 216; 333; 415]

- sort (fun a b -> a < b) langs
  - [ “algol”; “c”; “fortran” ]
- sort (fun a b -> a > b) langs
  - [ “fortran”; “c”; “algol” ]
- sort (fun a b -> strlen a < strlen b) langs
  - [ “c”; “algol”; “fortran” ]
- sort (fun a b -> match is_odd a, is_odd b with
  • | true, false -> true (* odd numbers first *)
  • | false, true -> false (* even numbers last *)
  • | _, _ -> a < b (* otherwise ascending *)
) courses
  - [ 333 ; 415 ; 216 ]

Java uses Inner Classes for this.

This last sort is a hint for PA0/1.
Partial Application and Currying

let myadd x y = x + y
• val myadd : int -> int -> int
myadd 3 5 = 8
let addtwo = myadd 2
  - How do we know what this means? We use referential transparency! Basically, just substitute it in.
• val addtwo : int -> int
addtwo 77 = 79
• **Currying**: “if you fix some arguments, you get a function of the remaining arguments”
• ML, Python and Ruby all support functional programming
  - closures, anonymous functions, etc.
• ML has strong static typing and type inference (as in this lecture)
• Ruby and Python have “strong” dynamic typing (or duck typing)
• All three combine OO and Functional
  - ... although it is rare to use both.
Homework

- Next: Backus Speedcoding
- Wed Sep 2: PA0 due
Pre-Game Warm-Up

• Punctuate and capitalize the following words to form a valid, sensical English paragraph:

that that is is that that is not is not is not that it