

lab post-mortem / HCL

Changelog

17 September: broken counter circuit: correct wire voltage diagrams to match values at time 0/1/2/3 better

last time

RISC v CISC

typical RISC properties

- separate memory access instructions, few ways to specify operands
- more registers (make up for separate memory access)
- simpler instructions generally
- hope: simple instructions are fast

Y86-64: kinda RISC-like variation on x86-64

Y86-64 encoding/decoding

the lab generally

yes, took a lot more time than my estimate

overall assessment: difficulty of parts of lab mostly okay (assuming issues with clarity of benchmark materials fixed), but too much stuff

lab part 1: yas

common misconception re: yas

basically produces **an executable**, not object file

yes, the extension `.yo` is confusing!

so, no symbol table+relocations — everything ready to load into memory

use in lab: to get correct machine code

...but *you* needed to handle relocations+symbols

lab part 1: relocations/symbol table entry

symbol table entries to remember where labels are:

```
foo: /* ←- label definition, becomes symbol table entry */
```

relocations to fill in addresses:

```
call foo /* ←- label use, needs relocation */  
jmp foo /* ←- label use, needs relocation */  
irmovq $foo, %rax /* ←- label use, needs relocation */
```

lab part 2: register availability

8-bit instructions, 3-bit opcode

instruction encoding with two operands:

opcode			rA		rB		???
0	1	2	3	4	5	6	7

e.g if opcode for add is 001, then add %r2, %r3
might become 001 10 11 0

2-bit register numbers → **4 total registers**

(maybe we should consider adding a bit to the opcode?)

lab part 2: register availability (alt)

8-bit instructions, 3-bit opcode

instruction encoding with three operands:

opcode			rA	rB	rC	???	
0	1	2	3	4	5	6	7

e.g. `add %r1, %r1, %r0` might become

`001 1 1 0 00`

1 bit register numbers → **2 total registers**

(probably, we should reconsider this instruction size)

lab part 3a: more instructions

question about: will compiler use more/less instructions with less registers

example of more instructions (**not what lab asked for**):
(function accumulating into a local variable, then returning it)

```
/* with less registers */  
...  
movq $0, 8(%rsp)  
... lots of computation ...  
pushq %r8  
movq 8(%rsp), %r8  
subq %r8, 128(%rdi)  
popq %r8  
... lots of computation ...  
addq %rax, 8(%rsp)  
movq 8(%rsp), %rax  
subq $16, %rsp  
ret
```

```
/* with more registers */  
...  
movq $0, %r15  
... lots of computation ...  
/* omitted instruction */  
/* omitted instruction */  
subq %r15, 128(%rdi)  
/* omitted instruction */  
... lots of computation ...  
addq %r15, %rax  
/* omitted instruction */  
subq $8, %rsp  
ret
```

lab part 3a: less instructions

question about: will compiler use more/less instructions with less registers

example of less instructions (but probably slower/longer machine code!)

(function accumulating into a global variable)

```
/* with less registers */  
/* omitted instruction */  
... lots of computation ...  
addq %rax, global_variable  
... lots of computation ...  
addq %rax, global_variable  
... lots of computation ...  
addq %rax, global_variable  
/* omitted instruction */  
ret
```

```
/* with more registers */  
movq global_variable, %r15  
... lots of computation ...  
addq %rax, %r15  
... lots of computation ...  
addq %rax, %r15  
... lots of computation ...  
addq %rax, %r15  
movq %r15, global_variable  
ret
```

lab part 3b (1)

anonymous feedback:

The expectation that we spend hours pouring over terribly formatted data just to do a long series of basic math operations as part of a hour lab is a waste of our time and teaches nothing.

I have a problem of opening text files in very wide text editors and not thinking about how it would look on smaller screens.

I probably should have added an “instructions that read or write memory” row to the data files

I probably should have written “instructions that don’t use memory” instead of “instructions that only use registers”

I expected (all instructions - instructions using memory)

lab part 3b (2)

anonymous feedback:

The expectation that we spend hours pouring over terribly formatted data just to do a long series of basic math operations as part of a hour lab is a waste of our time and teaches nothing.

Hope was to help with intuition for answering “does doing X help performance”

Also reveal that giving compilers a lot more registers isn't super helpful

Expected this to be a pretty quick part of the lab

Y86-64 decoding

```
20 10 60 20 61 37 72 84 00 00 00 00 00 00 00
20 12 20 01 70 68 00 00 00 00 00 00 00
```

byte:	0	1	2	3	4	5	6	7	8	9
halt	0	0								
nop	1	0								
rrmovq/cmovCC rA, rB	2	cc	rA	rB						
irmovq V, rB	3	0	F	rB	V					
rmmovq rA, D(rB)	4	0	rA	rB	D					
mrmovq D(rB), rA	5	0	rA	rB	D					
OPq rA, rB	6	fn	rA	rB						
jCC Dest	7	cc	Dest							
call Dest	8	0	Dest							
ret	9	0								
pushq rA	A	0	rA	F						
popq rA	B	0	rA	F						

Y86-64 decoding

```
20 10 60 20 61 37 72 84 00 00 00 00 00 00 00
20 12 20 01 70 68 00 00 00 00 00 00 00
```

exercise: types of first three instructions?

byte:	0	1	2	3	4	5	6	7	8	9
halt	0	0								
nop	1	0								
rrmovq/cmovCC rA, rB	2	cc	rA	rB						
irmovq V, rB	3	0	F	rB	V					
rmmovq rA, D(rB)	4	0	rA	rB	D					
mrmovq D(rB), rA	5	0	rA	rB	D					
OPq rA, rB	6	fn	rA	rB						
jCC Dest	7	cc	Dest							
call Dest	8	0	Dest							
ret	9	0								
pushq rA	A	0	rA	F						
popq rA	B	0	rA	F						

Y86-64 decoding

20 10 60 20 61 37 72 84 00 00 00 00 00 00 00
20 12 20 01 70 68 00 00 00 00 00 00 00

byte:	0	1	2	3	4	5	6	7	8	9
halt	0	0								
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rrmovq/cmovCC rA, rB	2	cc	rA	rB						
irmovq V, rB	3	0	F	rB	V					
rmmovq rA, D(rB)	4	0	rA	rB	D					
mrmovq D(rB), rA	5	0	rA	rB	D					
OPq rA, rB	6	fn	rA	rB						
jCC Dest	7	cc	Dest							
call Dest	8	0	Dest							
ret	9	0								
pushq rA	A	0	rA	F						
popq rA	B	0	rA	F						

Y86-64 decoding

20 10 60 20 61 37 72 84 00 00 00 00 00 00 00
20 12 20 01 70 68 00 00 00 00 00 00 00

rrmovq %rcx, %rax

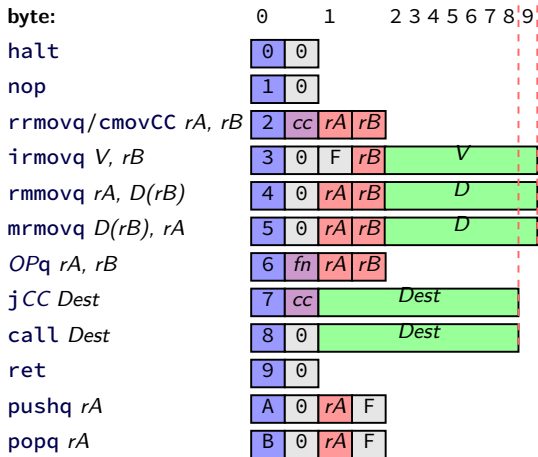
- ▶ 0 as cc: always
- ▶ 1 as reg: %rcx
- ▶ 0 as reg: %rax

byte:	0	1	2	3	4	5	6	7	8	9
halt	0	0								
nop	1	0								
rrmovq/cmovCC rA, rB	2	cc	rA	rB						
irmovq V, rB	3	0	F	rB	V					
rmmovq rA, D(rB)	4	0	rA	rB	D					
mrmovq D(rB), rA	5	0	rA	rB	D					
OPq rA, rB	6	fn	rA	rB						
jCC Dest	7	cc	Dest							
call Dest	8	0	Dest							
ret	9	0								
pushq rA	A	0	rA	F						
popq rA	B	0	rA	F						

Y86-64 decoding

```
20 10 60 20 61 37 72 84 00 00 00 00 00 00 00
20 12 20 01 70 68 00 00 00 00 00 00 00
```

```
rrmovq %rcx, %rax
addq   %rdx, %rax
subq   %rbx, %rdi
▶ 0 as fn: add
▶ 1 as fn: sub
```



Y86-64 decoding

```
20 10 60 20 61 37 72 84 00 00 00 00 00 00 00
20 12 20 01 70 68 00 00 00 00 00 00 00
```

```
rrmovq %rcx, %rax
addq   %rdx, %rax
subq   %rbx, %rdi
```

```
jl     0x84
```

- ▶ 2 as cc: ^l (less than)
- ▶ hex 8400... as little endian *Dest*:
0x84

byte:	0	1	2	3	4	5	6	7	8	9
halt	0	0								
nop	1	0								
rrmovq/cmovCC rA, rB	2	cc	rA	rB						
irmovq V, rB	3	0	F	rB	V					
rrmovq rA, D(rB)	4	0	rA	rB	D					
mrmovq D(rB), rA	5	0	rA	rB	D					
OPq rA, rB	6	fn	rA	rB						
jCC Dest	7	cc	Dest							
call Dest	8	0	Dest							
ret	9	0								
pushq rA	A	0	rA	F						
popq rA	B	0	rA	F						

Y86-64 decoding

20 10 60 20 61 37 72 84 00 00 00 00 00 00 00
20 12 20 01 70 68 00 00 00 00 00 00 00

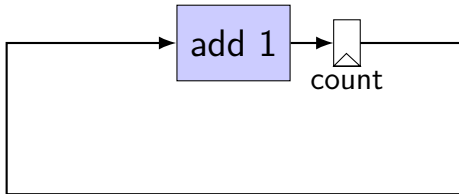
rrmovq %rcx, %rax
addq %rdx, %rax
subq %rbx, %rdi
jl 0x84
rrmovq %rcx, %rdx
rrmovq %rax, %rcx
jmp 0x68

byte:	0	1	2	3	4	5	6	7	8	9
halt	0	0								
nop	1	0								
rrmovq/cmovCC rA, rB	2	cc	rA	rB						
irmovq V, rB	3	0	F	rB	V					
rmmovq rA, D(rB)	4	0	rA	rB	D					
mrmovq D(rB), rA	5	0	rA	rB	D					
OPq rA, rB	6	fn	rA	rB						
jCC Dest	7	cc	Dest							
call Dest	8	0	Dest							
ret	9	0								
pushq rA	A	0	rA	F						
popq rA	B	0	rA	F						

describing hardware

how do we describe hardware?

pictures?



circuits with pictures?

yes, something you can do

such commercial tools exist, but...

not commonly used for processors

hardware description language

programming language for hardware

(typically) text-based representation of circuit

often abstracts away details like:

- how to build arithmetic operations from gates

- how to build registers from transistors

- how to build memories from transistors

- how to build MUXes from gates

- ...

those details also not a topic in this course

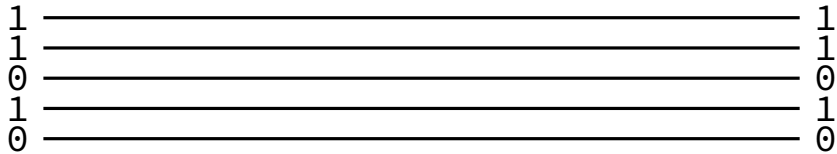
our tool: HCLRS

built for this course

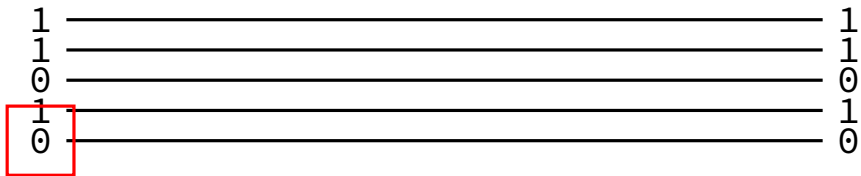
assumes you're making a processor

somewhat different from textbook's HCL

circuits: wires

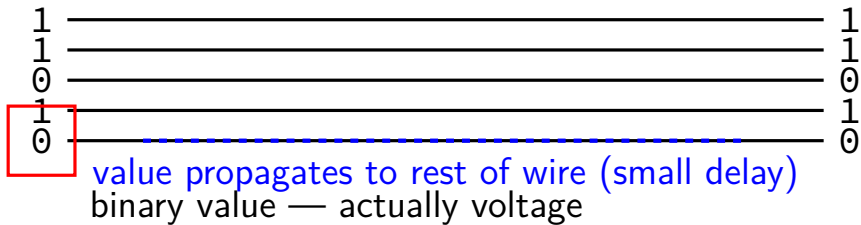


circuits: wires

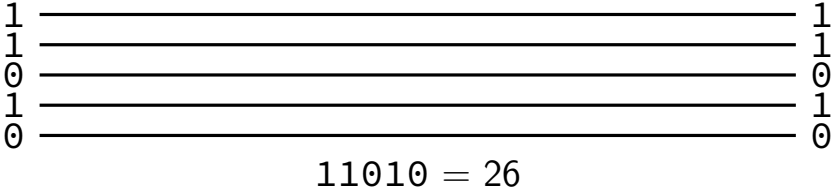


binary value — actually voltage

circuits: wires



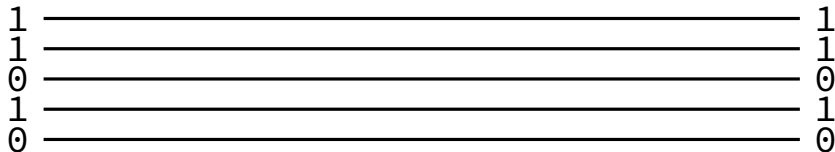
circuits: wire bundles



circuits: wire bundles



same as



$$11010 = 26$$

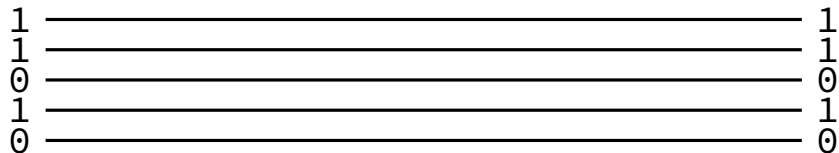
circuits: wire bundles



same as

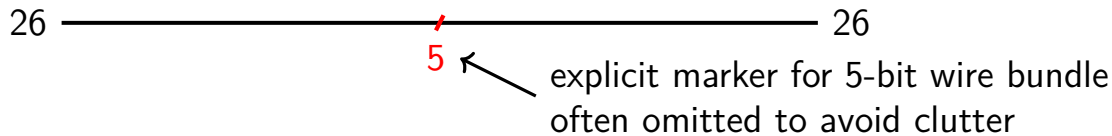


same as



$$11010 = 26$$

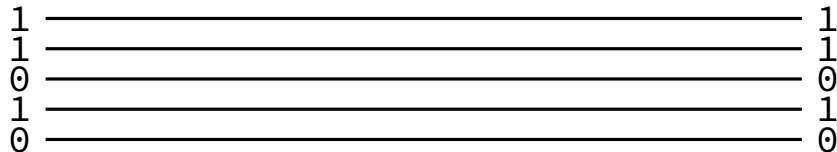
circuits: wire bundles



same as

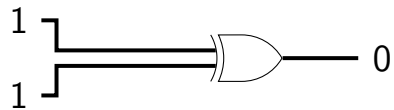
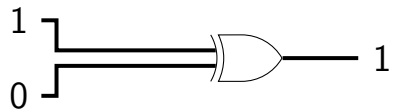
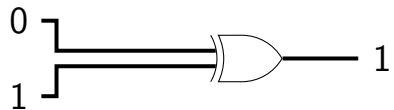
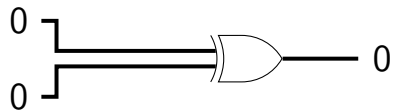


same as



$$11010 = 26$$

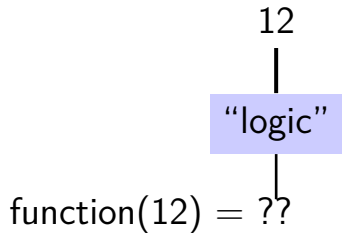
circuits: gates



circuits: logic

want to do calculations?

generalize gates:

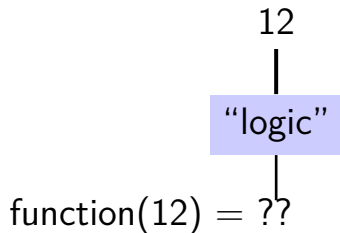


circuits: logic

want to do calculations?

generalize gates:

output wires contain result of function on input
changes as input changes (with delay)



circuits: logic

want to do calculations?

generalize gates:

output wires contain result of function on input
changes as input changes (with delay)

need not be same width as output

12

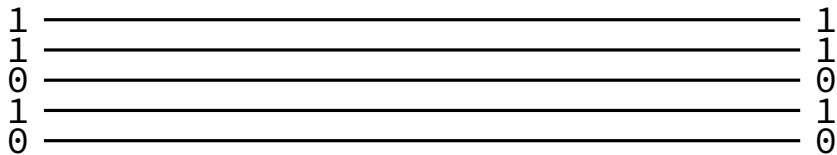


“logic”



function(12) = ??

HCLRS: wire (bundle)s

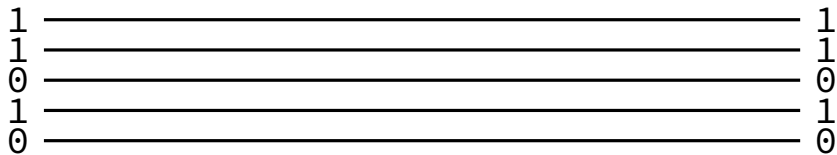


`wire foo : 5; foo = 0b11010;` *OR*

`wire foo : 5; foo = 26;` *OR*

`wire foo : 5; foo = 0x1a;`

HCLRS: wire (bundle)s



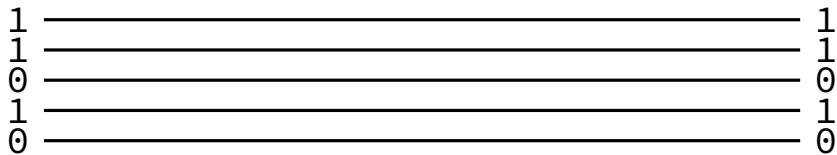
```
wire foo : 5; foo = 0b11010;    OR
```

```
wire foo : 5; foo = 26;        OR
```

```
wire foo : 5; foo = 0x1a;
```

name

HCLRS: wire (bundle)s



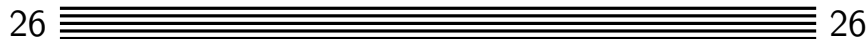
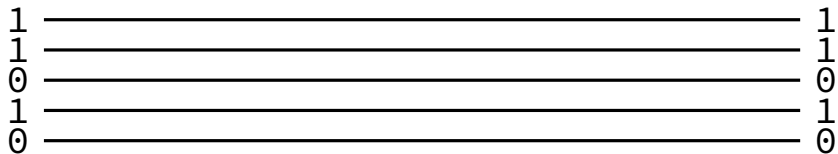
wire foo : 5; foo = 0b11010; OR

wire foo : 5; foo = 26; OR

wire foo : 5; foo = 0x1a;

width (in bits)

HCLRS: wire (bundle)s



wire foo : 5; foo = 0b11010; *OR*

wire foo : 5; foo = 26; *OR*

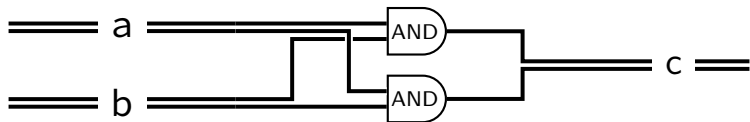
wire foo : 5; foo = 0x1a;

assignment

indicates wire is *connected* to value

HCLRS: gates + calculations (1)

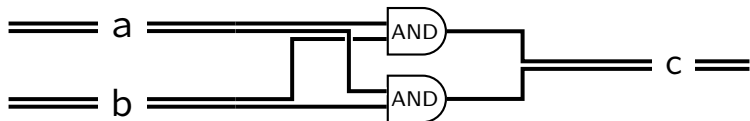
```
wire a : 2; wire b : 2; wire c : 2;  
c = b & a;  
a = 0b10;  
b = 0b11;
```



HCLRS: gates + calculations (1)

```
wire a : 2; wire b : 2; wire c : 2;  
c = b & a; }  
a = 0b10; } same as { a = 0b10;  
b = 0b11; }          { b = 0b11;  
                    { c = b & a;
```

order doesn't matter
connected or not



HCLRS: gates + calculations (1)

```
wire a : 2; wire b : 2; wire c : 2;
```

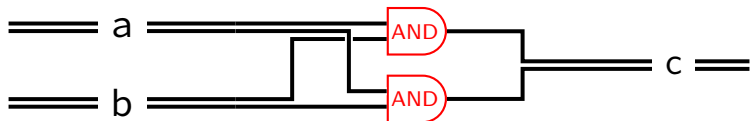
```
c = b & a;
```

```
a = 0b10;
```

```
b = 0b11;
```

C-like expressions supported

$0b10 \ \& \ 0b11 = 0b10$



HCLRS: gates + calculations (2)

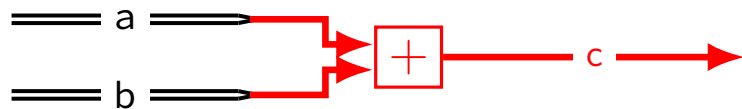
```
wire a : 2; wire b : 2; wire c : 2;
```

```
c = b + a; /* was b & a */
```

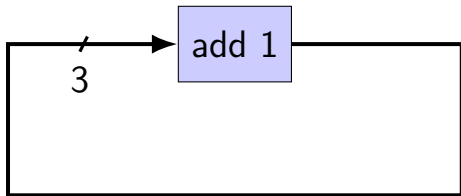
```
a = 0b10;
```

```
b = 0b11;
```

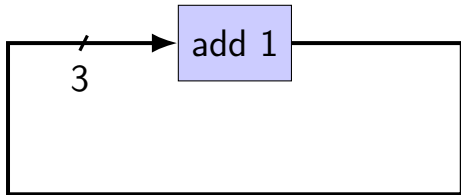
more than bitwise operators supported
 $0b10 + 0b11 = 0b101 \rightarrow 0b01$ (extra bits lost)



example: (broken) counter circuit (1)

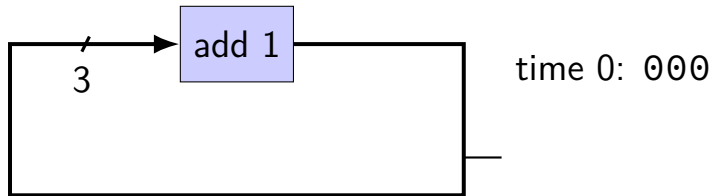


example: (broken) counter circuit (1)



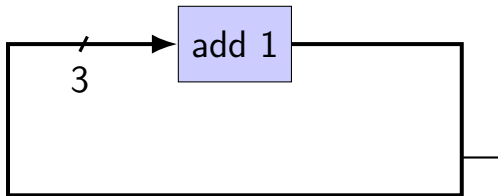
```
wire x : 3;  
x = x + 1;
```

example: (broken) counter circuit (1)



```
wire x : 3;  
x = x + 1;
```

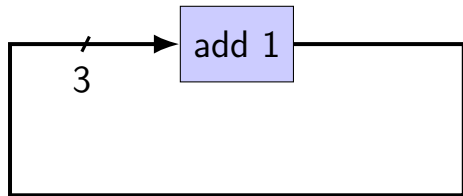
example: (broken) counter circuit (1)



time 0: 000 ← set how???

```
wire x : 3;  
x = x + 1;
```

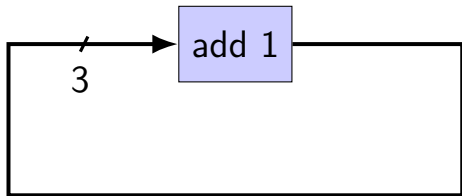
example: (broken) counter circuit (1)



time 0: 000
time 1: 001?
time 2: 010?
time 3: 011?

```
wire x : 3;  
x = x + 1;
```

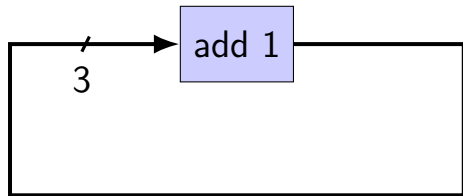

example: (broken) counter circuit (2)



~~wire x : 3;
x = x + 1;~~

HCLRS: compile error
"Circular dependency detected:
x depends on x"

example: (broken) counter circuit (3)

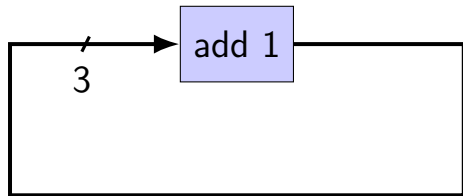


time 0: 000
time 1: 001?
time 2: 010?
time 3: 011?

```
wire x : 3;  
x = x + 1;
```



example: (broken) counter circuit (3)



time 0: 000

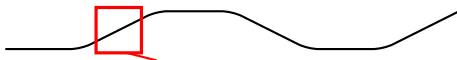
time 1: 001?

time 2: 010?

time 3: 011?

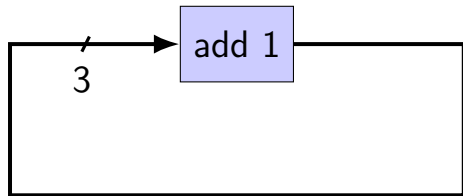
```
wire x : 3;
```

```
x = x + 1;
```



problem 1: how will "add 1" react to this value?
(not zero or one) ...

example: (broken) counter circuit (3)

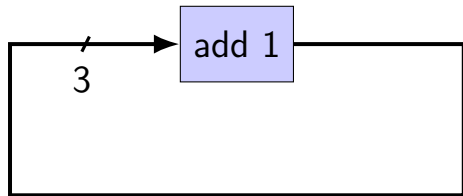


time 0: 000
time 1: 001?
time 2: 010?
time 3: 011?

```
wire x : 3;  
x = x + 1;
```

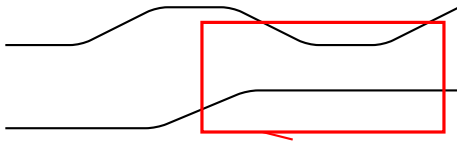


example: (broken) counter circuit (3)



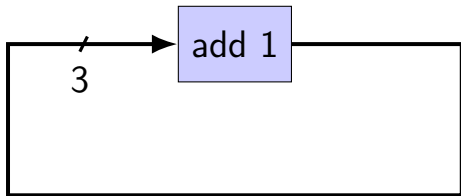
time 0: 000
time 1: 001?
time 2: 010?
time 3: 011?

```
wire x : 3;  
x = x + 1;
```



problem 2: changes not in sync?

example: (broken) counter circuit (4)



```
wire x : 3;  
x = x + 1;
```

~~time 0: 000
time 1: 001?
time 2: 010?
time 3: 011?~~

circuit is **not stable**
transient values during changes
hard to predict behavior

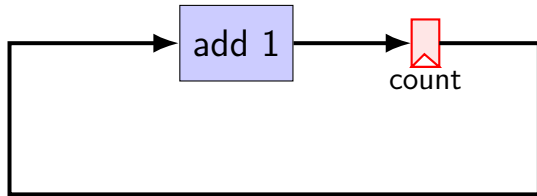
circuits: state

logic performs calculations all the time

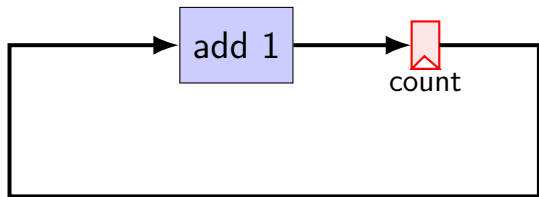
never stores values!

need **extra elements** to store values
registers, memory

example: counter circuit (corrected)



example: counter circuit (corrected)



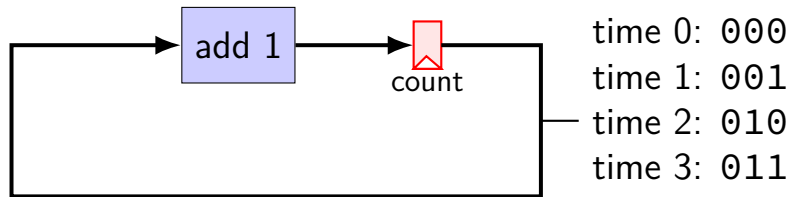
time 0: 000

time 1: 001

time 2: 010

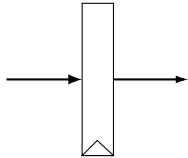
time 3: 011

example: counter circuit (corrected)

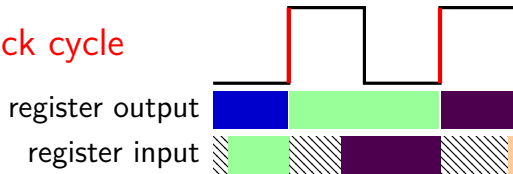


add **register** to store current count
updates based on “clock signal” (not shown)
avoids intermediate updates

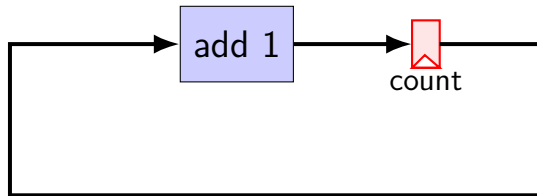
registers



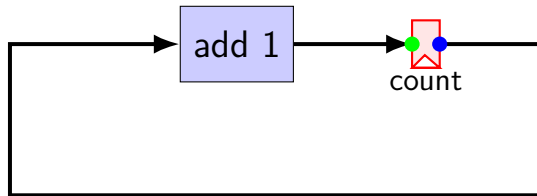
updates every **clock cycle**



example: counter circuit (real HCLRS)

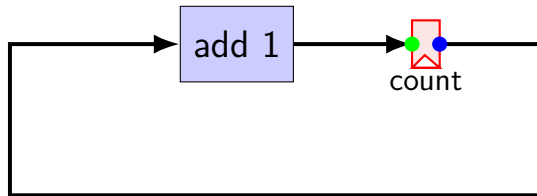


example: counter circuit (real HCLRS)



```
register xY {  
    count : 3 = 0b000 ;  
}  
x_count = Y_count + 0b001;
```

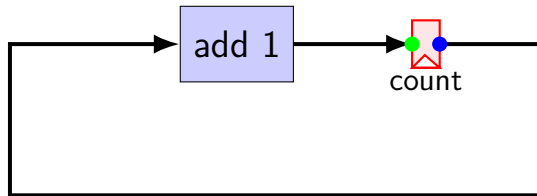
example: counter circuit (real HCLRS)



```
register xY {  
    count : 3 = 0b000 ;  
}  
x_count = Y_count + 0b001;
```

register “bank”
can have multiple (related) registers

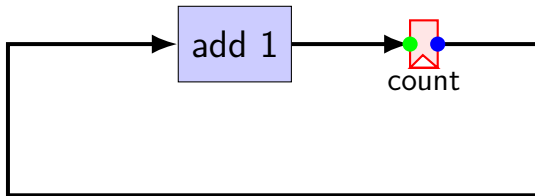
example: counter circuit (real HCLRS)



```
register xY {  
    count : 3 = 0b000 ;  
}  
x_count = Y_count + 0b001;
```

label for left/right side of registers
x: label for input side (always lowercase)
Y: label for output side (always uppercase)

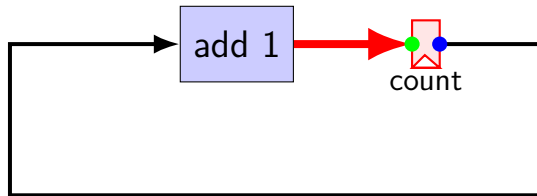
example: counter circuit (real HCLRS)



```
register xY {  
    count : 3 = 0b000 ;  
}  
x_count = Y_count + 0b001;
```

register "name"
input/output = *prefix_name*

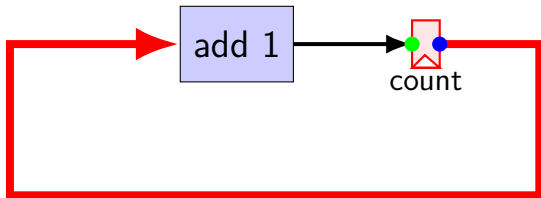
example: counter circuit (real HCLRS)



```
register xY {  
    count : 3 = 0b000 ;  
}  
x_count = Y_count + 0b001;
```

input wire to register

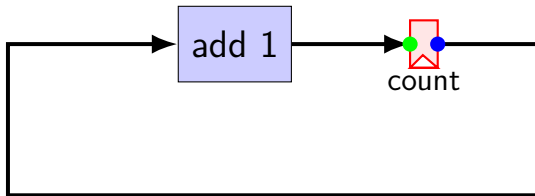
example: counter circuit (real HCLRS)



```
register xY {  
    count : 3 = 0b000 ;  
}  
x_count = Y_count + 0b001;
```

output wire of register

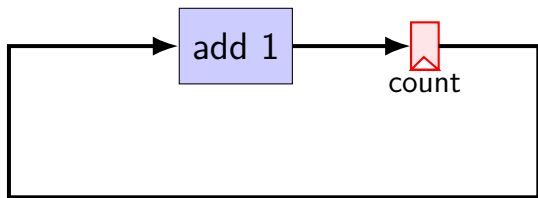
example: counter circuit (real HCLRS)



initial value of register
first value for output wire (Y_count)

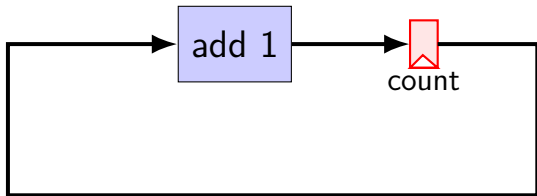
```
register xY {  
    count : 3 = 0b000;  
}  
x_count = Y_count + 0b001;
```

example: counter circuit



```
register xY {  
    count : 3 = 0b000 ;  
}  
x_count = Y_count + 0b001;
```

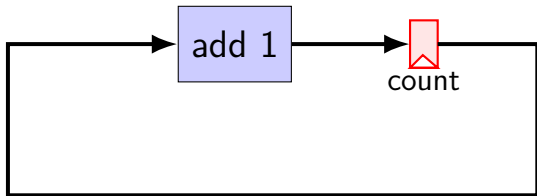
example: counter circuit



```
register xY {  
    count : 3 = 0b000 ;  
}  
x_count = Y_count + 0b001;
```

time	Y_count	x_count
start	000	001
start + 1 rising edge	001	010
start + 2 rising edges	010	011
start + 3 rising edges	011	100
...

example: counter circuit



```
register xY {  
    count : 3 = 0b000 ;  
}  
x_count = Y_count + 0b001;
```

time	Y_count	x_count
start	000	001
start + 1 rising edge	001	010
start + 2 rising edges	010	011
start + 3 rising edges	011	100
...

HCL circuit with registers

```
register xY {  
    a : 4 = 1; /* <-- initial Y_a */  
    b : 4 = 1; /* <-- initial Y_b */  
}
```

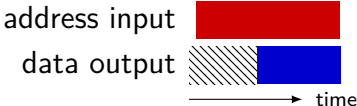
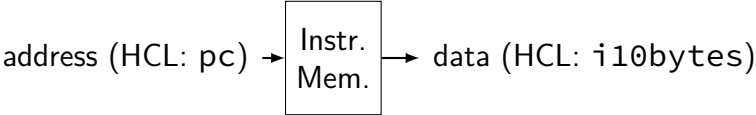
$x_b = x_a + Y_a;$

$x_a = Y_a + Y_b;$

exercise: value of Y_a , Y_b after two rising edges of clock?

- A. $Y_a = 2, Y_b = 3$
- B. $Y_a = 2, Y_b = 2$
- C. $Y_a = 3, Y_b = 5$
- D. $Y_a = 3, Y_b = 7$
- E. $Y_a = 3, Y_b = 11$
- F. $Y_a = 5, Y_b = 7$
- G. $Y_a = 7, Y_b = 11$
- H. none of the above

instruction memory



Stat signal

how do we stop the simulated machine?

hard-wired mechanism — Stat wire

possible values:

STAT_AOK — keep going

STAT_HLT — stop, normal shutdown

STAT_INS — invalid instruction

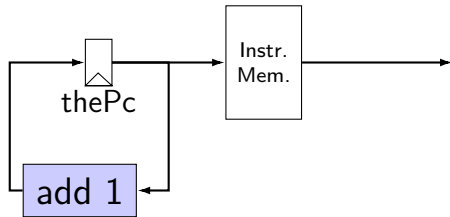
...(and more errors)

(predefined 3-bit constants)

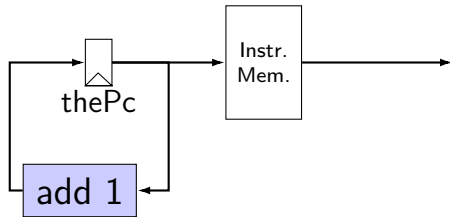
must be set

determines if **simulator** keeps going

nop CPU

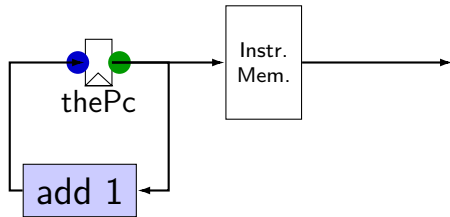


nop CPU



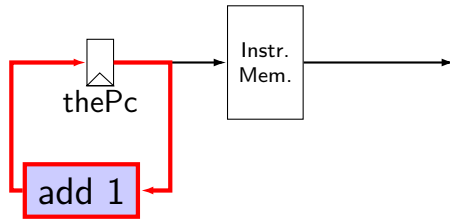
```
register pF {  
    thePc : 64 = 0;  
}
```

nop CPU



```
register pF {  
    thePc : 64 = 0;  
}
```

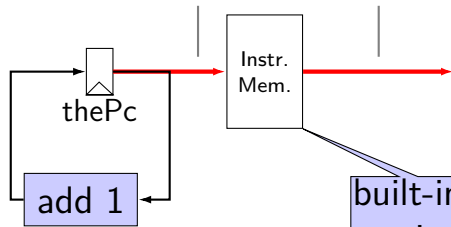
nop CPU



```
register pF {  
    thePc : 64 = 0;  
}  
p_thePc = F_thePc + 1;
```

nop CPU

“pc” “i10bytes”

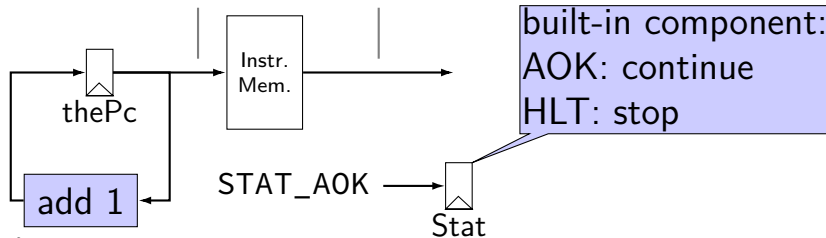


built-in component
use is mandatory

```
register pF {  
    thePc : 64 = 0;  
}  
p_thePc = F_thePc + 1;  
pc = F_thePc;
```

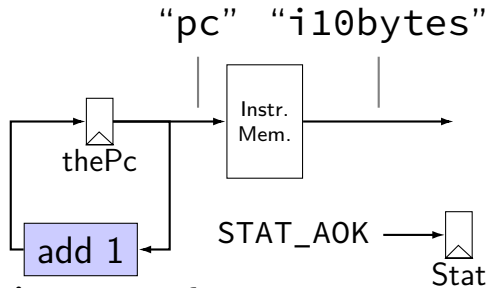
nop CPU

“pc” “i10bytes”



```
register pF {  
    thePc : 64 = 0;  
}  
p_thePc = F_thePc + 1;  
pc = F_thePc;  
Stat = STAT_AOK;
```

nop CPU



```
register pF {  
    thePc : 64 = 0;  
}  
p_thePc = F_thePc + 1;  
pc = F_thePc;  
Stat = STAT_AOK;
```


nop CPU: running

need a program in memory

.yo file

tools/yas — convert .ys to .yo

tools/yis — reference interpreter for .yo files

if your processor doesn't do the same thing...

can build tools by running make

nop CPU: creating a program

create assembly file: nops.ya:

```
nop  
nop  
nop  
nop  
nop
```

assemble using `tools/yas nops.ya` or `make nops.yo`

nop.yo

more readable/simpler than normal executables:

0x000: 10		nop
0x001: 10		nop
0x002: 10		nop
0x003: 10		nop
0x004: 10		nop

loaded into data and program memory

parts left of | just comments

running a simulator (1)

Usage: ./hclrs [options] HCL-FILE [YO-FILE [TIMEOUT]]

Runs HCL_FILE on YO-FILE. If --check is specified, no YO-FILE may be supplied.
Default timeout is 9999 cycles.

Options:

-c, --check	check syntax only
-d, --debug	output wire values after each cycle and other debug output
-q, --quiet	only output state at the end
-t, --testing	do not output custom register banks (for autograding)
-h, --help	print this help menu
-i, --interactive	prompt after each cycle
--trace-assignments	show assignments in the order they are simulated
--version	print version number

running a simulator (2)

```
$ ./hclrs nop_cpu.hcl nops.yo
```

```
+----- between cycles      0 and      1 -----+
| RAX:                0   RCX:                0   RDX:                0   |
| RBX:                0   RSP:                0   RBP:                0   |
| RSI:                0   RDI:                0   R8:                 0   |
| R9:                 0   R10:               0   R11:               0   |
| R12:                0   R13:               0   R14:               0   |
| register pF(N)      thePc=000000000000000000 |
| used memory:       _0 _1 _2 _3 _4 _5 _6 _7   _8 _9 _a _b _c _d _e _f |
| 0x00000000_:      10 10 10 10 10             |
+-----+

```

```
pc = 0x0; loaded [10 : nop]
```

```
....
```

```
+----- timed out after 9999 cycles in state: -----+
| RAX:                0   RCX:                0   RDX:                0   |
| RBX:                0   RSP:                0   RBP:                0   |
| RSI:                0   RDI:                0   R8:                 0   |
| R9:                 0   R10:               0   R11:               0   |
| R12:                0   R13:               0   R14:               0   |
| register pF(N)      thePc=0000000000000270f |
| used memory:       _0 _1 _2 _3 _4 _5 _6 _7   _8 _9 _a _b _c _d _e _f |
| 0x00000000_:      10 10 10 10 10             |
+-----+

```

running a simulator (2)

```
$ ./hclrs nop_cpu.hcl nops.yo
```

```
+----- between cycles      0 and      1 -----+
| RAX:                0    RCX:                0    RDX:                0    |
| RBX:                0    RSP:                0    RBP:                0    |
| RSI:                0    RDI:                0    R8:                 0    |
| R9:                 0    R10:               0    R11:               0    |
| R12:               0    R13:               0    R14:               0    |
| register pF(N)      thePc=0000000000000000 |
| used memory:      _0 _1 _2 _3 _4 _5 _6 _7  _8 _9 _a _b _c _d _e _f |
| 0x00000000_:    10 10 10 10 10 |
+-----+
```

```
pc = 0x0; loaded [10 : nop]
```

```
....
```

```
+----- timed out after 9999 cycles in state: -----+
| RAX:                0    RCX:                0    RDX:                0    |
| RBX:                0    RSP:                0    RBP:                0    |
| RSI:                0    RDI:                0    R8:                 0    |
| R9:                 0    R10:               0    R11:               0    |
| R12:               0    R13:               0    R14:               0    |
| register pF(N)      thePc=0000000000000270f |
| used memory:      _0 _1 _2 _3 _4 _5 _6 _7  _8 _9 _a _b _c _d _e _f |
| 0x00000000_:    10 10 10 10 10 |
+-----+
```

running a simulator (2)

```
$ ./hclrs nop_cpu.hcl nops.yo
```

```
+----- between cycles      0 and      1 -----+
| RAX:          0   RCX:          0   RDX:          0   |
| RBX:          0   RSP:          0   RBP:          0   |
| RSI:          0   RDI:          0   R8:           0   |
| R9:           0   R10:         0   R11:          0   |
| R12:          0   R13:         0   R14:          0   |
| register pF(N)  thePc=000000000000000000          |
| used memory:   _0 _1 _2 _3 _4 _5 _6 _7   _8 _9 _a _b _c _d _e _f |
| 0x00000000_:  10 10 10 10 10          |
+-----+

```

```
pc = 0x0; loaded [10 : nop]
```

```
....
```

```
+----- timed out after 9999 cycles in state: -----+
| RAX:          0   RCX:          0   RDX:          0   |
| RBX:          0   RSP:          0   RBP:          0   |
| RSI:          0   RDI:          0   R8:           0   |
| R9:           0   R10:         0   R11:          0   |
| R12:          0   R13:         0   R14:          0   |
| register pF(N)  thePc=0000000000000270f          |
| used memory:   _0 _1 _2 _3 _4 _5 _6 _7   _8 _9 _a _b _c _d _e _f |
| 0x00000000_:  10 10 10 10 10          |
+-----+

```

running a simulator (2)

```
$ ./hclrs nop_cpu.hcl nops.yo
```

```
+----- between cycles      0 and      1 -----+
| RAX:                0    RCX:                0    RDX:                0    |
| RBX:                0    RSP:                0    RBP:                0    |
| RSI:                0    RDI:                0    R8:                0    |
| R9:                 0    R10:               0    R11:               0    |
| R12:               0    R13:               0    R14:               0    |
| register pF(N)      thePc=000000000000000000 |
| used memory:      _0 _1 _2 _3 _4 _5 _6 _7   _8 _9 _a _b _c _d _e _f |
| 0x00000000_:     10 10 10 10 10              |
+-----+

```

```
pc = 0x0; loaded [10 : nop]
```

```
....
```

```
+----- timed out after 9999 cycles in state: -----+
| RAX:                0    RCX:                0    RDX:                0    |
| RBX:                0    RSP:                0    RBP:                0    |
| RSI:                0    RDI:                0    R8:                0    |
| R9:                 0    R10:               0    R11:               0    |
| R12:               0    R13:               0    R14:               0    |
| register pF(N)      thePc=0000000000000270f |
| used memory:      _0 _1 _2 _3 _4 _5 _6 _7   _8 _9 _a _b _c _d _e _f |
| 0x00000000_:     10 10 10 10 10              |
+-----+

```


backup slides

Y86-64 instruction formats

byte:	0	1	2	3	4	5	6	7	8	9
halt	0	0								
nop	1	0								
rrmovq/cmovCC <i>rA</i> , <i>rB</i>	2	cc	<i>rA</i>	<i>rB</i>						
irmovq <i>V</i> , <i>rB</i>	3	0	F	<i>rB</i>	<i>V</i>					
rmmovq <i>rA</i> , <i>D(rB)</i>	4	0	<i>rA</i>	<i>rB</i>	<i>D</i>					
mrmovq <i>D(rB)</i> , <i>rA</i>	5	0	<i>rA</i>	<i>rB</i>	<i>D</i>					
OPq <i>rA</i> , <i>rB</i>	6	fn	<i>rA</i>	<i>rB</i>						
jCC <i>Dest</i>	7	cc	<i>Dest</i>							
call <i>Dest</i>	8	0	<i>Dest</i>							
ret	9	0								
pushq <i>rA</i>	A	0	<i>rA</i>	F						
popq <i>rA</i>	B	0	<i>rA</i>	F						

Secondary opcodes: `cmovcc/jcc`

byte:	0	1	2	3	4	5	6	7	8	9
halt	0	0								
nop	1	0								
<code>rrmovq/cmovCC rA, rB</code>	2	cc	rA	rB						
<code>irmovq V, rB</code>	3	0	F	rB						
<code>rmmovq rA, D(rB)</code>	4	0	rA	rB						
<code>mrmovq D(rB), rA</code>	5	0	rA	rB						
<code>OPq rA, rB</code>	6	fn	rA	rB						
<code>jCC Dest</code>	7	cc								
<code>call Dest</code>	8	0								
ret	9	0								
<code>pushq rA</code>	A	0	rA	F						
<code>popq rA</code>	B	0	rA	F						

0	<i>always</i> (<code>jmp/rrmovq</code>)
1	<code>le</code>
2	<code>l</code>
3	<code>e</code>
4	<code>ne</code>
5	<code>ge</code>
6	<code>g</code>

Secondary opcodes: OPq

byte:	0	1	2	3	4	5	6	7	8	9
halt	0	0								
nop	1	0								
rmmovq/cmovCC rA, rB	2	cc	rA	rB						
irmovq V, rB	3	0	F	rB			V			
rmmovq $rA, D(rB)$	4	0	rA	rB	0	add	D			
mrmovq $D(rB), rA$	5	0	rA	rB	1	sub	D			
OPq rA, rB	6	fn	rA	rB	2	and				
jCC $Dest$	7	cc			3	xor				
call $Dest$	8	0					$Dest$			
ret	9	0								
pushq rA	A	0	rA	F						
popq rA	B	0	rA	F						

Registers: rA , rB

byte:	0	1	2
halt	0	0	
nop	1	0	
rmmovq/cmovCC rA , rB	2	cc	rA rB
irmovq V , rB	3	0	F rB
rmmovq rA , $D(rB)$	4	0	rA rB
mrmmovq $D(rB)$, rA	5	0	rA rB
OPq rA , rB	6	ff	rA rB
jCC Dest	7	cc	
call Dest	8	0	
ret	9	0	
pushq rA	A	0	rA F
popq rA	B	0	rA F

0	%rax	8	%r8
1	%rcx	9	%r9
2	%rdx	A	%r10
3	%rbx	B	%r11
4	%rsp	C	%r12
5	%rbp	D	%r13
6	%rsi	E	%r14
7	%rdi	F	none

Immediates: *V, D, Dest*

byte:	0	1	2	3	4	5	6	7	8	9
halt	0	0								
nop	1	0								
rrmovq/cmovCC <i>rA, rB</i>	2	cc	<i>rA</i>	<i>rB</i>						
irmovq <i>V, rB</i>	3	0	F	<i>rB</i>	<i>V</i>					
rmmovq <i>rA, D(rB)</i>	4	0	<i>rA</i>	<i>rB</i>	<i>D</i>					
mrmovq <i>D(rB), rA</i>	5	0	<i>rA</i>	<i>rB</i>	<i>D</i>					
<i>OPq rA, rB</i>	6	<i>fn</i>	<i>rA</i>	<i>rB</i>						
<i>jCC Dest</i>	7	cc	<i>Dest</i>							
call <i>Dest</i>	8	0	<i>Dest</i>							
ret	9	0								
pushq <i>rA</i>	A	0	<i>rA</i>	F						
popq <i>rA</i>	B	0	<i>rA</i>	F						

Immediates: *V, D, Dest*

byte:	0	1	2	3	4	5	6	7	8	9
halt	0	0								
nop	1	0								
rrmovq/cmovCC <i>rA, rB</i>	2	cc	rA	rB						
irmovq <i>V, rB</i>	3	0	F	rB	<i>V</i>					
rmmovq <i>rA, D(rB)</i>	4	0	rA	rB	<i>D</i>					
mrmovq <i>D(rB), rA</i>	5	0	rA	rB	<i>D</i>					
OPq <i>rA, rB</i>	6	fn	rA	rB						
jCC <i>Dest</i>	7	cc	<i>Dest</i>							
call <i>Dest</i>	8	0	<i>Dest</i>							
ret	9	0								
pushq <i>rA</i>	A	0	rA	F						
popq <i>rA</i>	B	0	rA	F						

differences from book

wire not **bool** or **int**

book uses names like `val C` — not required!

author's environment limited adding new wires

MUXes must have default (`1 : something`) case

implement your own ALU

differences from book

wire not **bool** or **int**

book uses names like `valC` — not required!

author's environment limited adding new wires

MUXes must have default (`1 : something`) case

implement your own ALU

differences from book

wire not **bool** or **int**

book uses names like `val C` — not required!

author's environment limited adding new wires

MUXes must have default (`1 : something`) case

implement your own ALU

comparing to yis

```
$ ./hclrs nopjmp_cpu.hcl nopjmp.yo
```

```
...  
...
```

```
+----- (end of halted state) -----+
```

```
Cycles run: 7
```

```
$ ./tools/yis nopjmp.yo
```

```
Stopped in 7 steps at PC = 0x1e.  Status 'HLT', CC Z=1 S=0 O=0
```

```
Changes to registers:
```

```
Changes to memory:
```

HCLRS summary

declare/assign values to **wires**

MUXes with

```
[ test1: value1; test2: value2; 1: default; ]
```

register banks with **register** `i0`:

next value on `i_name`; current value on `O_name`

fixed functionality

register file (15 registers; 2 read + 2 write)

memories (data + instruction)

Stat register (start/stop/error)