

compilation pipeline / C

changelog

Changes since first lecture:

2 September 2021: what's in these files? make indexing consistently 0-based

6 September 2021: edit what's in the files slides to clarify why .Lstr isn't in symbol table; add explicit .global declaration to main

last time

Linux x86-64 calling convention

agreement between function writers/users (not set by hardware)
certain args for registers/return value; overflow on stack
callee-saved = functions must restore orig value before returning
caller-saved = must handle function call changing value

condition codes

result of last arithmetic (ZF = was zero? SF = was negative?)
cmp = sub without storing; test = and without storing
jXX instructions named after compare result to 0

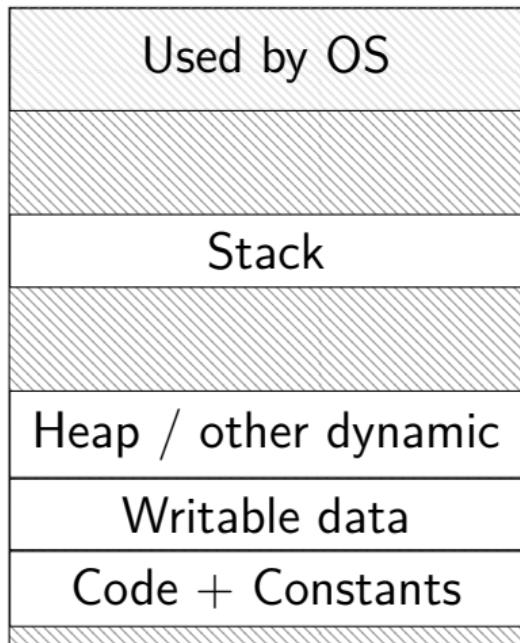
mapping C to assembly

one strategy: convert to C with gotos, then to assembly

compiling switches

series of if/then OR binary search OR jump table

program memory (x86-64 Linux)



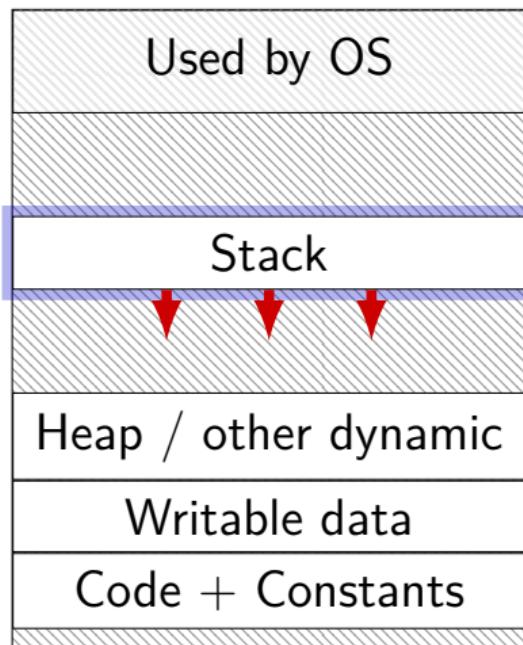
0xFFFF FFFF FFFF FFFF

0xFFFF 8000 0000 0000

0x7F...

0x0000 0000 0040 0000

program memory (x86-64 Linux)



0xFFFF FFFF FFFF FFFF

0xFFFF 8000 0000 0000

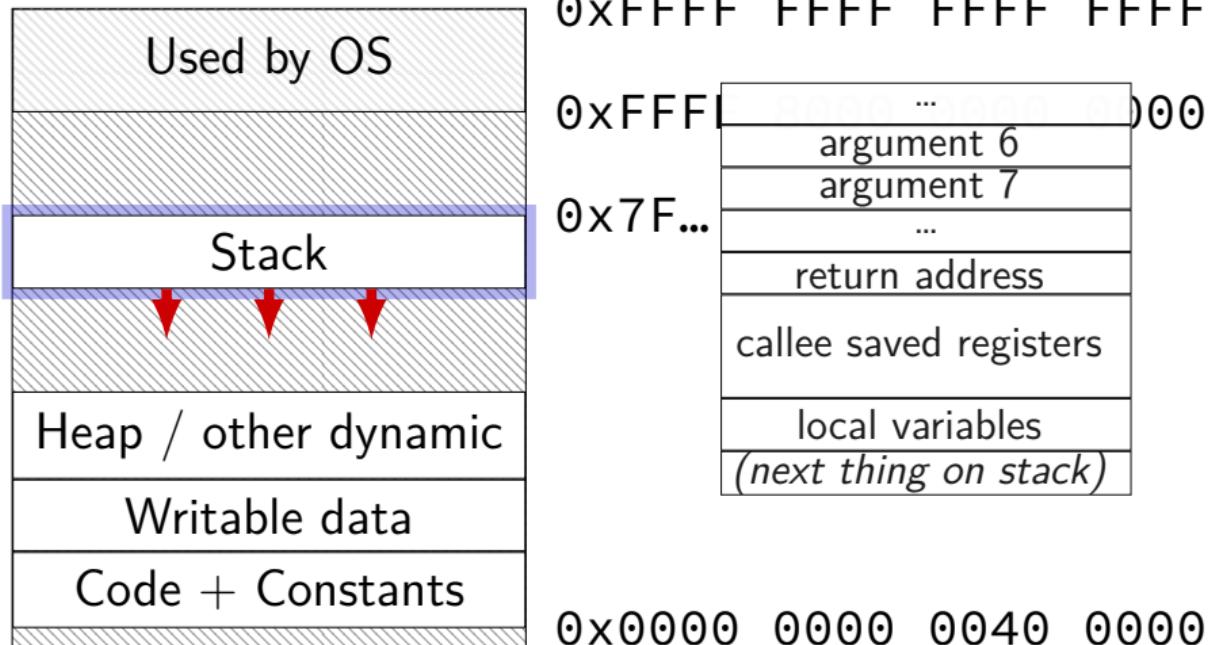
0x7F...

stack *grows down*

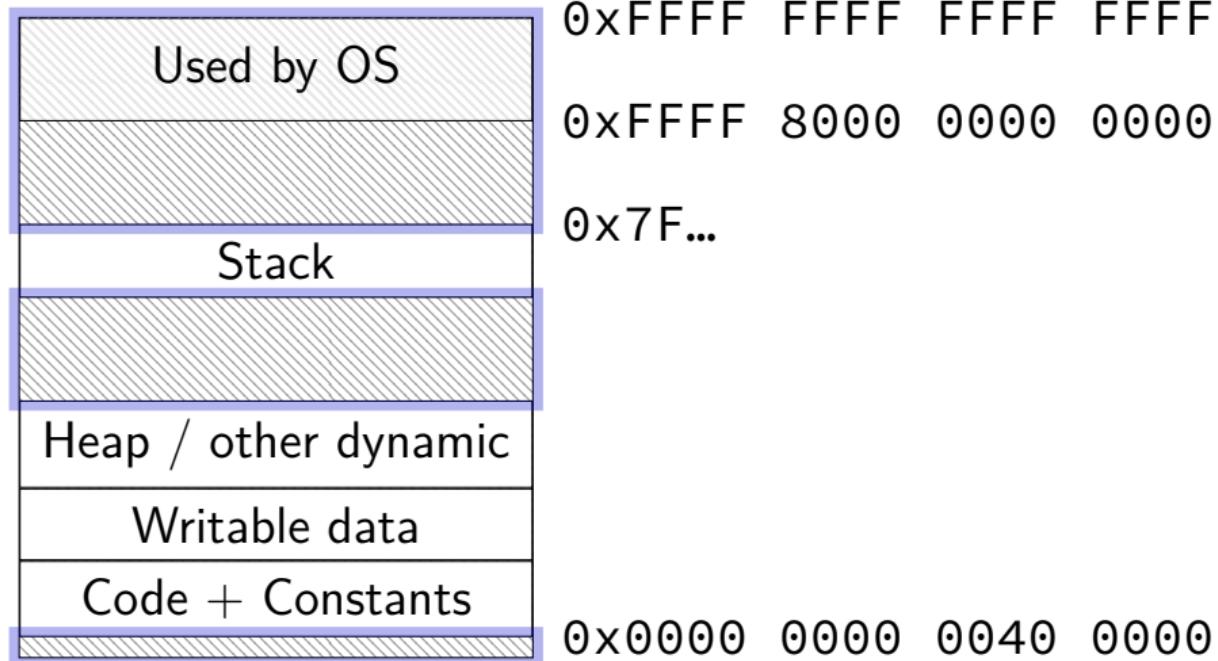
“top” has smallest address

0x0000 0000 0040 0000

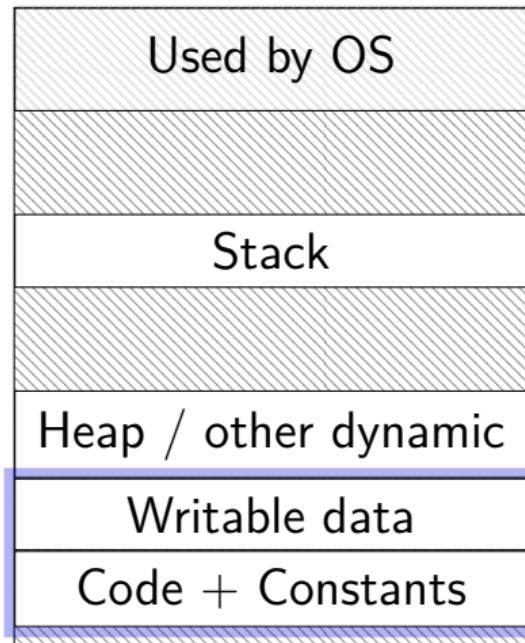
program memory (x86-64 Linux)



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0xFFFF FFFF FFFF FFFF

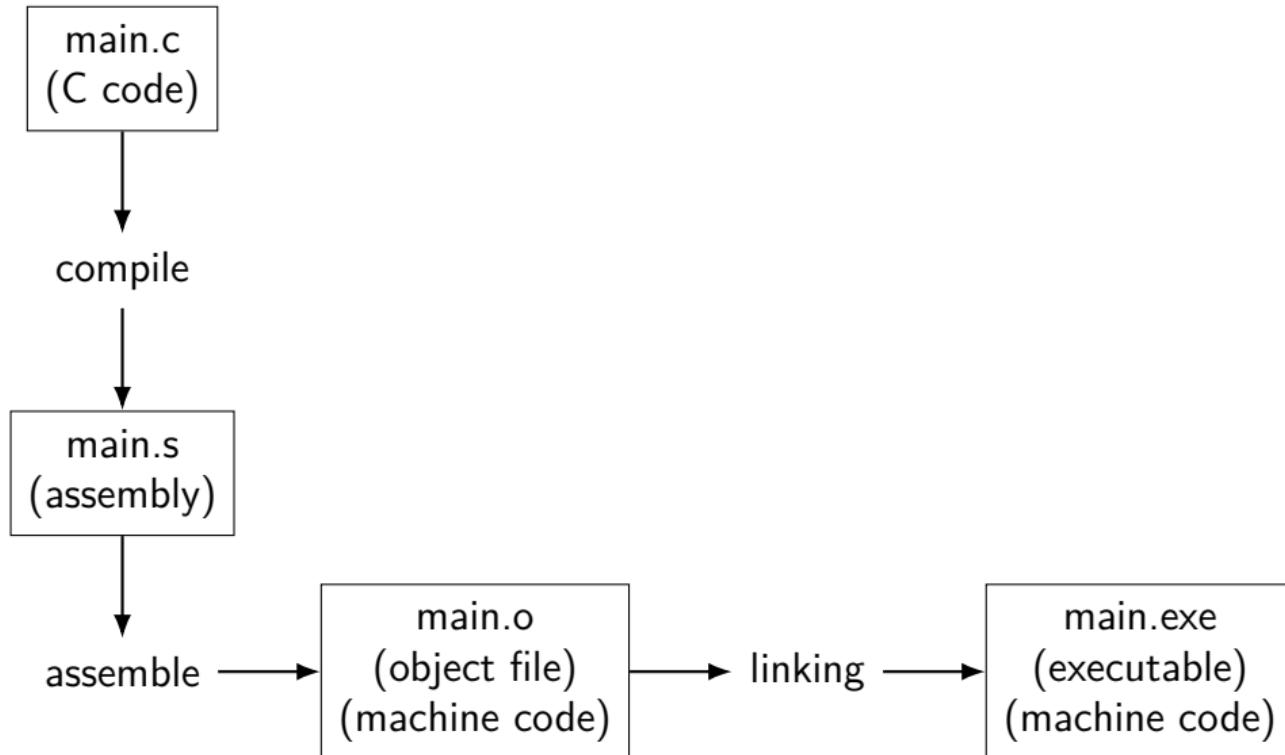
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0x7F...

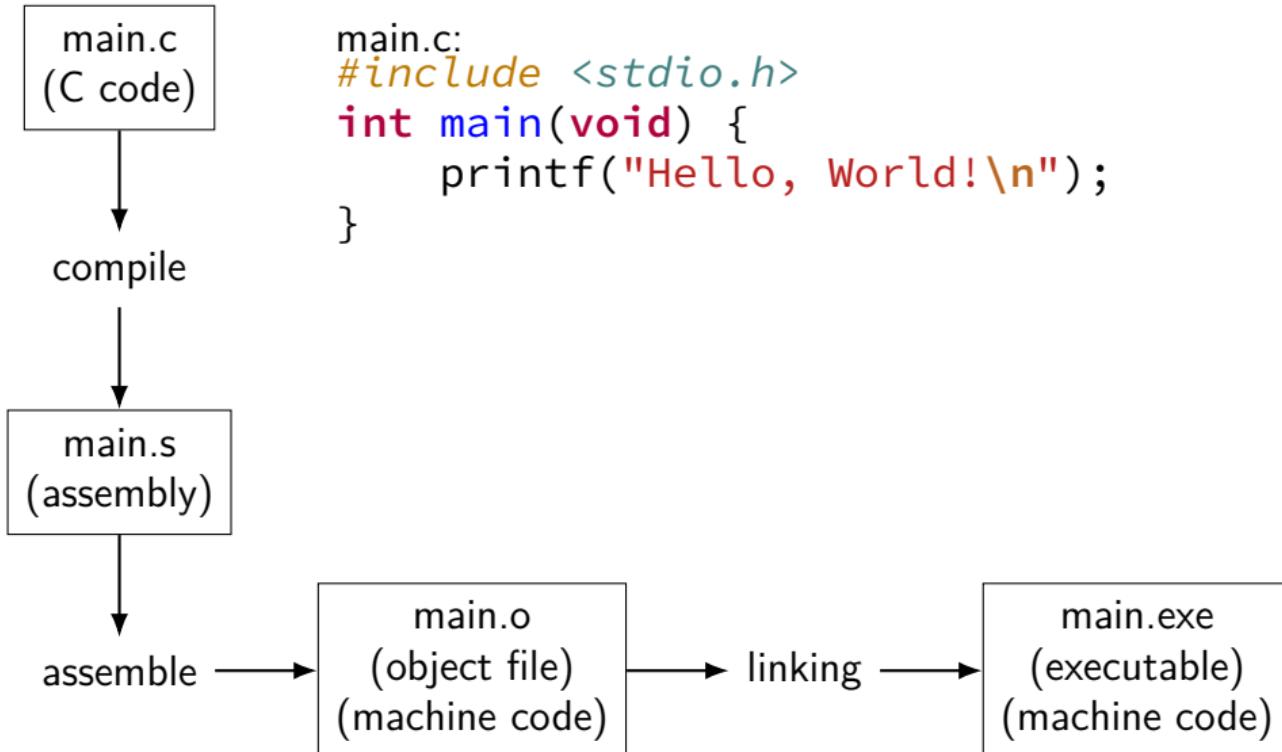
loaded from executable file

0x0000 0000 0040 0000

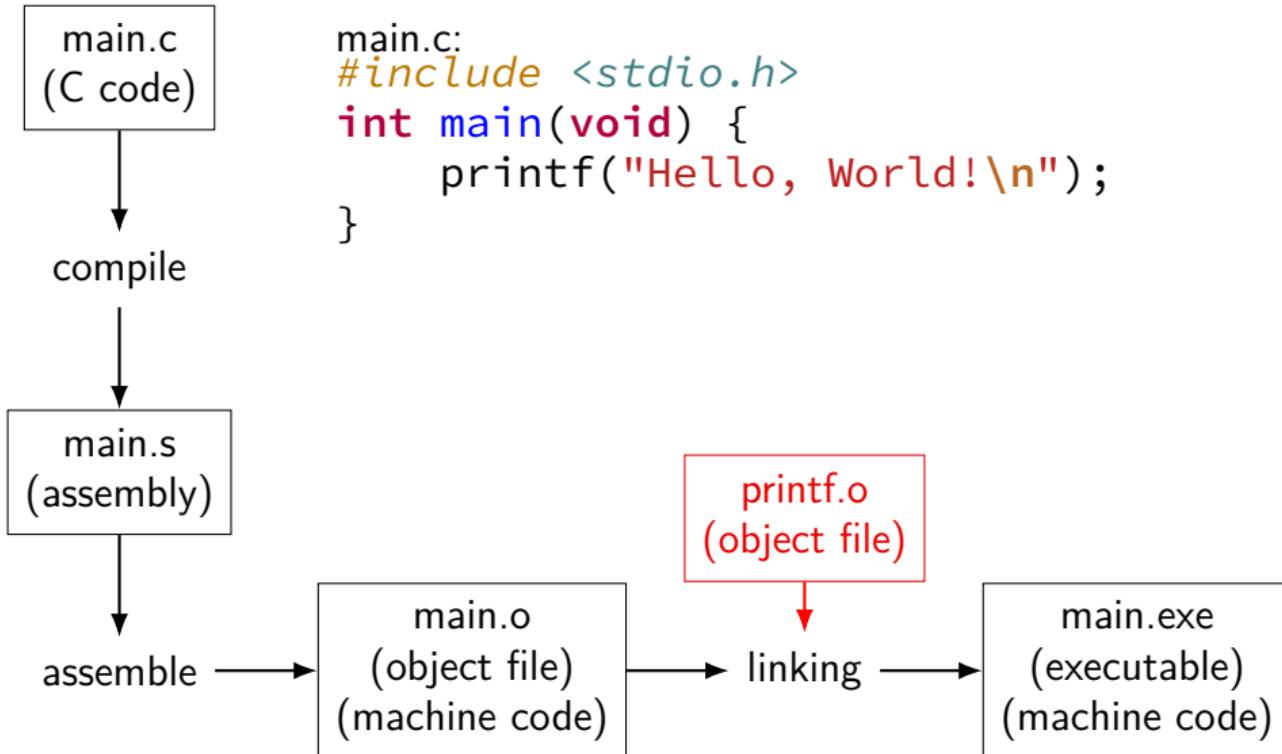
compilation pipeline



compilation pipeline



compilation pipeline



compilation commands

compile: gcc -S file.c ⇒ file.s (assembly)

assemble: gcc -c file.s ⇒ file.o (object file)

link: gcc -o file file.o ⇒ file (executable)

c+a: gcc -c file.c ⇒ file.o

c+a+l: gcc -o file file.c ⇒ file

...

what's in those files?

hello.c

```
#include <stdio.h>
int main(void) {
    puts("Hello, World!");
    return 0;
}
```

what's in those files?

hello.c

```
#include <stdio.h>
int main(void) {
    puts("Hello, World!");
    return 0;
}
```

hello.s

```
.text
.global main
main:
    sub    $8, %rsp
    mov    $.Lstr, %rdi
    call   puts
    xor    %eax, %eax
    add    $8, %rsp
    ret

.data
.Lstr: .string "Hello, World!"
```

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    ret

.data
.Lstr: .string "Hello, World!"
```

hello.s (Intel syntax)

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#include <stdio.h>
int main(void) {
    puts("Hello, World!");
    return 0;
}
```

hello.s

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.text
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main:
    sub $8, %rsp
    mov $.Lstr, %rdi
    call puts
    xor %eax, %eax
    add $8, %rsp
    ret

.data
.Lstr: .string "Hello, World!"
```

Linux x86-64
calling convention:
stack addr. must be
multiple of 16

what's in those files?

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int main(void) {
    puts("Hello, World!");
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}
```

hello.s

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    xor %eax, %eax ←
    add $8, %rsp
    ret

.data
.Lstr: .string "Hello, World!"
```

sets eax to 0
(shorter machine code than mov)

what's in those files?

hello.c

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#include <stdio.h>
int main(void) {
    puts("Hello, World!");
    return 0;
}
```

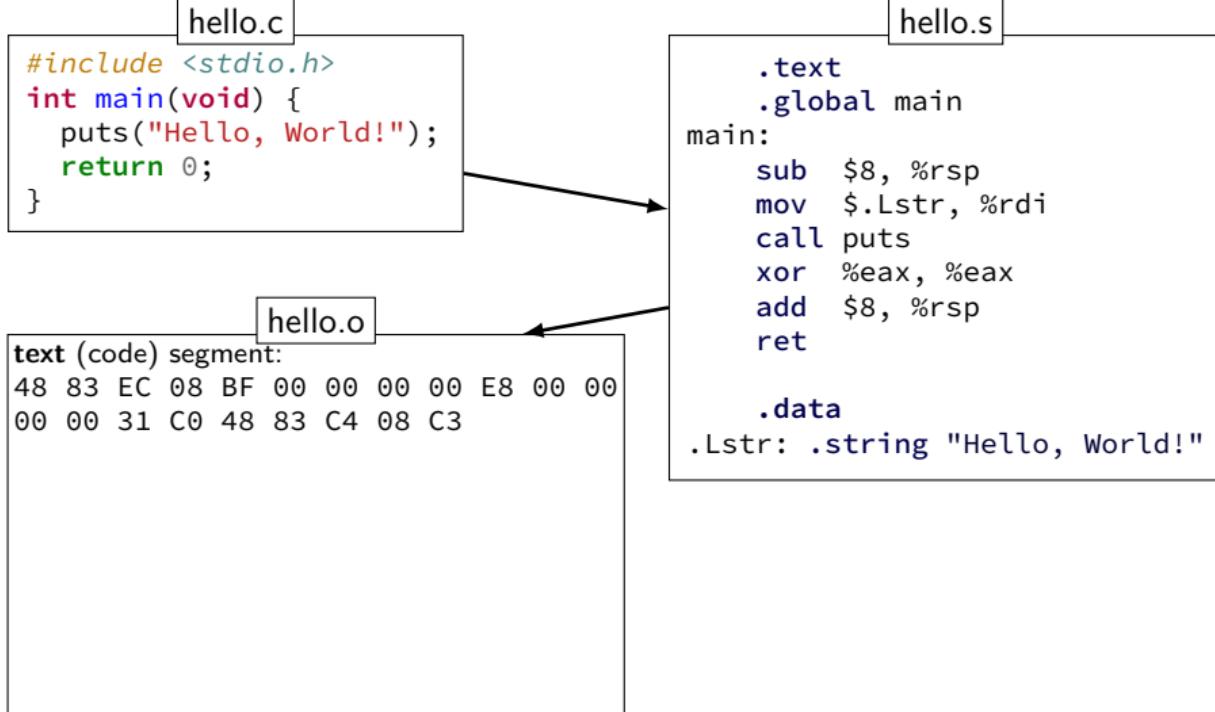
hello.s

```
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    xor    %eax, %eax
    add    $8, %rsp
    ret

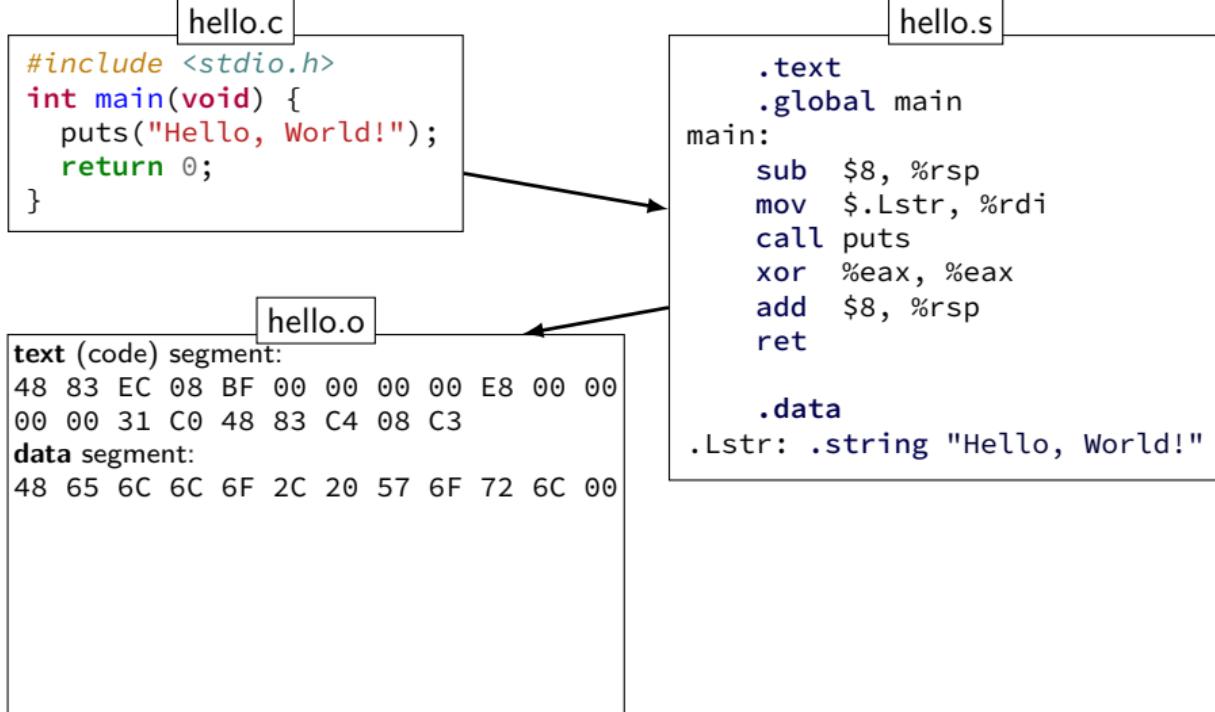
.data
.Lstr: .string "Hello, World!"
```

mark used by other files

what's in those files?



what's in those files?



what's in those files?

hello.c

```
#include <stdio.h>
int main(void) {
    puts("Hello, World!");
    return 0;
}
```

hello.s

```
.text
.global main
main:
    sub $8, %rsp
    mov $.Lstr, %rdi
    call puts
    xor %eax, %eax
    add $8, %rsp
    ret
```

hello.o

text (code) segment:

48 83 EC 08 BF 00 00 00 00 E8 00 00
00 00 31 C0 48 83 C4 08 C3

data segment:

48 65 6C 6C 6F 2C 20 57 6F 72 6C 00

what's in those files?

hello.c

```
#include <stdio.h>
int main(void) {
    puts("Hello, World!");
    return 0;
}
```

hello.s

```
.text
.global main
main:
    sub $8, %rsp
    mov $.Lstr, %rdi
    call puts
    xor %eax, %eax
    add $8, %rsp
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```

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48 65 6C 6C 6F 2C 20 57 6F 72 6C 00

relocations:

take 0s at and replace with
text, byte 5 () data segment, byte 0
text, byte 10 () address of puts

what's in those files?

hello.c

```
#include <stdio.h>
int main(void) {
    puts("Hello, World!");
    return 0;
}
```

hello.s

```
.text
.global main
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    call puts
    xor %eax, %eax
    add $8, %rsp
    ret
```

hello.o

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data segment:

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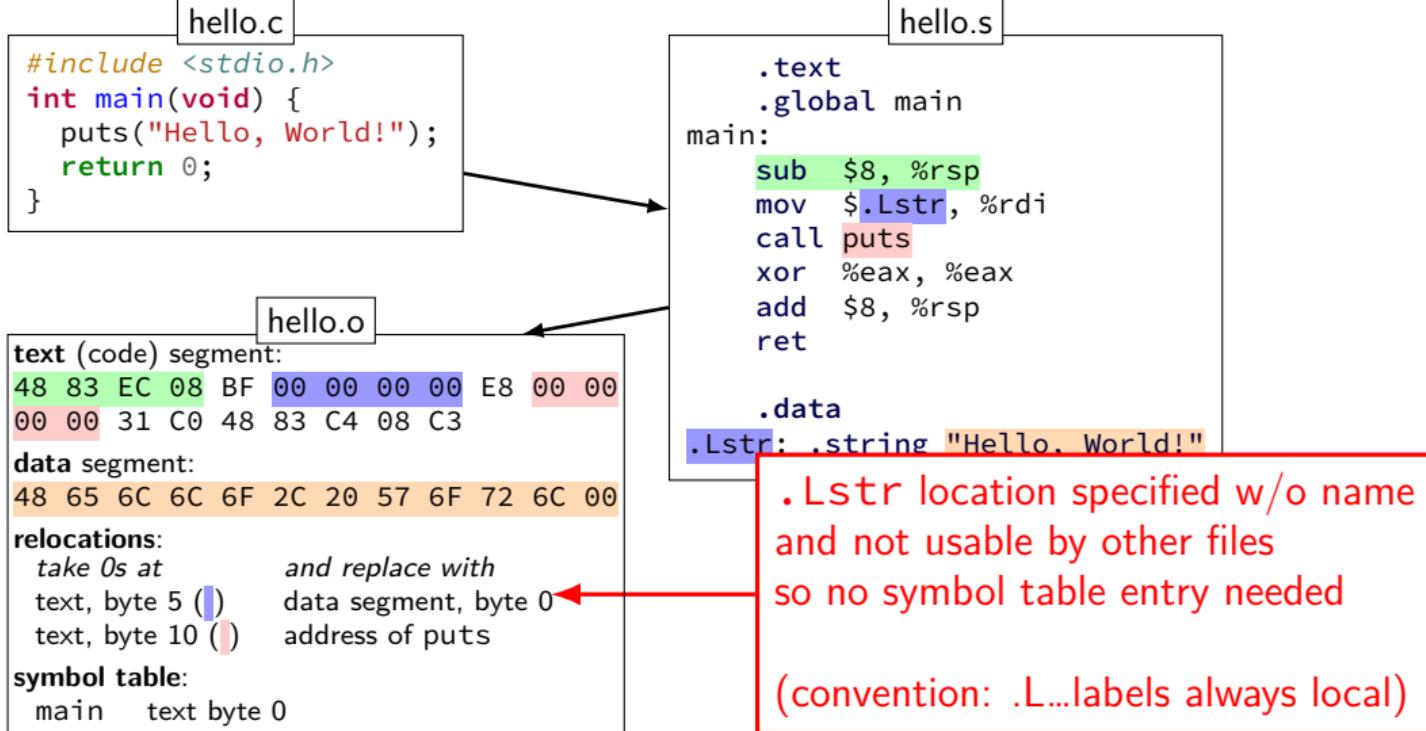
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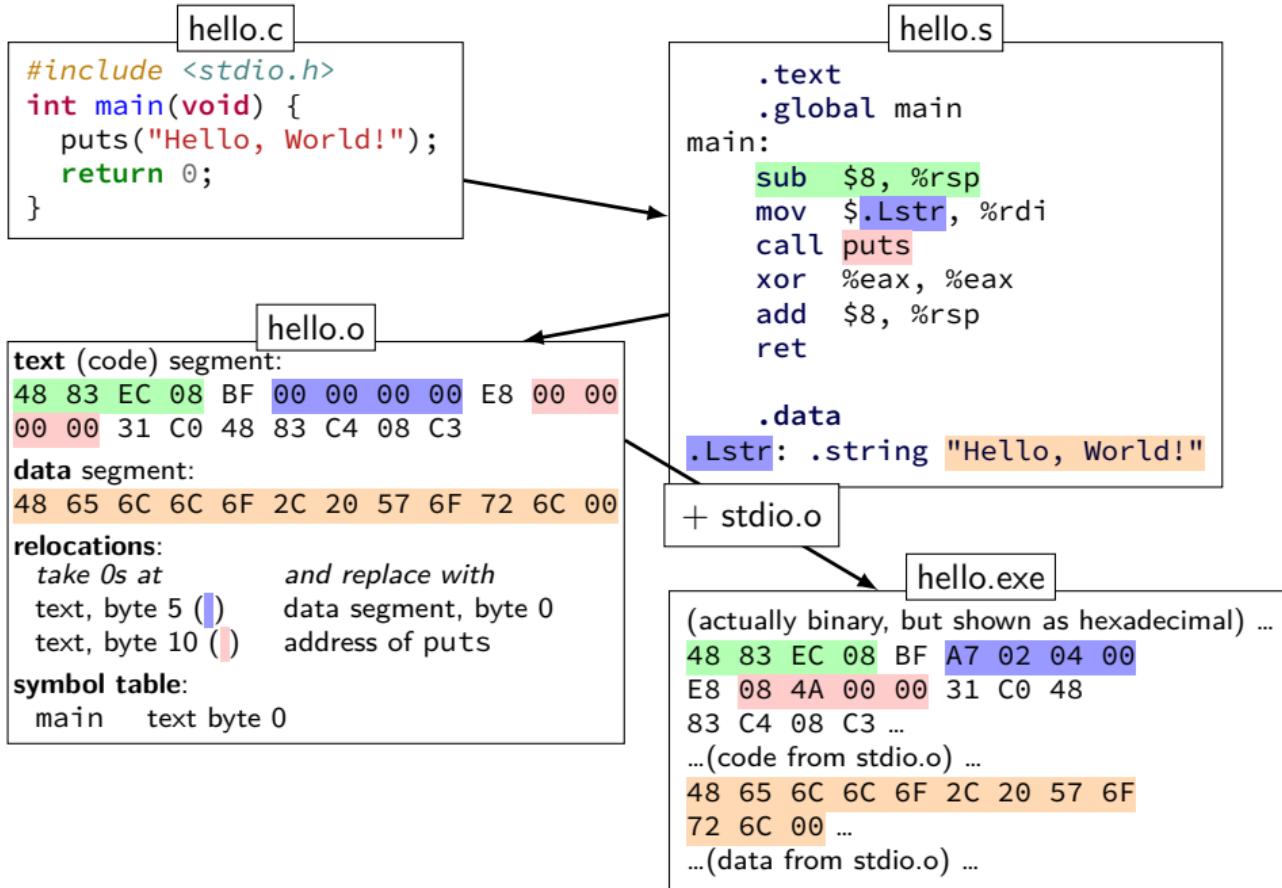
symbol table:

main text byte 0

what's in those files?



what's in those files?



hello.s

```
.section      .rodata.str1.1,"aMS",@progbits
.LC0:
    .string "Hello, World!"
    .text
    .globl  main
main:
    subq    $8, %rsp
    movl    $.LC0, %edi
    call    puts
    movl    $0, %eax
    addq    $8, %rsp
    ret
```

exercise (1)

main.c:

```
1 #include <stdio.h>
2 void sayHello(void) {
3     puts("Hello, World!");
4 }
5 int main(void) {
6     sayHello();
7 }
```

Which files contain the **memory address** of sayHello?

- A. main.s (assembly) D. B and C
- B. main.o (object) E. A, B and C
- C. main.exe (executable) F. something else

exercise (2)

main.c:

```
1 #include <stdio.h>
2 void sayHello(void) {
3     puts("Hello, World!");
4 }
5 int main(void) {
6     sayHello();
7 }
```

Which files contain the **literal ASCII string** of Hello, World!?

- A. main.s (assembly) D. B and C
- B. main.o (object) E. A, B and C
- C. main.exe (executable) F. something else

dynamic linking (very briefly)

dynamic linking — done **when application is loaded**

idea: don't have N copies of `printf` on disk

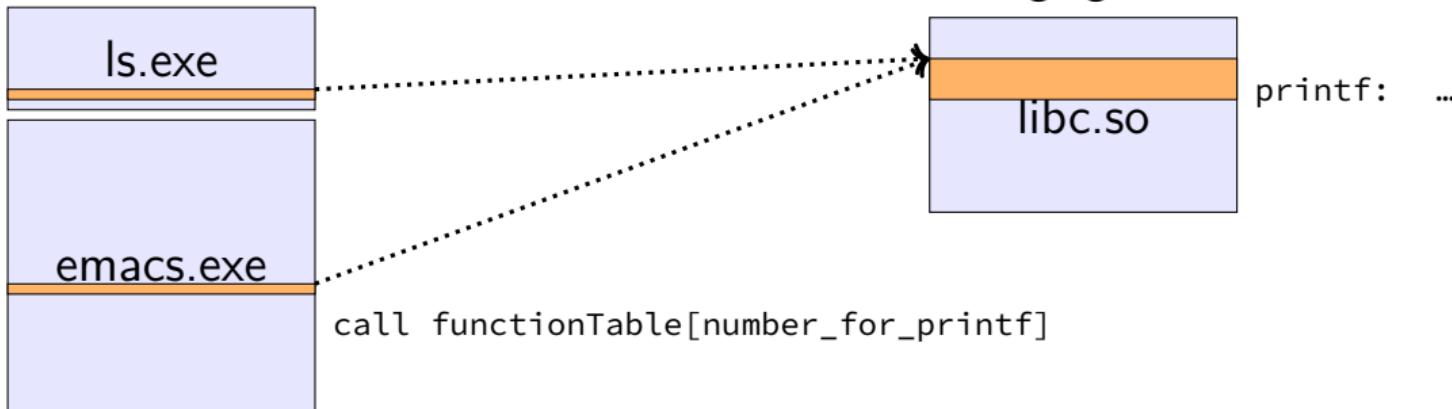
other type of linking: *static* (`gcc -static`)

load executable file + its libraries into memory when app starts

often extra indirection:

`call functionTable[number_for_printf]`

linker fills in `functionTable` instead of changing `calls`



ldd /bin/ls

```
$ ldd /bin/ls
linux-vdso.so.1 => (0x00007ffcc9d8000)
libselinux.so.1 => /lib/x86_64-linux-gnu/libselinux.so.1
                     (0x00007f851756f000)
libc.so.6 => /lib/x86_64-linux-gnu/libc.so.6
                     (0x00007f85171a5000)
libpcre.so.3 => /lib/x86_64-linux-gnu/libpcre.so.3
                     (0x00007f8516f35000)
libdl.so.2 => /lib/x86_64-linux-gnu/libdl.so.2
                     (0x00007f8516d31000)
/lib64/ld-linux-x86-64.so.2 (0x00007f8517791000)
libpthread.so.0 => /lib/x86_64-linux-gnu/libpthread.so.0
                     (0x00007f8516b14000)
```

relocation types

machine code doesn't always use addresses as is

“call function 4303 bytes later”

linker needs to compute “4303”

extra field on relocation list

C Data Types

Varies between machines(!). For **this course**:

type	size (bytes)
char	1
short	2
int	4
long	8

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type	size (bytes)
char	1
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float	4
double	8

C Data Types

Varies between machines(!). For **this course**:

type	size (bytes)
char	1
short	2
int	4
long	8
float	4
double	8
void *	8
<i>anything</i> *	8

truth

bøol

truth

bool

x == 4 is an int

1 if true; 0 if false

false values in C

0

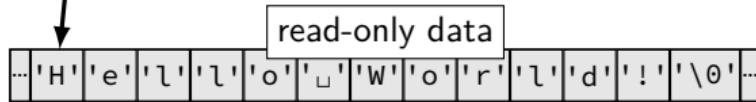
including null pointers — 0 cast to a pointer

strings in C

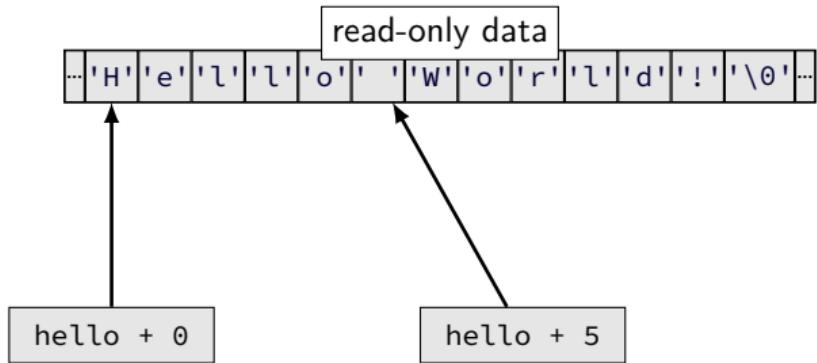
hello (on stack/register)

0x4005C0

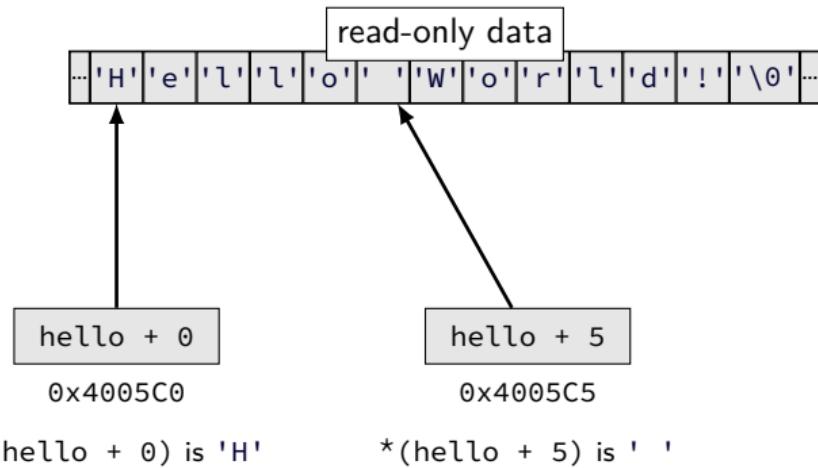
```
int main() {
    const char *hello = "Hello World!";
    ...
}
```



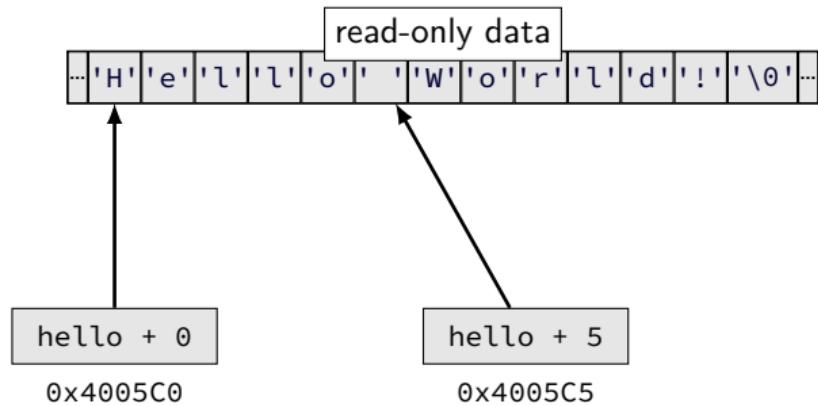
pointer arithmetic



pointer arithmetic



pointer arithmetic



`*(hello + 0)` is 'H'

`*(hello + 5)` is ' '

`hello[0]` is 'H'

`hello[5]` is ' '

arrays and pointers

$\ast(\text{foo} + \text{bar})$ exactly the same as `foo[bar]`

arrays 'decay' into pointers

arrays of non-bytes

array[2] and *(array + 2) still the same

```
1 int numbers[4] = {10, 11, 12, 13};  
2 int *pointer;  
3 pointer = numbers;  
4 *pointer = 20; // numbers[0] = 20;  
5 pointer = pointer + 2;  
6 /* adds 8 (2 ints) to address */  
7 *pointer = 30; // numbers[2] = 30;  
8 // numbers is 20, 11, 30, 13
```

arrays of non-bytes

array[2] and *(array + 2) still the same

```
1 int numbers[4] = {10, 11, 12, 13};  
2 int *pointer;  
3 pointer = numbers;  
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5 pointer = pointer + 2;  
/* adds 8 (2 ints) to address */  
7 *pointer = 30; // numbers[2] = 30;  
8 // numbers is 20, 11, 30, 13
```

exercise

```
1 char foo[4] = "foo";  
2 // {'f', 'o', 'o', '\0'}  
3 char *pointer;  
4 pointer = foo;  
5 *pointer = 'b';  
6 pointer = pointer + 2;  
7 pointer[0] = 'z';  
8 *(foo + 1) = 'a';
```

Final value of foo?

- A. "fao" D. "bao"
- B. "zao" E. something else/crash
- C. "baz"

exercise

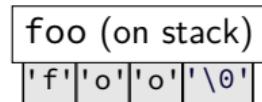
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Final value of foo?

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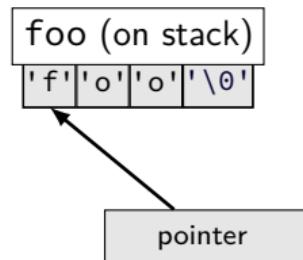
exercise explanation

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1 char foo[4] = "foo";
2     // {'f', 'o', 'o', '\0'}
3 char *pointer;
4 pointer = foo;
5 *pointer = 'b';
6 pointer = pointer + 2;
7 pointer[0] = 'z';
8 *(foo + 1) = 'a';
```



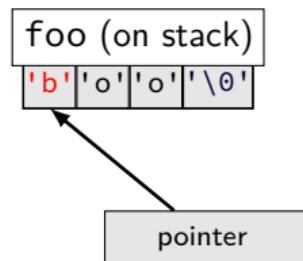
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6 pointer = pointer + 2;
7 pointer[0] = 'z';
8 *(foo + 1) = 'a';
```



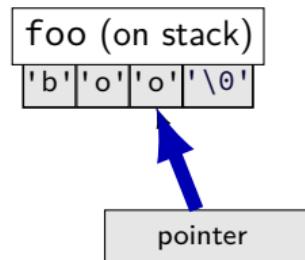
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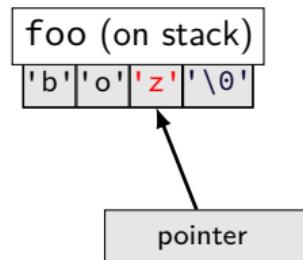
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```



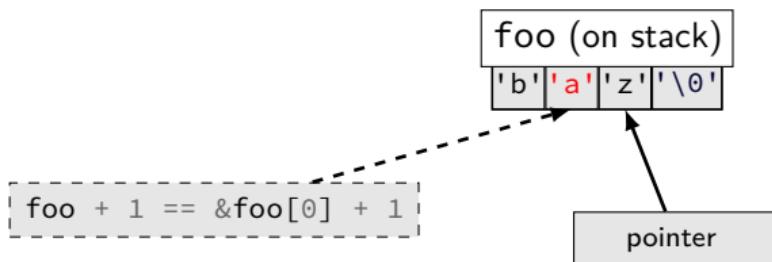
exercise explanation

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1 char foo[4] = "foo";
2     // {'f', 'o', 'o', '\0'}
3 char *pointer;
4 pointer = foo;
5 *pointer = 'b';
6 pointer = pointer + 2;
7 pointer[0] = 'z';    better style: *pointer = 'z';
8 *(foo + 1) = 'a';
```



exercise explanation

```
1 char foo[4] = "foo";
2     // {'f', 'o', 'o', '\0'}
3 char *pointer;
4 pointer = foo;
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6 pointer = pointer + 2;
7 pointer[0] = 'z';    better style: *pointer = 'z';
8 *(foo + 1) = 'a';    better style: foo[1] = 'a';
```



arrays: not quite pointers (1)

```
int array[100];  
int *pointer;
```

Legal: `pointer = array;`
same as `pointer = &(array[0]);`

arrays: not quite pointers (1)

```
int array[100];  
int *pointer;
```

Legal: pointer = array;
same as pointer = &(array[0]);

Illegal: ~~array = pointer;~~

arrays: not quite pointers (2)

```
int array[100];  
int *pointer = array;
```

sizeof(array) == 400

size of all elements

arrays: not quite pointers (2)

```
int array[100];  
int *pointer = array;
```

sizeof(array) == 400
size of all elements

sizeof(pointer) == 8
size of address

arrays: not quite pointers (2)

```
int array[100];
```

```
int *pointer = array;
```

sizeof(array) == 400

size of all elements

sizeof(pointer) == 8

size of address

sizeof(&array[0]) == ???

(&array[0] same as &(array[0]))

C evolution and standards

1978: Kernighan and Ritchie publish *The C Programming Language*
— “K&R C”

very different from modern C

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compiler option: -ansi, -std=c90

looks mostly like modern C

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compiler option: -ansi, -std=c90

looks mostly like modern C

1999: ISO (and ANSI) update C standard — C99

compiler option: -std=c99

adds: declare variables in middle of block

adds: // comments

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adds: // comments

2011, 2017: Second/Third ISO update — C11, C17

undefined behavior example (1)

```
#include <stdio.h>
#include <limits.h>
int test(int number) {
    return (number + 1) > number;
}

int main(void) {
    printf("%d\n", test(INT_MAX));
}
```

undefined behavior example (1)

```
#include <stdio.h>
#include <limits.h>
int test(int number) {
    return (number + 1) > number;
}

int main(void) {
    printf("%d\n", test(INT_MAX));
}
```

without optimizations: 0

undefined behavior example (1)

```
#include <stdio.h>
#include <limits.h>
int test(int number) {
    return (number + 1) > number;
}

int main(void) {
    printf("%d\n", test(INT_MAX));
}
```

without optimizations: 0

with optimizations: 1

undefined behavior example (2)

```
int test(int number) {  
    return (number + 1) > number;  
}
```

Optimized:

```
test:  
    movl    $1, %eax      # eax ← 1  
    ret
```

Less optimized:

```
test:  
    leal    1(%rdi), %eax # eax ← rdi + 1  
    cmpl    %eax, %edi  
    setl    %al             # al ← eax < edi  
    movzbl  %al, %eax     # eax ← al (pad with zeros)  
    ret
```

undefined behavior

compilers can do **whatever they want**

- what you expect
- crash your program

...

common types:

- signed* integer overflow/underflow
- out-of-bounds pointers
- integer divide-by-zero
- writing read-only data
- out-of-bounds shift

undefined behavior

why undefined behavior?

different architectures work differently

- allow compilers to expose whatever processor does “naturally”
- don’t encode any particular machine in the standard

flexibility for optimizations

backup slides

objdump -sx test.o (Linux) (1)

```
test.o:      file format elf64-x86-64
test.o
architecture: i386:x86-64, flags 0x00000011:
HAS_RELOC, HAS_SYMS
start address 0x0000000000000000
```

Sections:

Idx	Name	Size	VMA	LMA	File off	Algn
0	.text	00000000	0000000000000000	0000000000000000	00000040	2**0
			CONTENTS, ALLOC, LOAD, READONLY, CODE			
1	.data	00000000	0000000000000000	0000000000000000	00000040	2**0
			CONTENTS, ALLOC, LOAD, DATA			
2	.bss	00000000	0000000000000000	0000000000000000	00000040	2**0
			ALLOC			
3	.rodata.str1.1	0000000e	0000000000000000	0000000000000000	00000040	2**0
			CONTENTS, ALLOC, LOAD, READONLY, DATA			
4	.text.startup	00000014	0000000000000000	0000000000000000	0000004e	2**0
			CONTENTS, ALLOC, LOAD, RELOC, READONLY, CODE			
5	.comment	0000002b	0000000000000000	0000000000000000	00000062	2**0
			CONTENTS, READONLY			
6	.note.GNU-stack	00000000	0000000000000000	0000000000000000	0000008d	2**0
			CONTENTS, READONLY			
7	.eh_frame	00000030	0000000000000000	0000000000000000	00000090	2**3
			CONTENTS, ALLOC, LOAD, RELOC, READONLY, DATA			

objdump -sx test.o (Linux) (2)

SYMBOL TABLE:

0000000000000000	l	df	*ABS*	0000000000000000	test.c
0000000000000000	l	d	.text	0000000000000000	.text
0000000000000000	l	d	.data	0000000000000000	.data
0000000000000000	l	d	.bss	0000000000000000	.bss
0000000000000000	l	d	.rodata.str1.1	0000000000000000	.rodata.str1.1
0000000000000000	l	d	.text.startup	0000000000000000	.text.startup
0000000000000000	l	d	.note.GNU-stack	0000000000000000	.note.GNU-stack
0000000000000000	l	d	.eh_frame	0000000000000000	.eh_frame
0000000000000000	l		.rodata.str1.1	0000000000000000	.LC0
0000000000000000	l	d	.comment	0000000000000000	.comment
0000000000000000	g	F	.text.startup	000000000000014	main
0000000000000000			*UND*	0000000000000000	_GLOBAL_OFFSET_TABLE_
0000000000000000			*UND*	0000000000000000	puts

columns:

memory address (not yet assigned, so 0)

flags: l=local, g=global, F=function, ...

section (.text, .data, .bss, ...)

offset in section

name of symbol

objdump -sx test.o (Linux) (3)

RELOCATION RECORDS FOR [.text.startup]:

OFFSET	TYPE	VALUE
0000000000000003	R_X86_64_PC32	.LC0-0x0000000000000004
000000000000000c	R_X86_64_PLT32	puts-0x0000000000000004

RELOCATION RECORDS FOR [.eh_frame]:

OFFSET	TYPE	VALUE
0000000000000020	R_X86_64_PC32	.text.startup

Contents of section .rodata.str1.1:

0000 48656c6c 6f2c2057 6f726c64 2100	Hello, World!.
--------------------------------------	----------------

Contents of section .text.startup:

0000 488d3d00 00000048 83ec08e8 00000000	H.=....H.....
0010 31c05ac3	1.Z.

Contents of section .comment:

0000 00474343 3a202855 62756e74 7520372e	.GCC: (Ubuntu 7.
0010 332e302d 32377562 756e7475 317e3138	3.0-27ubuntu1~18
0020 2e303429 20372e33 2e3000	.04) 7.3.0.

Contents of section .eh_frame:

0000 14000000 00000000 017a5200 01781001zR...x..
0010 1b0c0708 90010000 14000000 1c000000
0020 00000000 14000000 004b0e10 480e0800K..H...

compiling switches (1)

```
switch (a) {  
    case 1: ...; break;  
    case 2: ...; break;  
    ...  
    default: ...  
}  
  
// same as if statement?  
cmpq $1, %rax  
je code_for_1  
cmpq $2, %rax  
je code_for_2  
cmpq $3, %rax  
je code_for_3  
...  
jmp code_for_default
```

compiling switches (2)

```
switch (a) {  
    case 1: ...; break;  
    case 2: ...; break;  
    ...  
    case 100: ...; break;  
    default: ...  
}  
  
// binary search  
cmpq $50, %rax  
jl code_for_less_than_50  
cmpq $75, %rax  
jl code_for_50_to_75  
...  
code_for_less_than_50:  
    cmpq $25, %rax  
    jl less_than_25_cases  
...
```

compiling switches (3a)

```
switch (a) {  
    case 1: ...; break;  
    case 2: ...; break;  
    ...  
    case 100: ...; break;  
    default: ...  
}
```

```
// jump table  
cmpq $100, %rax  
jg code_for_default  
cmpq $1, %rax  
jl code_for_default  
jmp *table - 8(%rax, 8)
```

table:

```
// not instructions  
// .quad = 64-bit (4 x 16) constant  
.quad code_for_1  
.quad code_for_2  
.quad code_for_3  
.quad code_for_4  
...
```

compiling switches (3b)

```
jmp *table-8(,%rax,8)
```

suppose RAX = 2,
table located at 0x12500

compiling switches (3b)

```
jmp *table-8(,%rax,8)
```

address	value
...	...
0x124F8	...
table 0x12500	0x13008
table + 0x08 0x12508	0x130A0
table + 0x10 0x12510	0x130C8
table + 0x18 0x12518	0x13110
...	...

suppose RAX = 2,
table located at 0x12500

} table — list of code addresses

...	...
code_for_1 0x13008	...
...	...
...	...
code_for_2 0x130A0	...
...	...

compiling switches (3b)

```
jmp *table-8(,%rax,8)
```

address	value
...	...
0x124F8	...
table 0x12500	0x13008
table + 0x08 0x12508	0x130A0
table + 0x10 0x12510	0x130C8
table + 0x18 0x12518	0x13110
...	...

suppose RAX = 2,
table located at 0x12500

$$(table - 8) + rax \times 8 = \\ 0x124F8 + 0x10 = 0x12508$$



...	...
code_for_1 0x13008	...
...	...
...	...
code_for_2 0x130A0	...
...	...

compiling switches (3b)

```
jmp *table-8(,%rax,8)
```

address	value
...	...
0x124F8	...
table 0x12500	0x13008
table + 0x08 0x12508	0x130A0
table + 0x10 0x12510	0x130C8
table + 0x18 0x12518	0x13110
...	...

...	...
code_for_1 0x13008	...
...	...
...	...
code_for_2 0x130A0	...
...	...

suppose RAX = 2,
table located at 0x12500

pointer to machine code

computed jumps

```
cmpq $100, %rax
jg code_for_default
cmpq $1, %rax
jl code_for_default
// jump to memory[table + rax * 8]
// table of pointers to instructions
jmp *table(,%rax,8)
// intel: jmp QWORD PTR[rax*8 + table]
...
table:
.quad code_for_1
.quad code_for_2
.quad code_for_3
...
```

struct

```
struct rational {  
    int numerator;  
    int denominator;  
};  
// ...  
struct rational two_and_a_half;  
two_and_a_half.numerator = 5;  
two_and_a_half.denominator = 2;  
struct rational *pointer = &two_and_a_half;  
printf("%d/%d\n",  
    pointer->numerator,  
    pointer->denominator);
```

struct

```
struct rational {  
    int numerator;  
    int denominator;  
};  
// ...  
struct rational two_and_a_half;  
two_and_a_half.numerator = 5;  
two_and_a_half.denominator = 2;  
struct rational *pointer = &two_and_a_half;  
printf("%d/%d\n",  
    pointer->numerator,  
    pointer->denominator);
```

typedef

instead of writing:

```
...
unsigned int a;
unsigned int b;
unsigned int c;
```

can write:

```
typedef unsigned int uint;
```

```
...
uint a;
uint b;
uint c;
```

typedef struct (1)

```
struct other_name_for_rational {
    int numerator;
    int denominator;
};

typedef struct other_name_for_rational rational;
// ...
rational two_and_a_half;
two_and_a_half.numerator = 5;
two_and_a_half.denominator = 2;
rational *pointer = &two_and_a_half;
printf("%d/%d\n",
       pointer->numerator,
       pointer->denominator);
```

typedef struct (1)

```
struct other_name_for_rational {
    int numerator;
    int denominator;
};

typedef struct other_name_for_rational rational;
// ...
rational two_and_a_half;
two_and_a_half.numerator = 5;
two_and_a_half.denominator = 2;
rational *pointer = &two_and_a_half;
printf("%d/%d\n",
       pointer->numerator,
       pointer->denominator);
```

```
typedef struct (2)
struct other_name_for_rational {
    int numerator;
    int denominator;
};
typedef struct other_name_for_rational rational;
// same as:
typedef struct other_name_for_rational {
    int numerator;
    int denominator;
} rational;
```

```
typedef struct (2)
struct other_name_for_rational {
    int numerator;
    int denominator;
};
typedef struct other_name_for_rational rational;
// same as:
typedef struct other_name_for_rational {
    int numerator;
    int denominator;
} rational;
```

```
typedef struct (2)
struct other_name_for_rational {
    int numerator;
    int denominator;
};
typedef struct other_name_for_rational rational;
// same as:
typedef struct other_name_for_rational {
    int numerator;
    int denominator;
} rational;
// almost the same as:
typedef struct {
    int numerator;
    int denominator;
} rational;
```

```
typedef struct (3)
struct other_name_for_rational {
    int numerator;
    int denominator;
};

typedef struct other_name_for_rational rational;

valid ways to declare an instance:
struct other_name_for_rational some_variable;
rational some_variable;

INVALID ways:
/* INVALID: */ struct rational some_variable;
/* INVALID: */ other_name_for_rational some_variable;
```

typedef struct (3)

```
struct other_name_for_rational {  
    int numerator;  
    int denominator;  
};  
typedef struct other_name_for_rational rational;
```

valid ways to declare an instance:

```
struct other_name_for_rational some_variable;  
rational some_variable;
```

INVALID ways:

```
/* INVALID: */ struct rational some_variable;  
/* INVALID: */ other_name_for_rational some_variable;
```

```
typedef struct (3)
struct other_name_for_rational {
    int numerator;
    int denominator;
};

typedef struct other_name_for_rational rational;

valid ways to declare an instance:
struct other_name_for_rational some_variable;
rational some_variable;

INVALID ways:
/* INVALID: */ struct rational some_variable;
/* INVALID: */ other_name_for_rational some_variable;
```

structs aren't references

```
typedef struct {  
    long a; long b; long c;  
} triple;  
...
```

```
triple foo;  
foo.a = foo.b = foo.c = 3;  
triple bar = foo;  
bar.a = 4;  
// foo is 3, 3, 3  
// bar is 4, 3, 3
```

...
return address
callee saved
registers
foo.c
foo.b
foo.a
bar.c
bar.b
bar.a