

program layout / C stuff

# changelog

5 September 2022: adjust “what’s in those files” to indicate that we’re putting in the *address* of the data segment

# last time

lea (load effective address)

effective address = address computed in middle of running instruction

lea computes address, places in destination register

processor doesn't check/care if "address" is valid in memory

lea often used for non-address computation

condition codes

special 1-bit registers w/ results of "last arithmetic"

ZF = was zero?; SF = was negative [sign bit]?

also OF, CF for overflow

cmp = sub but only set condition codes

jXX — named after comparing to 0 (or subtracting)

converting control flow to assembly

# while-to-assembly (1)

```
while (x >= 0) {  
    foo()  
    x--;  
}
```

---

# while-to-assembly (1)

```
while (x >= 0) {  
    foo()  
    x--;  
}
```

---

```
start_loop:  
    if (x < 0) goto end_loop;  
    foo()  
    x--;  
    goto start_loop;  
end_loop:
```

## while-to-assembly (2)

```
start_loop:  
    if (x < 0) goto end_loop;  
    foo()  
    x--;  
    goto start_loop;  
end_loop:
```

---

```
start_loop:  
    cmpq $0, %r12  
    jl end_loop // jump if r12 - 0 < 0  
    call foo  
    subq $1, %r12  
    jmp start_loop
```

# while — levels of optimization

```
while (b < 10) { foo(); b += 1; }
```

```
start_loop:  
    cmpq $10, %rbx  
    jge end_loop  
    call foo  
    addq $1, %rbx  
    jmp start_loop
```

```
end_loop:
```

```
    ...  
    ...  
    ...  
    ...
```

# while — levels of optimization

```
while (b < 10) { foo(); b += 1; }
```

```
start_loop:  
    cmpq $10, %rbx  
    jge end_loop  
    call foo  
    addq $1, %rbx  
    jmp start_loop  
end_loop:  
    ...  
    ...  
    ...  
    ...
```

```
        cmpq $10, %rbx  
        jge end_loop  
start_loop:  
    call foo  
    addq $1, %rbx  
    cmpq $10, %rbx  
    jne start_loop  
end_loop:  
    ...  
    ...  
    ...
```

# while — levels of optimization

```
while (b < 10) { foo(); b += 1; }
```

```
start_loop:  
    cmpq $10, %rbx  
    jge end_loop  
    call foo  
    addq $1, %rbx  
    jmp start_loop  
end_loop:  
    ...  
    ...  
    ...  
    ...
```

```
        cmpq $10, %rbx  
        jge end_loop  
start_loop:  
    call foo  
    addq $1, %rbx  
    cmpq $10, %rbx  
    jne start_loop  
end_loop:  
    ...  
    ...  
    ...
```

```
        cmpq $10, %rbx  
        jge end_loop  
        movq $10, %rax  
        subq %rbx, %rax  
        movq %rax, %rbx  
start_loop:  
    call foo  
    decq %rbx  
    jne start_loop  
    movq $10, %rbx  
end_loop:
```

# compiling switches (1)

```
switch (a) {  
    case 1: ...; break;  
    case 2: ...; break;  
    ...  
    default: ...  
}  
  
// same as if statement?  
cmpq $1, %rax  
je code_for_1  
cmpq $2, %rax  
je code_for_2  
cmpq $3, %rax  
je code_for_3  
...  
jmp code_for_default
```

## compiling switches (2)

```
switch (a) {  
    case 1: ...; break;  
    case 2: ...; break;  
    ...  
    case 100: ...; break;  
    default: ...  
}  
  
// binary search  
cmpq $50, %rax  
jl code_for_less_than_50  
cmpq $75, %rax  
jl code_for_50_to_75  
...  
code_for_less_than_50:  
    cmpq $25, %rax  
    jl less_than_25_cases  
    ...
```

# compiling switches (3a)

```
switch (a) {  
    case 1: ...; break;  
    case 2: ...; break;  
    ...  
    case 100: ...; break;  
    default: ...  
}
```

```
// jump table  
cmpq $100, %rax  
jg code_for_default  
cmpq $1, %rax  
jl code_for_default  
jmp *table - 8(,%rax,8)
```

table:

```
// not instructions  
// .quad = 64-bit (4 x 16) constant  
.quad code_for_1  
.quad code_for_2  
.quad code_for_3  
.quad code_for_4  
...
```

## compiling switches (3b)

```
jmp *table-8(,%rax,8)
```

suppose RAX = 2,  
table located at 0x12500

# compiling switches (3b)

```
jmp *table-8(,%rax,8)
```

address	value
...	...
0x124F8	...
table 0x12500	0x13008
table + 0x08 0x12508	0x130A0
table + 0x10 0x12510	0x130C8
table + 0x18 0x12518	0x13110
...	...

suppose RAX = 2,  
table located at 0x12500

} table — list of code addresses

...	...
code_for_1 0x13008	...
...	...
...	...
code_for_2 0x130A0	...
...	...

# compiling switches (3b)

```
jmp *table-8(,%rax,8)
```

address	value
...	...
0x124F8	...
table 0x12500	0x13008
table + 0x08 0x12508	0x130A0
table + 0x10 0x12510	0x130C8
table + 0x18 0x12518	0x13110
...	...

suppose RAX = 2,  
table located at 0x12500

$$(table - 8) + rax \times 8 = \\ 0x124F8 + 0x10 = 0x12508$$

...	...
code_for_1 0x13008	...
...	...
...	...
code_for_2 0x130A0	...
...	...

# compiling switches (3b)

```
jmp *table-8(,%rax,8)
```

address	value
...	...
0x124F8	...
table 0x12500	0x13008
table + 0x08 0x12508	0x130A0
table + 0x10 0x12510	0x130C8
table + 0x18 0x12518	0x13110
...	...
...	...
code_for_1 0x13008	[...]
...	...
...	...
code_for_2 0x130A0	[...]
...	...

suppose RAX = 2,  
table located at 0x12500

pointer to machine code

# computed jumps

```
cmpq $100, %rax
jg code_for_default
cmpq $1, %rax
jl code_for_default
// jump to memory[table + rax * 8]
// table of pointers to instructions
jmp *table(,%rax,8)
// intel: jmp QWORD PTR[rax*8 + table]
```

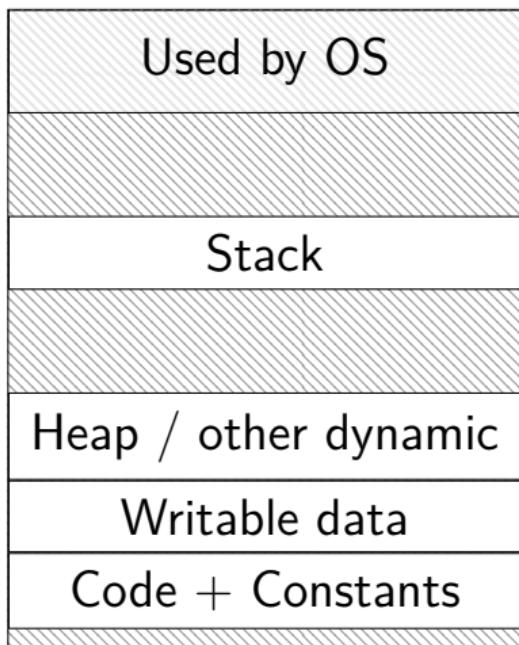
...

table:

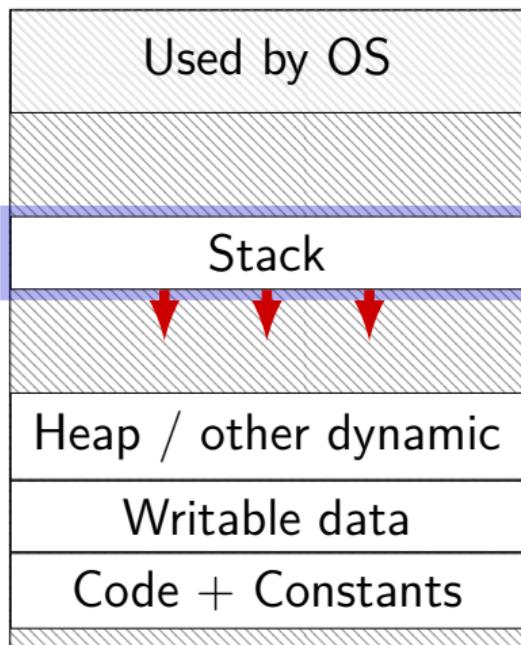
```
.quad code_for_1
.quad code_for_2
.quad code_for_3
```

...

# program memory (x86-64 Linux)



# program memory (x86-64 Linux)



0xFFFF FFFF FFFF FFFF

0xFFFF 8000 0000 0000

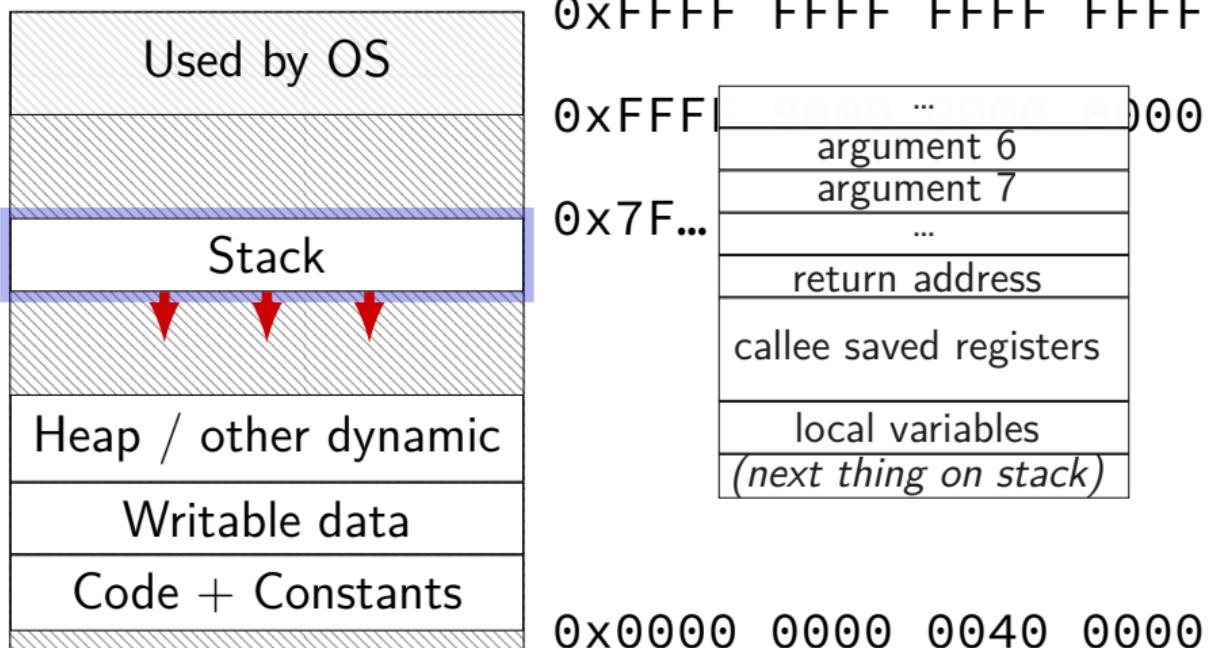
0x7F...

stack *grows down*

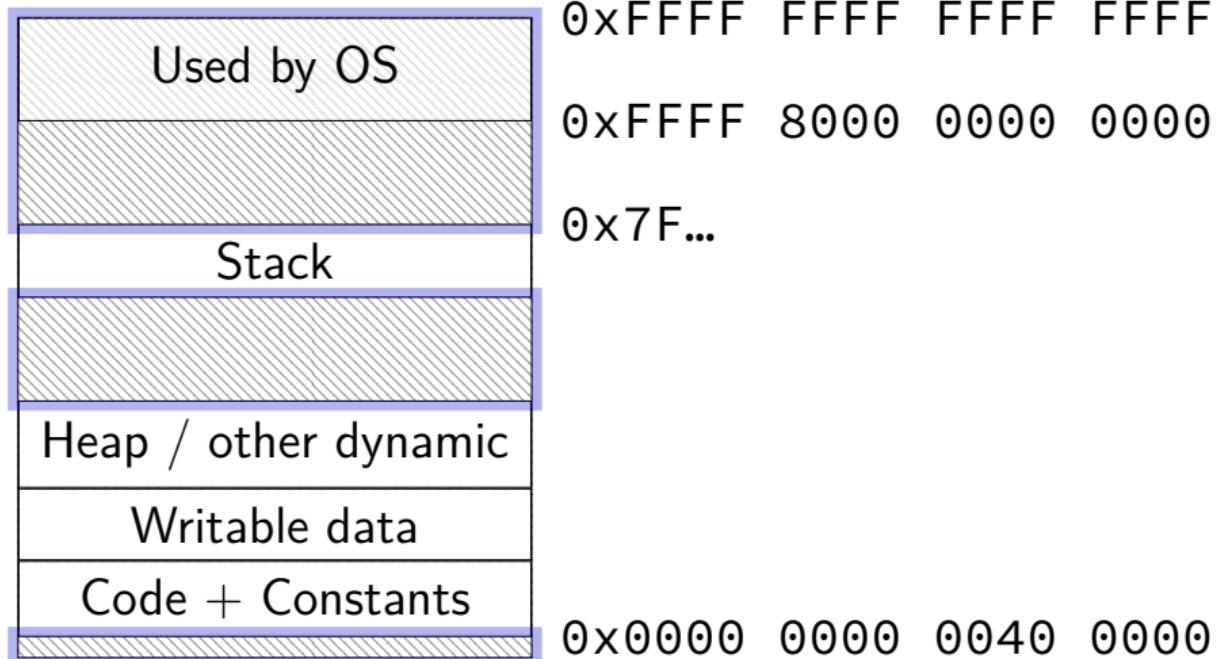
“top” has smallest address

0x0000 0000 0040 0000

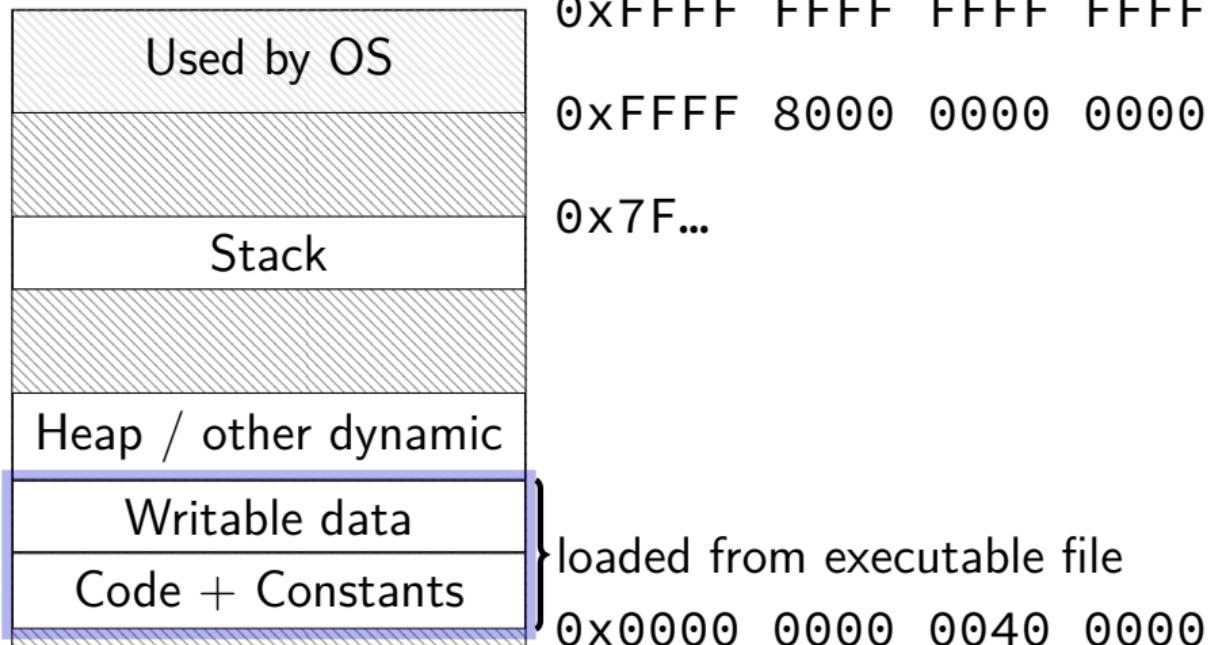
# program memory (x86-64 Linux)



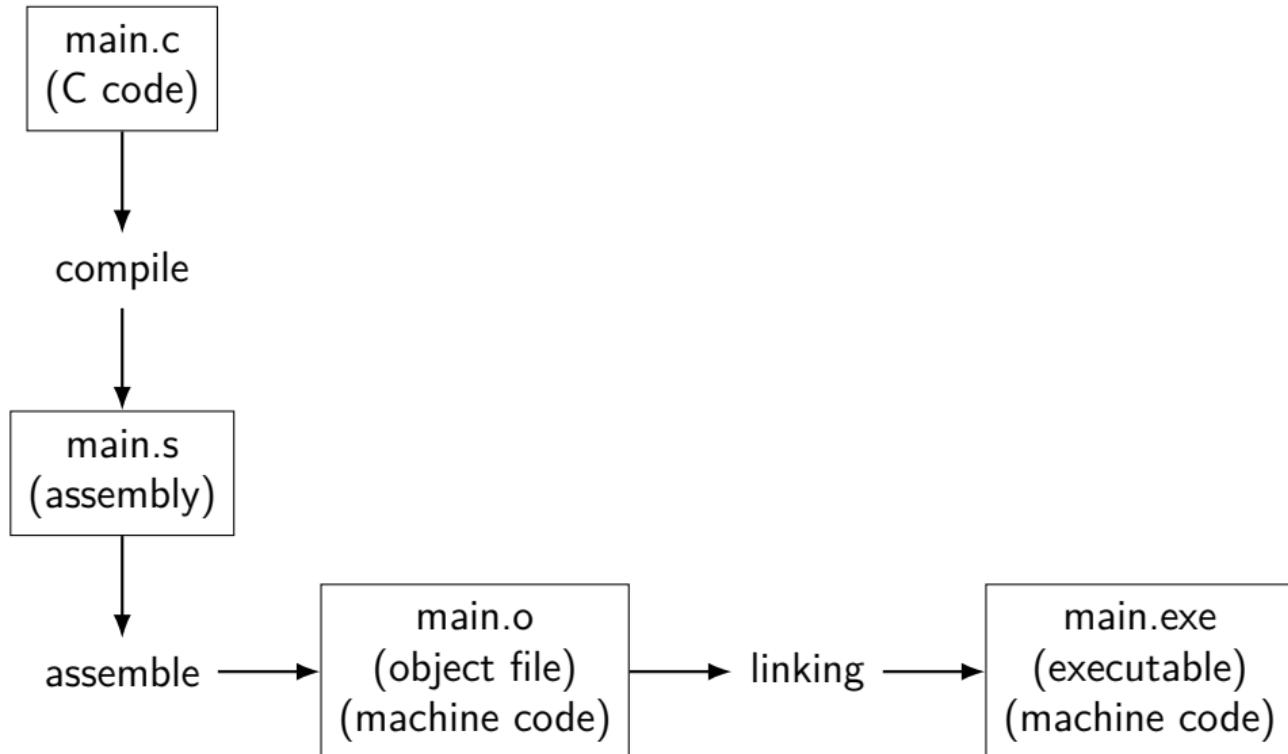
# program memory (x86-64 Linux)



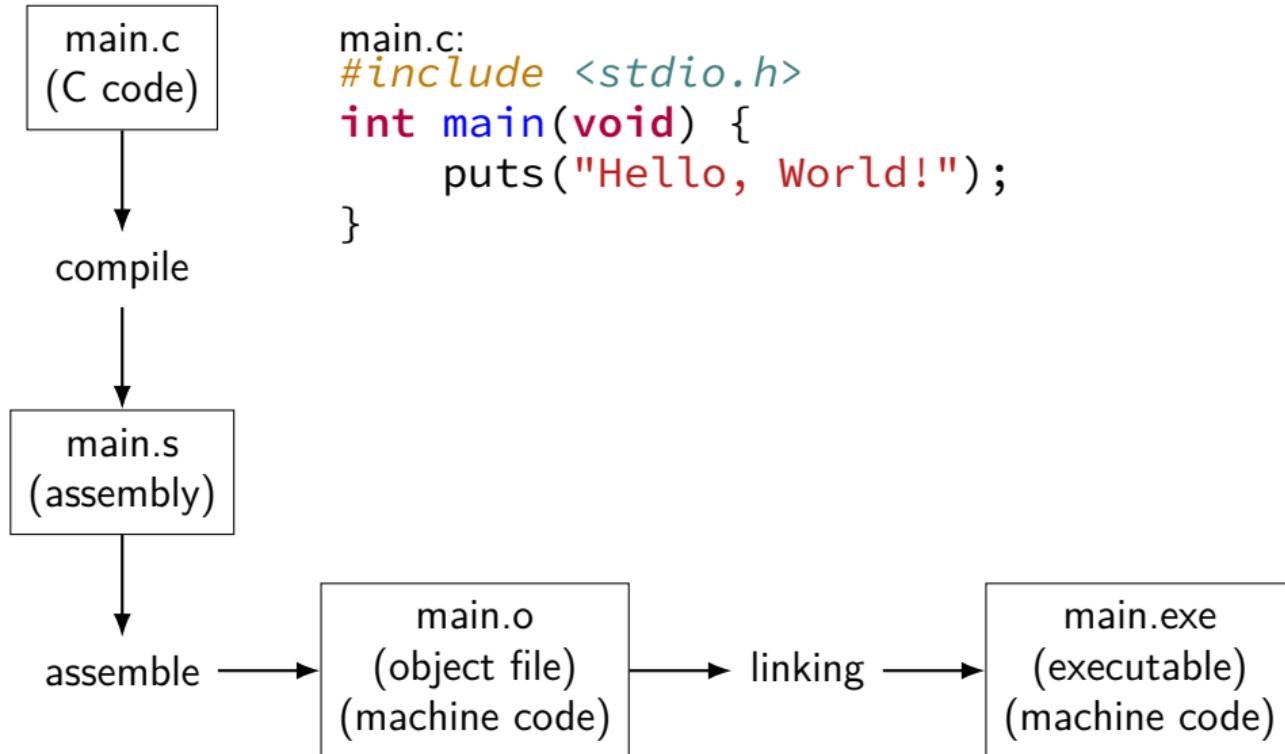
# program memory (x86-64 Linux)



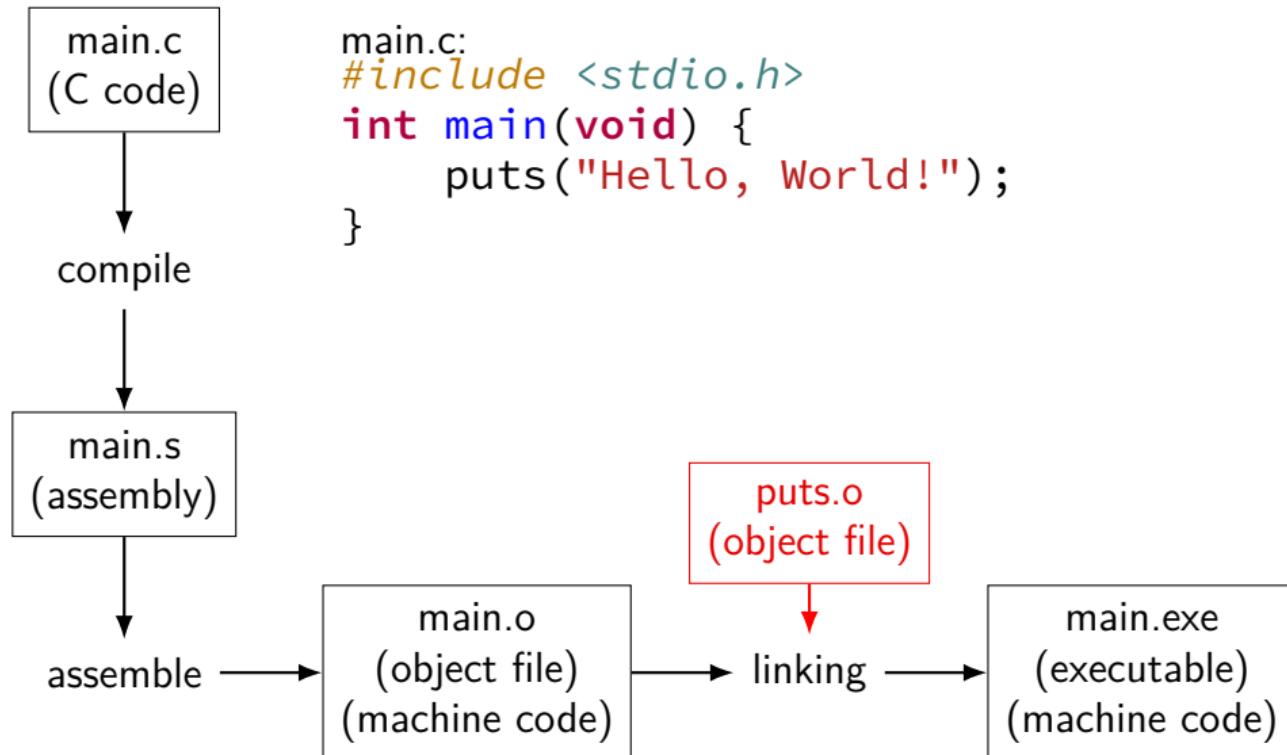
# compilation pipeline



# compilation pipeline



# compilation pipeline



# compilation commands

compile: `gcc -S file.c` ⇒ `file.s` (assembly)

assemble: `gcc -c file.s` ⇒ `file.o` (object file)

link: `gcc -o file file.o` ⇒ `file` (executable)

c+a: `gcc -c file.c` ⇒ `file.o`

c+a+l: `gcc -o file file.c` ⇒ `file`

...

# combining assembly files?

main.s

```
.text
.global main
main:
    mov $str, %rdi
    call puts
    ret
.data
str:
.string "Hello!"
```

puts.s

```
.text
.global puts
puts:
    ...
    call putchar
    ...
```

# combining assembly files?

main.s

```
.text
.global main
main:
    mov $str, %rdi
    call puts
    ret
.data
str:
.string "Hello!"
```

puts.s

```
.text
.global puts
puts:
...
    call putchar
...
```

combined?

```
.text
main:
    mov $str, %rdi
    call puts
    ret
.data
str:
.string "Hello"

.text
puts:
...
    call putchar
...
```

# combining assembly files?

main.s

```
.text  
.global main  
main:  
    mov $str, %rdi  
    call puts  
    ret  
.data  
str:  
    .string "Hello!"
```

puts.s

```
.text  
.global puts  
puts:  
    ...  
    call putchar  
    ...
```

combined?

```
.text  
main:  
    mov $str, %rdi  
    call puts  
    ret  
.data  
str:  
    .string "Hello"  
  
.text  
puts:  
    ...  
    call putchar  
    ...
```

problem:  
how many times  
do we generate  
library machine  
code?

# combining assembly files?

main.s

```
.text  
.global main  
main:  
    mov $str, %rdi  
    call puts  
    ret  
.data  
str:  
.string "Hello!"
```

combined?

```
.text  
main:  
    mov $str, %rdi  
    call puts  
    ret  
.data  
str:  
.string "Hello"
```

puts.s

```
.text  
.global puts  
puts:  
    ...  
    call putchar  
    ...
```

normal part of making machine code:  
choosing + filling in addresses

# combining assembly files?

main.s

```
.text  
.global main  
main:  
    mov $str, %rdi  
    call puts  
    ret  
.data  
str:  
    .string "Hello!"
```

puts.s

```
.text  
.global puts  
puts:  
    ...  
    call putchar  
    ...
```

combined?

```
.text  
main:  
    mov $str, %rdi  
    call puts  
    ret  
.data  
str:  
    .string "Hello"  
  
.text  
puts:  
    ...  
    call putchar  
    ...
```

(with addrs)

```
0x10000:  
    mov $0x20000, %rdi  
    call 0x10040  
    ret  
  
0x10040:  
    ...  
    call 0x10800  
    ...  
    ...  
  
0x20000:  
    .string "Hello"
```

# combining assembly files?

main.s

```
.text  
.global main  
main:  
    mov $str, %rdi  
    call puts  
    ret  
.data  
str:  
.string "Hello!"
```

puts.s

```
.text  
.global puts  
puts:  
    ...  
    call putchar  
    ...
```

main.s as machine code

```
mov $???str, %rdi  
call ???puts  
ret  
.string "Hello"
```

puts.s as machine code

```
...  
call ???putchar  
...
```

idea:  
translate each .s  
to machine code  
and combine later

problem:  
can't put labels  
in machine code

# what's in those files?

hello.c

```
#include <stdio.h>
int main(void) {
    puts("Hello, World!");
    return 0;
}
```

# what's in those files?

hello.c

```
#include <stdio.h>
int main(void) {
    puts("Hello, World!");
    return 0;
}
```

hello.s

```
.text
.global main
main:
    sub $8, %rsp
    mov $.Lstr, %rdi
    call puts
    xor %eax, %eax
    add $8, %rsp
    ret

.data
.Lstr: .string "Hello, World!"
```

# what's in those files?

hello.c

```
#include <stdio.h>
int main(void) {
    puts("Hello, World!");
    return 0;
}
```

hello.s

```
.text
.global main
main:
    sub $8, %rsp
    mov $.Lstr, %rdi
    call puts
    xor %eax, %eax
    add $8, %rsp
    ret

.data
.Lstr: .string "Hello, Wor
```

hello.s (Intel syntax)

```
.text
main:
    sub RSP, 8
    mov RDI, .Lstr
    call puts
    xor EAX, EAX
    add RSP, 8
    ret

.data
.Lstr: .string "Hello, Wor
```

# what's in those files?

hello.c

```
#include <stdio.h>
int main(void) {
    puts("Hello, World!");
    return 0;
}
```

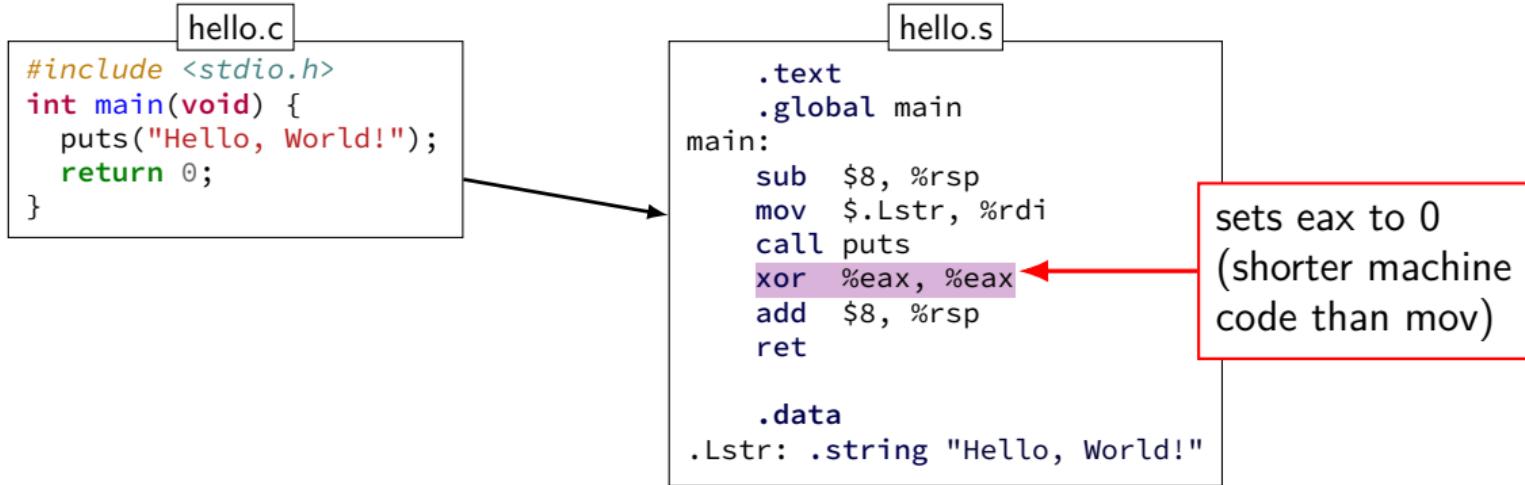
hello.s

```
.text
.global main
main:
    sub $8, %rsp
    mov $.Lstr, %rdi
    call puts
    xor %eax, %eax
    add $8, %rsp
    ret

.data
.Lstr: .string "Hello, World!"
```

Linux x86-64  
calling convention:  
stack addr. must be  
multiple of 16

# what's in those files?



# what's in those files?

hello.c

```
#include <stdio.h>
int main(void) {
    puts("Hello, World!");
    return 0;
}
```

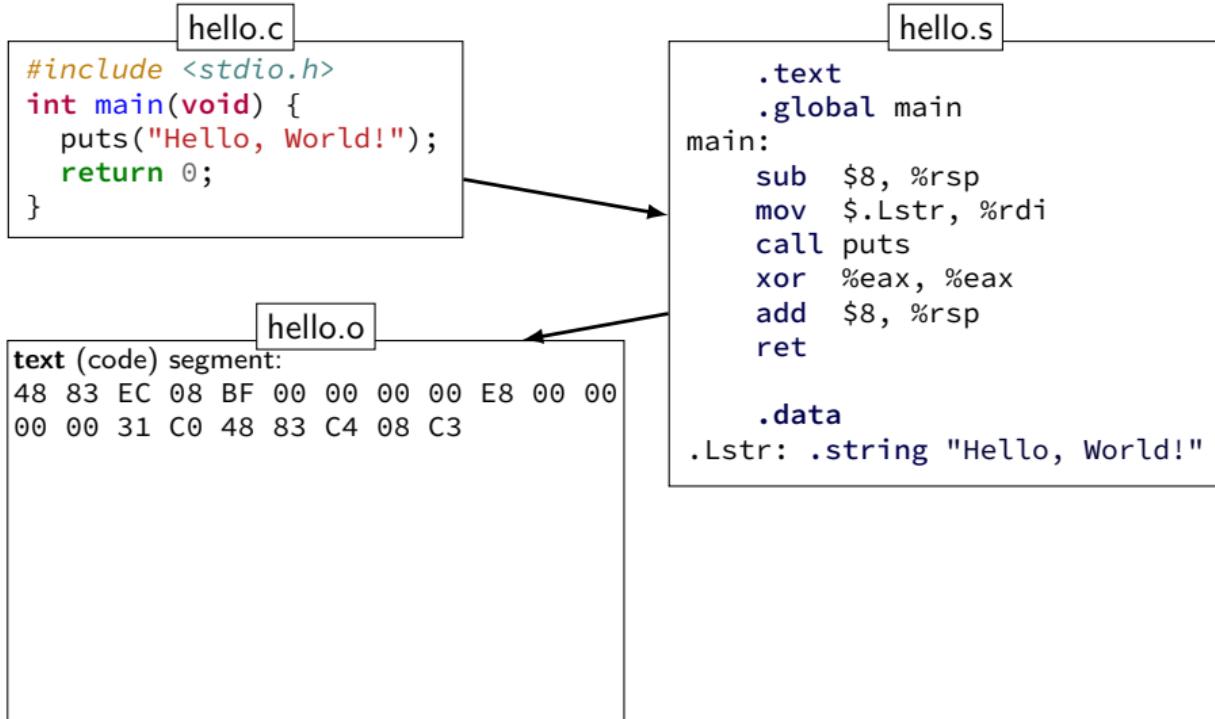
hello.s

```
.text
.global main
main:
    sub    $8, %rsp
    mov    $.Lstr, %rdi
    call   puts
    xor    %eax, %eax
    add    $8, %rsp
    ret

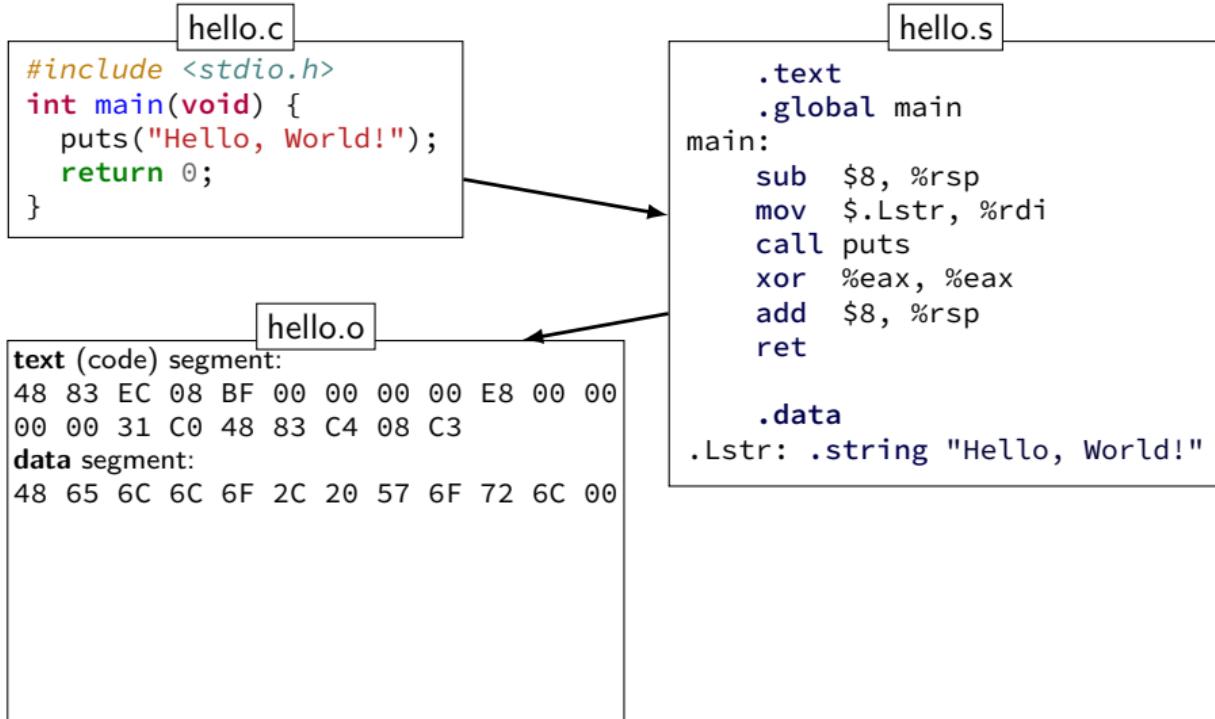
.data
.Lstr: .string "Hello, World!"
```

mark used by other files

# what's in those files?



# what's in those files?



# what's in those files?

hello.c

```
#include <stdio.h>
int main(void) {
    puts("Hello, World!");
    return 0;
}
```

hello.s

```
.text
.global main
main:
    sub $8, %rsp
    mov $.Lstr, %rdi
    call puts
    xor %eax, %eax
    add $8, %rsp
    ret
```

.data

```
.Lstr: .string "Hello, World!"
```

hello.o

text (code) segment:

```
48 83 EC 08 BF 00 00 00 00 E8 00 00
00 00 31 C0 48 83 C4 08 C3
```

data segment:

```
48 65 6C 6C 6F 2C 20 57 6F 72 6C 00
```

# what's in those files?

hello.c

```
#include <stdio.h>
int main(void) {
    puts("Hello, World!");
    return 0;
}
```

hello.s

```
.text
.global main
main:
    sub $8, %rsp
    mov $.Lstr, %rdi
    call puts
    xor %eax, %eax
    add $8, %rsp
    ret
```

.data

```
.Lstr: .string "Hello, World!"
```

hello.o

text (code) segment:

48 83 EC 08 BF 00 00 00 00 E8 00 00  
00 00 31 C0 48 83 C4 08 C3

data segment:

48 65 6C 6C 6F 2C 20 57 6F 72 6C 00

relocations:

take 0s at text, byte 5 () and replace with addr. of data segment, byte 0  
text, byte 10 () puts

# what's in those files?

hello.c

```
#include <stdio.h>
int main(void) {
    puts("Hello, World!");
    return 0;
}
```

hello.s

```
.text
.global main
main:
    sub $8, %rsp
    mov $.Lstr, %rdi
    call puts
    xor %eax, %eax
    add $8, %rsp
    ret
```

```
.data
.Lstr: .string "Hello, World!"
```

hello.o

text (code) segment:

48 83 EC 08 BF 00 00 00 00 E8 00 00  
00 00 31 C0 48 83 C4 08 C3

data segment:

48 65 6C 6C 6F 2C 20 57 6F 72 6C 00

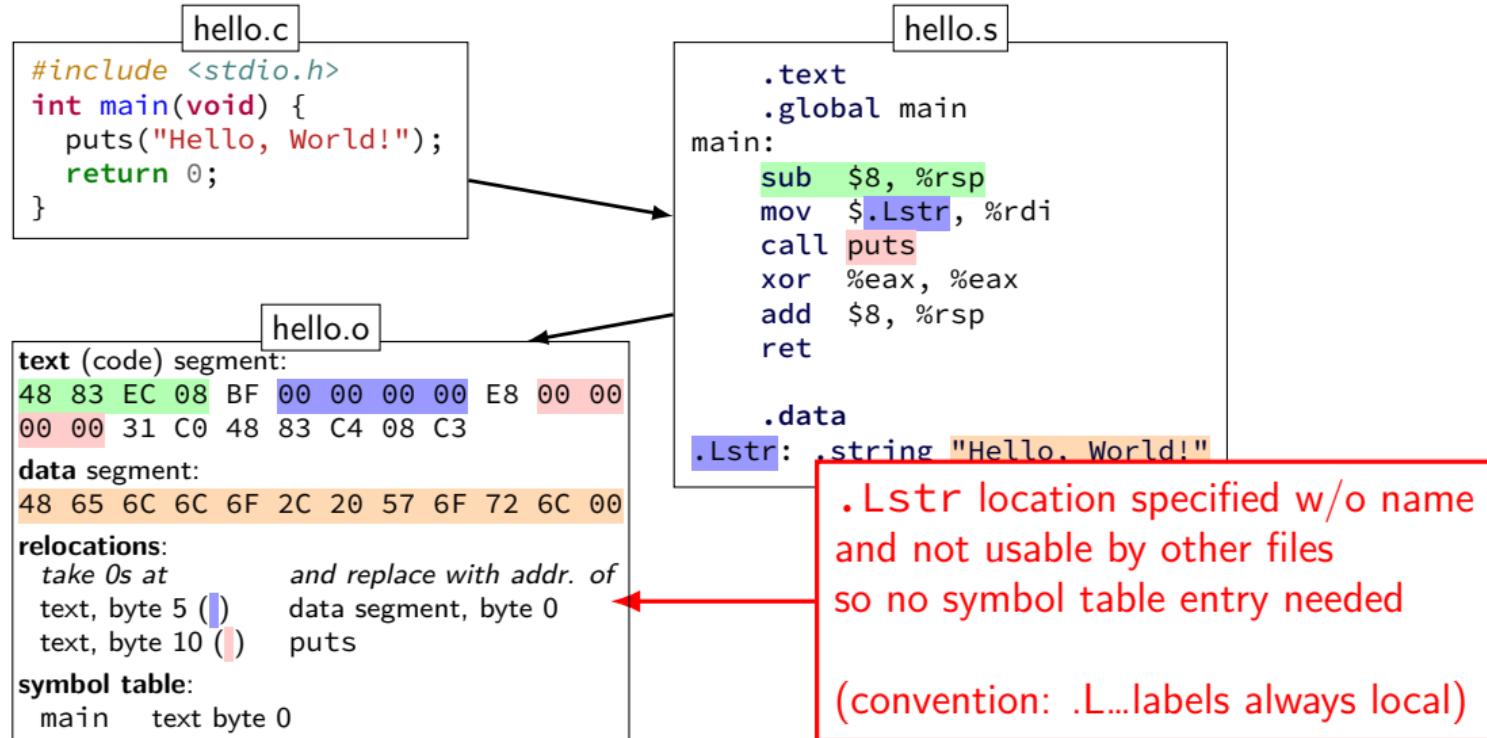
relocations:

take 0s at and replace with addr. of  
text, byte 5 () data segment, byte 0  
text, byte 10 () puts

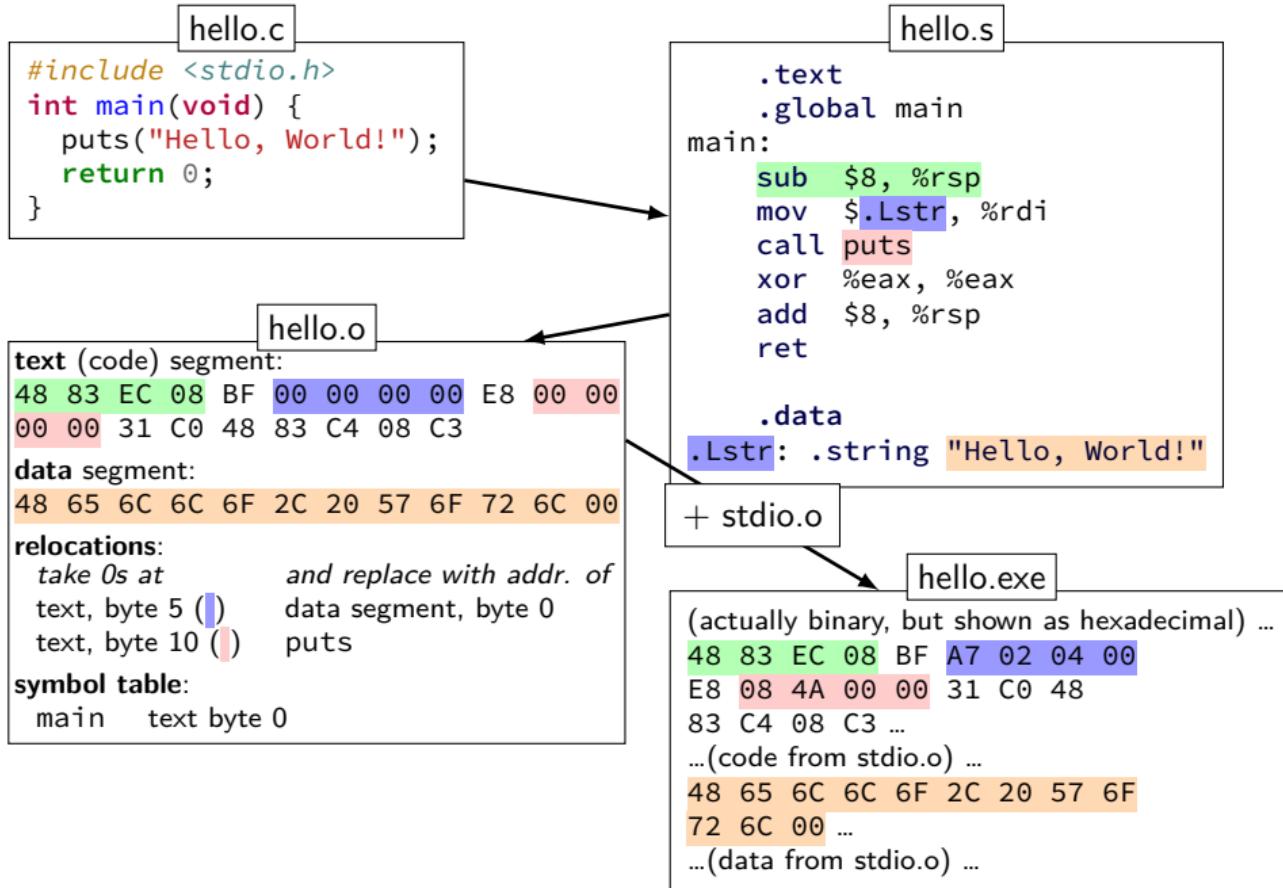
symbol table:

main text byte 0

# what's in those files?



# what's in those files?



# hello.s

```
.section      .rodata.str1.1,"aMS",@progb
.LC0:
.string "Hello, World!"
.text
.globl main
main:
    subq    $8, %rsp
    movl    $.LC0, %edi
    call    puts
    movl    $0, %eax
    addq    $8, %rsp
    ret
```

# exercise (1)

main.c:

```
1 #include <stdio.h>
2 void sayHello(void) {
3     puts("Hello, World!");
4 }
5 int main(void) {
6     sayHello();
7 }
```

Which files likely contain the **memory address** of sayHello?

- A. main.s (assembly)      D. B and C
- B. main.o (object)        E. A, B and C
- C. main.exe (executable) F. something else

## exercise (2)

main.c:

```
1 #include <stdio.h>
2 void sayHello(void) {
3     puts("Hello, World!");
4 }
5 int main(void) {
6     sayHello();
7 }
```

Which files likely contain **literal ASCII string** of Hello, World!?

- A. main.s (assembly)      D. B and C
- B. main.o (object)        E. A, B and C
- C. main.exe (executable) F. something else

# main.s contains it?

```
.text
.global sayHello
sayHello:
    mov $.Lstr, %rdi
    call puts
    ...
.data
.Lstr:
    .string "Hello, World!"
```

---

```
.text
.global sayHello
sayHello:
    mov $.Lstr, %rdi
    call puts
    ...
.data
.Lstr:
    .byte 72,101,108,108,111,44,32,87,111,114,108,100,33,0
```

## main.o contains it?

complaint: in hexadecimal, like we've shown?

most object file formats aim for **efficiency**

simpler for linker to copy raw bytes

similar argument for main.exe and program loading

# dynamic linking (very briefly)

*dynamic linking* — done **when application is loaded**

idea: don't have  $N$  copies of `printf` on disk

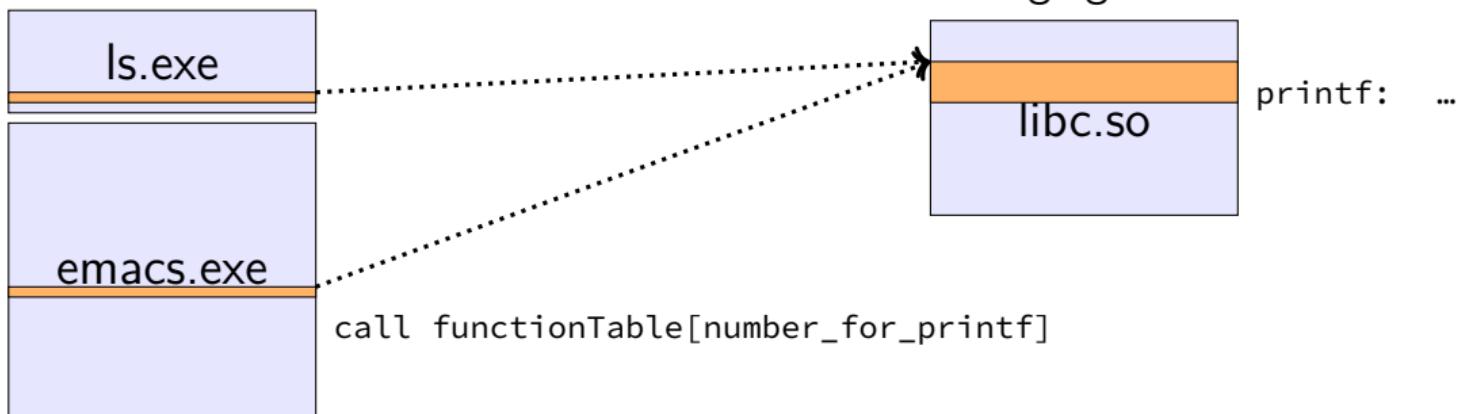
other type of linking: *static* (`gcc -static`)

load executable file + its libraries into memory when app starts

often extra indirection:

`call functionTable[number_for_printf]`

linker fills in `functionTable` instead of changing calls



# ldd /bin/ls

```
$ ldd /bin/ls
 linux-vdso.so.1 => (0x00007ffcca9d8000)
 libsdl.so.1 => /lib/x86_64-linux-gnu/libsdl.so.1
                  (0x00007f851756f000)
 libc.so.6 => /lib/x86_64-linux-gnu/libc.so.6
                  (0x00007f85171a5000)
 libpcre.so.3 => /lib/x86_64-linux-gnu/libpcre.so.3
                  (0x00007f8516f35000)
 libdl.so.2 => /lib/x86_64-linux-gnu/libdl.so.2
                  (0x00007f8516d31000)
 /lib64/ld-linux-x86-64.so.2 (0x00007f8517791000)
 libpthread.so.0 => /lib/x86_64-linux-gnu/libpthread.so.0
                  (0x00007f8516b14000)
```

## relocation types

machine code doesn't always use addresses as is

"call function 4303 bytes later"

linker needs to compute "4303"

extra field on relocation list

# C Data Types

Varies between machines(!). For this course:

type	size (bytes)
char	1
short	2
int	4
long	8

# C Data Types

Varies between machines(!). For **this course**:

type	size (bytes)
char	1
short	2
int	4
long	8
float	4
double	8

# C Data Types

Varies between machines(!). For **this course**:

type	size (bytes)
char	1
short	2
int	4
long	8
float	4
double	8
void *	8
<i>anything</i> *	8

truth

bool

# truth

bool

x == 4 is an int  
1 if true; 0 if false

## false values in C

0

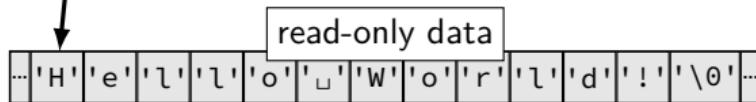
including null pointers — 0 cast to a pointer

# strings in C

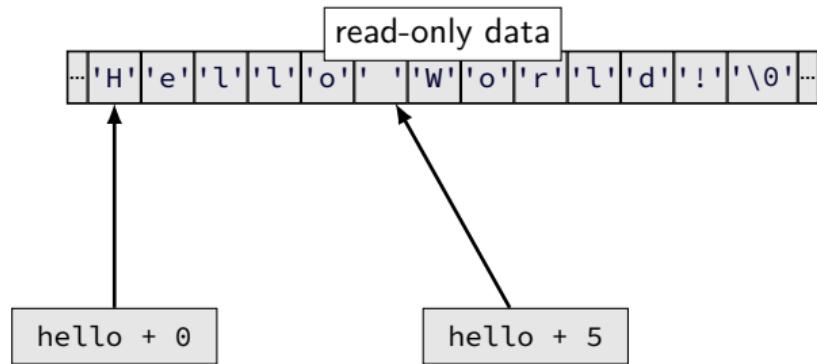
hello (on stack/register)

0x4005C0

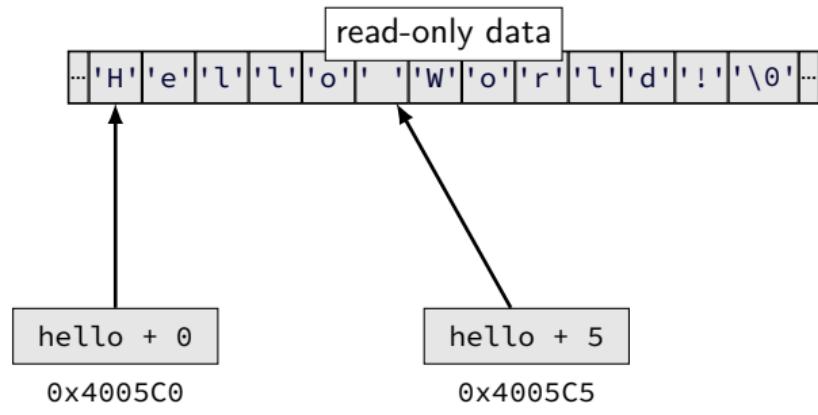
```
int main() {  
    const char *hello = "Hello World!";  
    ...  
}
```



# pointer arithmetic



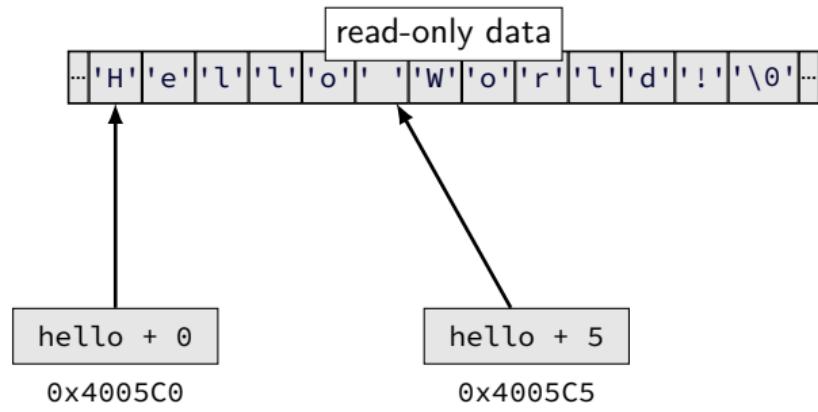
# pointer arithmetic



`*(hello + 0)` is 'H'

`*(hello + 5)` is ' '

# pointer arithmetic



`*(hello + 0)` is 'H'

`hello[0]` is 'H'

`*(hello + 5)` is ' '

`hello[5]` is ' '

# arrays and pointers

`*(foo + bar)` exactly the same as `foo[bar]`

arrays 'decay' into pointers

## arrays of non-bytes

array[2] and \*(array + 2) still the same

```
1 int numbers[4] = {10, 11, 12, 13};  
2 int *pointer;  
3 pointer = numbers;  
4 *pointer = 20; // numbers[0] = 20;  
5 pointer = pointer + 2;  
6 /* adds 8 (2 ints) to address */  
7 *pointer = 30; // numbers[2] = 30;  
8 // numbers is 20, 11, 30, 13
```

## arrays of non-bytes

array[2] and \*(array + 2) still the same

```
1 int numbers[4] = {10, 11, 12, 13};  
2 int *pointer;  
3 pointer = numbers;  
4 *pointer = 20; // numbers[0] = 20;  
5 pointer = pointer + 2;  
/* adds 8 (2 ints) to address */  
7 *pointer = 30; // numbers[2] = 30;  
8 // numbers is 20, 11, 30, 13
```

## exercise

```
1 char foo[4] = "foo";
2     // {'f', 'o', 'o', '\0'}
3 char *pointer;
4 pointer = foo;
5 *pointer = 'b';
6 pointer = pointer + 2;
7 pointer[0] = 'z';
8 *(foo + 1) = 'a';
```

Final value of foo?

- A. "fao"
- B. "zao"
- C. "baz"
- D. "bao"
- E. something else/crash

## exercise

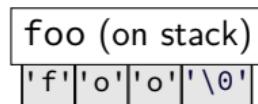
```
1 char foo[4] = "foo";
2     // {'f', 'o', 'o', '\0'}
3 char *pointer;
4 pointer = foo;
5 *pointer = 'b';
6 pointer = pointer + 2;
7 pointer[0] = 'z';
8 *(foo + 1) = 'a';
```

Final value of foo?

- A. "fao"
- B. "zao"
- C. "baz"
- D. "bao"
- E. something else/crash

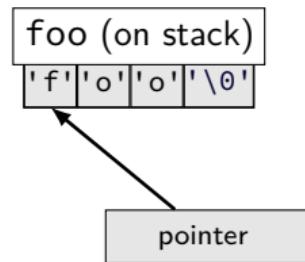
# exercise explanation

```
1 char foo[4] = "foo";
2     // {'f', 'o', 'o', '\0'}
3 char *pointer;
4 pointer = foo;
5 *pointer = 'b';
6 pointer = pointer + 2;
7 pointer[0] = 'z';
8 *(foo + 1) = 'a';
```



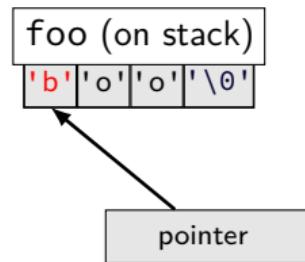
# exercise explanation

```
1 char foo[4] = "foo";
2 // {'f', 'o', 'o', '\0'}
3 char *pointer;
4 pointer = foo;
5 *pointer = 'b';
6 pointer = pointer + 2;
7 pointer[0] = 'z';
8 *(foo + 1) = 'a';
```



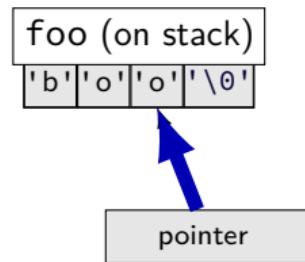
# exercise explanation

```
1 char foo[4] = "foo";
2     // {'f', 'o', 'o', '\0'}
3 char *pointer;
4 pointer = foo;
5 *pointer = 'b';
6 pointer = pointer + 2;
7 pointer[0] = 'z';
8 *(foo + 1) = 'a';
```



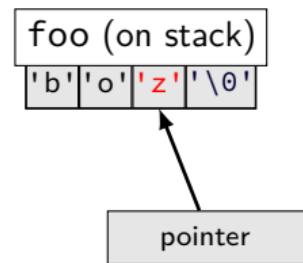
# exercise explanation

```
1 char foo[4] = "foo";
2     // {'f', 'o', 'o', '\0'}
3 char *pointer;
4 pointer = foo;
5 *pointer = 'b';
6 pointer = pointer + 2;
7 pointer[0] = 'z';
8 *(foo + 1) = 'a';
```



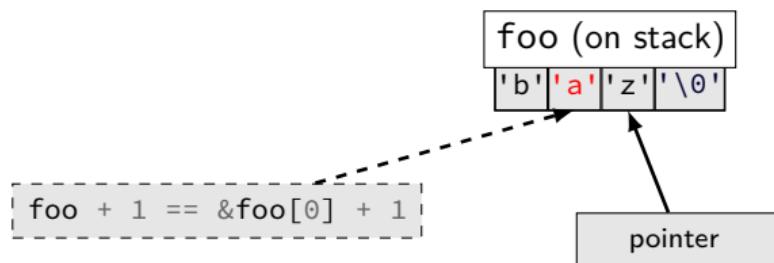
# exercise explanation

```
1 char foo[4] = "foo";
2     // {'f', 'o', 'o', '\0'}
3 char *pointer;
4 pointer = foo;
5 *pointer = 'b';
6 pointer = pointer + 2;
7 pointer[0] = 'z';    better style: *pointer = 'z';
8 *(foo + 1) = 'a';
```



# exercise explanation

```
1 char foo[4] = "foo";
2     // {'f', 'o', 'o', '\0'}
3 char *pointer;
4 pointer = foo;
5 *pointer = 'b';
6 pointer = pointer + 2;
7 pointer[0] = 'z';    better style: *pointer = 'z';
8 *(foo + 1) = 'a';    better style: foo[1] = 'a';
```



## arrays: not quite pointers (1)

```
int array[100];  
int *pointer;
```

Legal: pointer = array;  
same as pointer = &(array[0]);

## arrays: not quite pointers (1)

```
int array[100];  
int *pointer;
```

Legal: pointer = array;  
same as pointer = &(array[0]);

Illegal: ~~array = pointer;~~

## arrays: not quite pointers (2)

```
int array[100];  
int *pointer = array;
```

```
sizeof(array) == 400
```

size of all elements

## arrays: not quite pointers (2)

```
int array[100];  
int *pointer = array;
```

**sizeof(array) == 400**

size of all elements

**sizeof(pointer) == 8**

size of address

## arrays: not quite pointers (2)

```
int array[100];  
int *pointer = array;
```

**sizeof(array) == 400**

size of all elements

**sizeof(pointer) == 8**

size of address

**sizeof(&array[0]) == ???**

(&array[0] same as &(array[0]))

# struct

```
struct rational {  
    int numerator;  
    int denominator;  
};  
// ...  
struct rational two_and_a_half;  
two_and_a_half.numerator = 5;  
two_and_a_half.denominator = 2;  
struct rational *pointer = &two_and_a_half;  
printf("%d/%d\n",  
    pointer->numerator,  
    pointer->denominator);
```

# struct

```
struct rational {  
    int numerator;  
    int denominator;  
};  
// ...  
struct rational two_and_a_half;  
two_and_a_half.numerator = 5;  
two_and_a_half.denominator = 2;  
struct rational *pointer = &two_and_a_half;  
printf("%d/%d\n",  
      pointer->numerator,  
      pointer->denominator);
```

# typedef

instead of writing:

```
...  
unsigned int a;  
unsigned int b;  
unsigned int c;
```

can write:

```
typedef unsigned int uint;
```

```
...  
uint a;  
uint b;  
uint c;
```

# typedef struct (1)

```
struct other_name_for_rational {
    int numerator;
    int denominator;
};

typedef struct other_name_for_rational rational;
// ...
rational two_and_a_half;
two_and_a_half.numerator = 5;
two_and_a_half.denominator = 2;
rational *pointer = &two_and_a_half;
printf("%d/%d\n",
       pointer->numerator,
       pointer->denominator);
```

# typedef struct (1)

```
struct other_name_for_rational {
    int numerator;
    int denominator;
};

typedef struct other_name_for_rational rational;
// ...
rational two_and_a_half;
two_and_a_half.numerator = 5;
two_and_a_half.denominator = 2;
rational *pointer = &two_and_a_half;
printf("%d/%d\n",
       pointer->numerator,
       pointer->denominator);
```

```
typedef struct (2)
struct other_name_for_rational {
    int numerator;
    int denominator;
};
typedef struct other_name_for_rational rational;
// same as:
typedef struct other_name_for_rational {
    int numerator;
    int denominator;
} rational;
```

```
typedef struct (2)
struct other_name_for_rational {
    int numerator;
    int denominator;
};
typedef struct other_name_for_rational rational;
// same as:
typedef struct other_name_for_rational {
    int numerator;
    int denominator;
} rational;
```

## typedef struct (2)

```
struct other_name_for_rational {  
    int numerator;  
    int denominator;  
};  
typedef struct other_name_for_rational rational;  
// same as:  
typedef struct other_name_for_rational {  
    int numerator;  
    int denominator;  
} rational;  
// almost the same as:  
typedef struct {  
    int numerator;  
    int denominator;  
} rational;
```

## typedef struct (3)

```
struct other_name_for_rational {  
    int numerator;  
    int denominator;  
};  
typedef struct other_name_for_rational rational;
```

valid ways to declare an instance:

```
struct other_name_for_rational some_variable;  
rational some_variable;
```

INVALID ways:

```
/* INVALID: */ struct rational some_variable;  
/* INVALID: */ other_name_for_rational some_variable;
```

## typedef struct (3)

```
struct other_name_for_rational {  
    int numerator;  
    int denominator;  
};  
typedef struct other_name_for_rational rational;
```

valid ways to declare an instance:

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struct other_name_for_rational some_variable;  
rational some_variable;
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INVALID ways:

```
/* INVALID: */ struct rational some_variable;  
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```

## typedef struct (3)

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struct other_name_for_rational {  
    int numerator;  
    int denominator;  
};  
typedef struct other_name_for_rational rational;
```

valid ways to declare an instance:

```
struct other_name_for_rational some_variable;  
rational some_variable;
```

INVALID ways:

```
/* INVALID: */ struct rational some_variable;  
/* INVALID: */ other_name_for_rational some_variable;
```

# structs aren't references

```
typedef struct {  
    long a; long b; long c;  
} triple;  
...
```

```
triple foo;  
foo.a = foo.b = foo.c = 3;  
triple bar = foo;  
bar.a = 4;  
// foo is 3, 3, 3  
// bar is 4, 3, 3
```

...
return address
callee saved
registers
foo.c
foo.b
foo.a
bar.c
bar.b
bar.a

# C evolution and standards

1978: Kernighan and Ritchie publish *The C Programming Language* — “K&R C”

**very** different from modern C

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compiler option: -std=c99

adds: declare variables in middle of block

adds: **//** comments

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compiler option: -std=c99

adds: declare variables in middle of block

adds: // comments

2011, 2017: Second/Third ISO update — C11, C17

# undefined behavior example (1)

```
#include <stdio.h>
#include <limits.h>
int test(int number) {
    return (number + 1) > number;
}

int main(void) {
    printf("%d\n", test(INT_MAX));
}
```

# undefined behavior example (1)

```
#include <stdio.h>
#include <limits.h>
int test(int number) {
    return (number + 1) > number;
}

int main(void) {
    printf("%d\n", test(INT_MAX));
}
```

without optimizations: 0

# undefined behavior example (1)

```
#include <stdio.h>
#include <limits.h>
int test(int number) {
    return (number + 1) > number;
}

int main(void) {
    printf("%d\n", test(INT_MAX));
}
```

without optimizations: 0

with optimizations: 1

## undefined behavior example (2)

```
int test(int number) {  
    return (number + 1) > number;  
}
```

Optimized:

```
test:  
    movl    $1, %eax      # eax <- 1  
    ret
```

Less optimized:

```
test:  
    leal    1(%rdi), %eax # eax <- rdi + 1  
    cmpl    %eax, %edi  
    setl    %al            # al <- eax < edi  
    movzbl  %al, %eax     # eax <- al (pad with zeros)  
    ret
```

# undefined behavior

compilers can do **whatever they want**

what you expect

crash your program

...

common types:

*signed* integer overflow/underflow

out-of-bounds pointers

integer divide-by-zero

writing read-only data

out-of-bounds shift

# **undefined behavior**

why undefined behavior?

different architectures work differently

- allow compilers to expose whatever processor does “naturally”
- don’t encode any particular machine in the standard

flexibility for optimizations

# backup slides

# while exercise

```
while (b < 10) { foo(); b += 1; }
```

Assume b is in **callee-saved** register %rbx. Which are correct assembly translations?

*// version A*

```
start_loop:  
    call foo  
    addq $1, %rbx  
    cmpq $10, %rbx  
    jl start_loop
```

*// version B*

```
start_loop:  
    cmpq $10, %rbx  
    jge end_loop  
    call foo  
    addq $1, %rbx  
    jmp start_loop  
end_loop:
```

*// version C*

```
start_loop:  
    movq $10, %rax  
    subq %rbx, %rax  
    jge end_loop  
    call foo  
    addq $1, %rbx  
    jmp start_loop  
end_loop:
```

## while exercise: translating?

```
while (b < 10) {  
    foo();  
    b += 1;  
}
```

---

## while exercise: translating?

```
while (b < 10) {  
    foo();  
    b += 1;  
}
```

---

```
start_loop: if (b < 10) goto end_loop;  
            foo();  
            b += 1;  
            goto start_loop;  
end_loop:
```

# objdump -sx test.o (Linux) (1)

```
test.o:      file format elf64-x86-64
test.o
architecture: i386:x86-64, flags 0x00000011:
HAS_RELOC, HAS_SYMS
start address 0x0000000000000000
```

## Sections:

Idx	Name	Size	VMA	LMA	File off	Algn
0	.text	00000000	0000000000000000	0000000000000000	00000040	2**0
			CONTENTS, ALLOC, LOAD, READONLY, CODE			
1	.data	00000000	0000000000000000	0000000000000000	00000040	2**0
			CONTENTS, ALLOC, LOAD, DATA			
2	.bss	00000000	0000000000000000	0000000000000000	00000040	2**0
			ALLOC			
3	.rodata.str1.1	0000000e	0000000000000000	0000000000000000	00000040	2**0
			CONTENTS, ALLOC, LOAD, READONLY, DATA			
4	.text.startup	00000014	0000000000000000	0000000000000000	0000004e	2**0
			CONTENTS, ALLOC, LOAD, RELOC, READONLY, CODE			
5	.comment	0000002b	0000000000000000	0000000000000000	00000062	2**0
			CONTENTS, READONLY			
6	.note.GNU-stack	00000000	0000000000000000	0000000000000000	0000008d	2**0
			CONTENTS, READONLY			
7	.eh_frame	00000030	0000000000000000	0000000000000000	00000090	2**3
			CONTENTS, ALLOC, LOAD, RELOC, READONLY, DATA			

# objdump -sx test.o (Linux) (2)

## SYMBOL TABLE:

0000000000000000	l	df	*ABS*	0000000000000000	test.c
0000000000000000	l	d	.text	0000000000000000	.text
0000000000000000	l	d	.data	0000000000000000	.data
0000000000000000	l	d	.bss	0000000000000000	.bss
0000000000000000	l	d	.rodata.str1.1	0000000000000000	.rodata.str1.1
0000000000000000	l	d	.text.startup	0000000000000000	.text.startup
0000000000000000	l	d	.note.GNU-stack	0000000000000000	.note.GNU-stack
0000000000000000	l	d	.eh_frame	0000000000000000	.eh_frame
0000000000000000	l		.rodata.str1.1	0000000000000000	.LC0
0000000000000000	l	d	.comment	0000000000000000	.comment
0000000000000000	g	F	.text.startup	000000000000014	main
0000000000000000			*UND*	0000000000000000	_GLOBAL_OFFSET_TABLE_
0000000000000000			*UND*	0000000000000000	puts

## columns:

memory address (not yet assigned, so 0)

flags: l=local, g=global, F=function, ...

section (.text, .data, .bss, ...)

offset in section

name of symbol

# objdump -sx test.o (Linux) (3)

RELOCATION RECORDS FOR [.text.startup]:

OFFSET	TYPE	VALUE
0000000000000003	R_X86_64_PC32	.LC0—0x0000000000000004
000000000000000c	R_X86_64_PLT32	puts—0x0000000000000004

RELOCATION RECORDS FOR [.eh\_frame]:

OFFSET	TYPE	VALUE
0000000000000020	R_X86_64_PC32	.text.startup

Contents of section .rodata.str1.1:

0000 48656c6c 6f2c2057 6f726c64 2100	Hello, World!.
--------------------------------------	----------------

Contents of section .text.startup:

0000 488d3d00 00000048 83ec08e8 00000000	H.=....H.....
0010 31c05ac3	1.Z.

Contents of section .comment:

0000 00474343 3a202855 62756e74 7520372e	.GCC: (Ubuntu 7.
0010 332e302d 32377562 756e7475 317e3138	3.0—27ubuntu1~18
0020 2e303429 20372e33 2e3000	.04) 7.3.0.

Contents of section .eh\_frame:

0000 14000000 00000000 017a5200 01781001	.....zR..x..
0010 1b0c0708 90010000 14000000 1c000000	.....
0020 00000000 14000000 004b0e10 480e0800	.....K..H...

## example: C that is not C++

valid C and invalid C++:

```
char *str = malloc(100);
```

valid C and valid C++:

```
char *str = (char *) malloc(100);
```

valid C and invalid C++:

```
int class = 1;
```

# linked lists / dynamic allocation

```
typedef struct list_t {  
    int item;  
    struct list_t *next;  
} list;  
// ...
```

# linked lists / dynamic allocation

```
typedef struct list_t {  
    int item;  
    struct list_t *next;  
} list;  
// ...
```

# linked lists / dynamic allocation

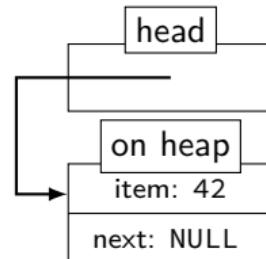
```
typedef struct list_t {  
    int item;  
    struct list_t *next;  
} list;  
// ...
```

```
list* head = malloc(sizeof(list));  
/* C++: new list; */  
head->item = 42;  
head->next = NULL;  
// ...  
free(head);  
/* C++: delete list */
```

# linked lists / dynamic allocation

```
typedef struct list_t {  
    int item;  
    struct list_t *next;  
} list;  
// ...
```

```
list* head = malloc(sizeof(list));  
/* C++: new list; */  
head->item = 42;  
head->next = NULL;  
// ...  
free(head);  
/* C++: delete list */
```

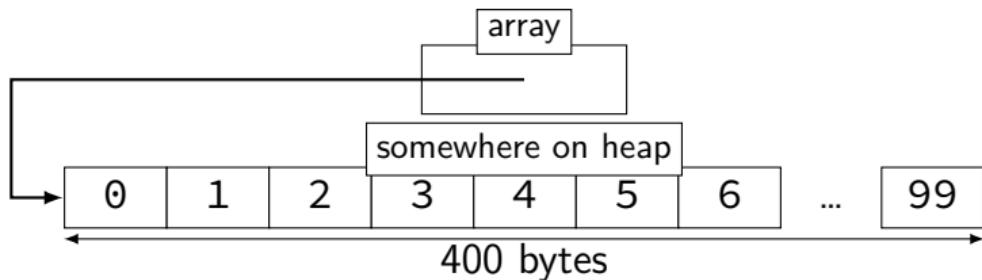


# dynamic arrays

```
int *array = malloc(sizeof(int)*100);
    // C++: new int[100]
for (i = 0; i < 100; ++i) {
    array[i] = i;
}
// ...
free(array); // C++: delete[] array
```

# dynamic arrays

```
int *array = malloc(sizeof(int)*100);
// C++: new int[100]
for (i = 0; i < 100; ++i) {
    array[i] = i;
}
// ...
free(array); // C++: delete[] array
```



## a note on precedence

`&foo[42]` is the same as `&(foo[42])` (*not* `(&foo)[42]`)

`*foo[42]` is the same as `*(foo[42])` (*not* `(*foo)[42]`)

`*foo++` is the same as `*(foo++)` (*not* `(*foo)++`)

# unsigned and signed types

type	min	max
<b>signed int = signed = int</b>	$-2^{31}$	$2^{31} - 1$
<b>unsigned int = unsigned</b>	0	$2^{32} - 1$
<b>signed long = long</b>	$-2^{63}$	$2^{63} - 1$
<b>unsigned long</b>	0	$2^{64} - 1$

:

## unsigned/signed comparison trap (1)

```
int x = -1;  
unsigned int y = 0;  
printf("%d\n", x < y);
```

## unsigned/signed comparison trap (1)

```
int x = -1;  
unsigned int y = 0;  
printf("%d\n", x < y);  
result is 0
```

# unsigned/signed comparison trap (1)

```
int x = -1;  
unsigned int y = 0;  
printf("%d\n", x < y);
```

result is 0

short solution: don't compare signed to unsigned:

```
(long) x < (long) y
```

## unsigned/sign comparison trap (2)

```
int x = -1;  
unsigned int y = 0;  
printf("%d\n", x < y);
```

compiler converts both to **same type** first

**int** if all possible values fit

otherwise: first operand (x, y) type from this list:

- unsigned long
- long
- unsigned int
- int

# stdio.h

C does not have <iostream>

instead <stdio.h>

# stdio

```
cr4bd@power1  
: /if22/cr4bd ; man stdio
```

...

STDIO(3)	Linux Programmer's Manual	STDIO(3)
----------	---------------------------	----------

## NAME

stdio - standard input/output library functions

## SYNOPSIS

```
#include <stdio.h>
```

```
FILE *stdin;  
FILE *stdout;  
FILE *stderr;
```

## DESCRIPTION

The standard I/O library provides a simple and efficient buffered stream I/O interface. Input and output is mapped into logical data streams and the physical I/O characteristics are concealed. The functions and macros are listed below; more information is available from the individual man pages.

# stdio

STDIO(3)

Linux Programmer's Manual

STDIO(3)

## NAME

stdio - standard input/output library functions

...

## List of functions

Function	Description
----------	-------------

---

clearerr	check and reset stream status
fclose	close a stream

...

printf	formatted output conversion
--------	-----------------------------

...

# printf

```
1 int custNo = 1000;
2 const char *name = "Jane Smith"
3     printf("Customer #%-d: %s\n" ,
4            custNo, name);
5 // "Customer #1000: Jane Smith"
6 // same as:
7 cout << "Customer #" << custNo
8      << ": " << name << endl;
```

# printf

```
1 int custNo = 1000;
2 const char *name = "Jane Smith"
3     printf("Customer # %d: %s\n" ,
4            custNo, name);
5 // "Customer #1000: Jane Smith"
6 // same as:
7 cout << "Customer #" << custNo
8      << ": " << name << endl;
```

# printf

```
1 int custNo = 1000;
2 const char *name = "Jane Smith"
3     printf("Customer #%-d: %s\n" ,
4             custNo, name);
5 // "Customer #1000: Jane Smith"
6 // same as:
7 cout << "Customer #" << custNo
8     << ": " << name << endl;
```

format string must **match types** of argument

# printf formats quick reference

Specifier	Argument Type	Example(s)
%s	char *	Hello, World!
%p	any pointer	0x4005d4
%d	int/short/char	42
%u	unsigned int/short/char	42
%x	unsigned int/short/char	2a
%ld	long	42
%f	double/float	42.000000 0.000000
%e	double/float	4.200000e+01 4.200000e-19
%g	double/float	42, 4.2e-19
%%	(no argument)	%

# printf formats quick reference

Specifier	Argument Type	Example(s)
%s	char *	Hello, World!
%p	any pointer	0x4005d4
%d	int/short/char	42
%u	unsigned int/short/char	42
%x	unsigned int/short/char	2a
%ld	detailed docs: man 3 printf	
%f	double/float	42.000000 0.000000
%e	double/float	4.200000e+01 4.200000e-19
%g	double/float	42, 4.2e-19
%%	(no argument)	%

# short-circuit (`||`)

```
1 #include <stdio.h>
2 int zero() { printf("zero()\n"); return 0; }
3 int one() { printf("one()\n"); return 1; }
4 int main() {
5     printf("> %d\n", zero() || one());
6     printf("> %d\n", one() || zero());
7     return 0;
8 }
```

zero()

one()

> 1

one()

> 1

	OR	false	true
false		false	true
true		true	true

# short-circuit (||)

```
1 #include <stdio.h>
2 int zero() { printf("zero()\n"); return 0; }
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zero()

one()

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one()

> 1

OR		false	true
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7     return 0;
8 }
```

`zero()`

`one()`

> 1

`one()`

> 1

OR		<code>false</code>	<code>true</code>
<code>false</code>		<code>false</code>	<code>true</code>
<code>true</code>		<code>true</code>	<code>true</code>

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8 }
```

`zero()`

`one()`

`> 1`

`one()`

`> 1`

OR	<b>false</b>	<b>true</b>
<b>false</b>	false	true
<b>true</b>	true	true

# short-circuit (`||`)

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5     printf("> %d\n", zero() || one());
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7     return 0;
8 }
```

zero()

one()

> 1

one()

> 1

	OR	false	true
false		false	true
true		true	true

# short-circuit (`&&`)

```
1 #include <stdio.h>
2 int zero() { printf("zero()\n"); return 0; }
3 int one() { printf("one()\n"); return 1; }
4 int main() {
5     printf("> %d\n", zero() && one());
6     printf("> %d\n", one() && zero());
7     return 0;
8 }
```

`zero()`

`> 0`

`one()`

`zero()`

`> 0`

AND	<b>false</b>	<b>true</b>
<b>false</b>	false	false
<b>true</b>	false	true

# short-circuit (`&&`)

```
1 #include <stdio.h>
2 int zero() { printf("zero()\n"); return 0; }
3 int one() { printf("one()\n"); return 1; }
4 int main() {
5     printf("> %d\n", zero() && one());
6     printf("> %d\n", one() && zero());
7     return 0;
8 }
```

`zero()`

`> 0`

`one()`

`zero()`

`> 0`

AND		<b>false</b>	<b>true</b>
<b>false</b>		false	false
<b>true</b>		false	true

# short-circuit (`&&`)

```
1 #include <stdio.h>
2 int zero() { printf("zero()\n"); return 0; }
3 int one() { printf("one()\n"); return 1; }
4 int main() {
5     printf("> %d\n", zero() && one());
6     printf("> %d\n", one() && zero());
7     return 0;
8 }
```

`zero()`

`> 0`

`one()`

`zero()`

`> 0`

	AND	<b>false</b>	<b>true</b>
<b>false</b>		false	false
<b>true</b>		false	true

# short-circuit (`&&`)

```
1 #include <stdio.h>
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7     return 0;
8 }
```

`zero()`

`> 0`

`one()`

`zero()`

`> 0`

AND	<b>false</b>	<b>true</b>
<b>false</b>	false	false
<b>true</b>	false	true

# short-circuit (`&&`)

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1 #include <stdio.h>
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8 }
```

zero()

> 0

one()

zero()

> 0

AND	false	true
false	false	false
true	false	true