Assembly part 2 / C part 1

last time

linking extras:

- different kinds of relocations — addresses versus offset to addresses
- dynamic linking (briefly)

AT&T syntax

- destination last
- `O(B, I, S)` — $B + I \times S + O$
- `jmp *`

`lea` — `mov`, but don't do memory access

if/do-while to assembly

condition codes — last arithmetic result

reminder: quiz

post-quiz — after this lecture

pre-quiz — before next lecture

demo

condition codes

x86 has condition codes

set by (almost) all arithmetic instructions

- addq, subq, imulq, etc.

store info about last arithmetic result

- was it zero? was it negative? etc.

condition codes and jumps

`jg`, `jle`, etc. read condition codes

named based on interpreting result of subtraction

0: equal; negative: less than; positive: greater than

condition codes example (1)

```
movq $-10, %rax
movq $20, %rbx
subq %rax, %rbx // %rbx - %rax = 30
  // result > 0: %rbx was > %rax
jle foo // not taken; 30 > 0
```

condition codes and cmpq

"last arithmetic result"???

then what is cmp, etc.?

`cmp` does subtraction (but doesn't store result)

similar `test` does bitwise-and

`testq %rax, %rax` — result is %rax

condition codes example (2)

```
movq $-10, %rax
movq $20, %rbx
cmpq %rax, %rbx
jle foo // not taken; %rbx - %rax > 0
```

omitting the cmp

```
    movq $99, %r12 // register for x
start_loop:
    call foo
    subq $1, %r12
    cmpq $0, %r12
    // compute r12 - 0 + sets cond. codes
    jge  start_loop  // r12 >= 0?
                     // or result >= 0?
```

```
    movq $99, %r12 // register for x
start_loop:
    call foo
    subq $1, %r12
    // new r12 = old r12 - 1 + sets cond. codes
    jge  start_loop  // old r12 >= 1?
                     // or result >= 0?
```

condition codes example (3)

```
movq $-10, %rax
movq $20, %rbx
subq %rax, %rbx
jle  foo // not taken, %rbx - %rax > 0 -> %rbx

movq $20, %rbx
addq $-20, %rbx
je   foo // taken, result is 0
         // x - y = 0 -> x = y
```

what sets condition codes

most instructions that compute something set condition codes

some instructions only set condition codes:

- `cmp` ∼ `sub`
- `test` ∼ `and` (bitwise and — later)
- `testq %rax, %rax` — result is %rax

some instructions don't change condition codes:

- `lea`, `mov`
- control flow: `jmp`, `call`, `ret`, `jle`, etc.

condition codes examples (4)

```
movq $20, %rbx
addq $-20, %rbx // result is 0
movq $1, %rax // irrelevant
je   foo // taken, result is 0
```

while-to-assembly (1)

```
while (x >= 0) {
    foo()
    x--;
}
```

---

while-to-assembly (1)

```
while (x >= 0) {
    foo()
    x--;
}
```

```
start_loop:
    if (x < 0) goto end_loop;
    foo()
    x--;
    goto start_loop:
end_loop:
```

while-to-assembly (2)

```
start_loop:
    if (x < 0) goto end_loop;
    foo()
    x--;
    goto start_loop:
end_loop:
```

```
start_loop:
    cmpq $0, %r12
    jl end_loop // jump if r12 - 0 >= 0
    call foo
    subq $1, %r12
    jmp start_loop
```

while exercise

```
while (b < 10) { foo(); b += 1; }
```

Assume b is in callee-saved register %rbx. Which are correct assembly translations?

```
// version A
start_loop:
  call foo
  addq $1, %rbx
  cmpq $10, %rbx
  jl start_loop
```

```
// version B
start_loop:
  cmpq $10, %rbx
  jge end_loop
  call foo
  addq $1, %rbx
  jmp start_loop
end_loop:
```

```
// version C
start_loop:
  movq $10, %rax
  subq %rbx, %rax
  jge end_loop
  call foo
  addq $1, %rbx
  jmp start_loop
end_loop:
```

while to assembly (1)

```
while (b < 10) {
    foo();
    b += 1;
}
```

---

while to assembly (1)

```
while (b < 10) {
    foo();
    b += 1;
}
```

```
start_loop: if (b < 10) goto end_loop;
            foo();
            b += 1;
            goto start_loop;
end_loop:
```

while — levels of optimization

```
while (b < 10) { foo(); b += 1; }
```

```
start_loop:
  cmpq $10, %rbx
  jge end_loop
  call foo
  addq $1, %rbx
  jmp start_loop
end_loop:
    ...
    ...
    ...
    ...
```

while — levels of optimization

```
while (b < 10) { foo(); b += 1; }
```

```
start_loop:
  cmpq $10, %rbx
  jge end_loop
  call foo
  addq $1, %rbx
  jmp start_loop
end_loop:
    ...
    ...
    ...
    ...
```

```
  cmpq $10, %rbx
  jge end_loop
start_loop:
  call foo
  addq $1, %rbx
  cmpq $10, %rbx
  jne start_loop
end_loop:
    ...
    ...
    ...
```

while — levels of optimization

```
while (b < 10) { foo(); b += 1; }
```

```
start_loop:
  cmpq $10, %rbx
  jge end_loop
  call foo
  addq $1, %rbx
  jmp start_loop
end_loop:
    ...
    ...
    ...
    ...
```

```
  cmpq $10, %rbx
  jge end_loop
start_loop:
  call foo
  addq $1, %rbx
  cmpq $10, %rbx
  jne start_loop
end_loop:
    ...
    ...
    ...
```

```
  cmpq $10, %rbx
  jge end_loop
  movq $10, %rax
  subq %rbx, %rax
  movq %rax, %rbx
start_loop:
  call foo
  decq %rbx
  jne start_loop
  movq $10, %rbx
end_loop:
```

condition codes: closer look

x86 condition codes:

`ZF` ("zero flag") — was result zero? (sub/cmp: equal)
`SF` ("sign flag") — was result negative? (sub/cmp: less)

(and some more, e.g. to handle overflow)

GDB: part of "eflags" register

set by cmp, test, arithmetic

condition codes example (2)

```
movq $-10, %rax
movq $20, %rbx
cmpq %rax, %rbx
jle foo // not taken; %rbx - %rax > 0
```

%rbx - %rax = 30 — SF = 0 (not negative), ZF = 0 (not zero)

condition codes examples (4)

```
movq $20, %rbx
addq $-20, %rbx // result is 0
movq $1, %rax // irrelevant
je   foo // taken, result is 0
```

20 + -20 = 0 — SF = 0 (not negative), ZF = 1 (zero)

condition codes: closer look

x86 condition codes:

`ZF` ("zero flag") — was result zero? (sub/cmp: equal)
`SF` ("sign flag") — was result negative? (sub/cmp: less)
`CF` ("carry flag") — did computation overflow (as unsigned)?
`OF` ("overflow flag") — did computation overflow (as signed)?
(and one more)

GDB: part of "eflags" register

set by cmp, test, arithmetic

closer look: condition codes (1)

```
  movq $-10, %rax
  movq $20, %rbx
  cmpq %rax, %rbx
// result = %rbx - %rax = 30
```

as signed: $20 - (-10) = 30$

as unsigned: $20 - (2^{64} - 10) = \cancel{-2^{64} - 30}\ 30$ (overflow!)

$\texttt{ZF} = 0$ (false) not zero rax and rbx not equal

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ZF $= 0$ (false)	not zero	rax and rbx not equal
SF $= 0$ (false)	not negative	rax $<=$ rbx

closer look: condition codes (1)

```
  movq $-10, %rax
  movq $20, %rbx
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as signed: $20 - (-10) = 30$

as unsigned: $20 - (2^{64} - 10) = \cancel{-2^{64} - 30}\ 30$ (overflow!)

ZF $= 0$ (false)	not zero	rax and rbx not equal
SF $= 0$ (false)	not negative	rax <= rbx
OF $= 0$ (false)	no overflow as signed	correct for signed

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```
  movq $-10, %rax
  movq $20, %rbx
  cmpq %rax, %rbx
// result = %rbx - %rax = 30
```

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ZF $= 0$ (false)	not zero	rax and rbx not equal
SF $= 0$ (false)	not negative	rax <= rbx
OF $= 0$ (false)	no overflow as signed	correct for signed
CF $= 1$ (true)	overflow as unsigned	incorrect for unsigned

exercise: condition codes (2)

```
  // 2^63 - 1
  movq $0x7FFFFFFFFFFFFFFF, %rax
  // 2^63 (unsigned); -2**63 (signed)
  movq $0x8000000000000000, %rbx
  cmpq %rax, %rbx
// result = %rbx - %rax
```

ZF = ?

SF = ?

OF = ?

CF = ?

closer look: condition codes (2)

```
  // 2**63 - 1
  movq $0x7FFFFFFFFFFFFFFF, %rax
  // 2**63 (unsigned); -2**63 (signed)
  movq $0x8000000000000000, %rbx
  cmpq %rax, %rbx
// result = %rbx - %rax
```

as signed: $-2^{63} - \left(2^{63} - 1\right) = \cancel{-2^{64} + 1}\ 1$ (overflow)

as unsigned: $2^{63} - \left(2^{63} - 1\right) = 1$

$\mathrm{ZF} = 0$ (false) not zero rax and rbx not equal

closer look: condition codes (2)

```
  // 2**63 - 1
  movq $0x7FFFFFFFFFFFFFFF, %rax
  // 2**63 (unsigned); -2**63 (signed)
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  cmpq %rax, %rbx
// result = %rbx - %rax
```

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as unsigned: $2^{63} - \left(2^{63} - 1\right) = 1$

ZF $= 0$ (false)	not zero	rax and rbx not equal
SF $= 0$ (false)	not negative	rax $<=$ rbx (if correct)

closer look: condition codes (2)

```
  // 2**63 - 1
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as unsigned: $2^{63} - \left(2^{63} - 1\right) = 1$

ZF $= 0$ (false)	not zero	rax and rbx not equal
SF $= 0$ (false)	not negative	rax <= rbx (if correct)
OF $= 1$ (true)	overflow as signed	incorrect for signed

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  // 2**63 - 1
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```

as signed: $-2^{63} - \left(2^{63} - 1\right) = \cancel{-2^{64} + 1}\ 1$ (overflow)

as unsigned: $2^{63} - \left(2^{63} - 1\right) = 1$

ZF $= 0$ (false)	not zero	rax and rbx not equal
SF $= 0$ (false)	not negative	rax <= rbx (if correct)
OF $= 1$ (true)	overflow as signed	incorrect for signed
CF $= 0$ (false)	no overflow as unsigned	correct for unsigned

closer look: condition codes (3)

```
  movq  $-1, %rax
  addq  $-2, %rax
// result = -3
```

as signed: $-1 + (-2) = -3$

as unsigned: $(2^{64} - 1) + (2^{64} - 2) = \cancel{2^{65} - 3}\ 2^{64} - 3$ (overflow)

ZF $= 0$ (false) not zero result not zero

closer look: condition codes (3)

```
  movq  $-1, %rax
  addq  $-2, %rax
// result = -3
```

as signed: $-1 + (-2) = -3$

as unsigned: $(2^{64} - 1) + (2^{64} - 2) = \cancel{2^{65} - 3}$ $2^{64} - 3$ (overflow)

ZF $= 0$ (false)	not zero	result not zero
SF $= 1$ (true)	negative	result is negative
OF $= 0$ (false)	no overflow as signed	correct for signed
CF $= 1$ (true)	overflow as unsigned	incorrect for unsigned

C Data Types

Varies between machines(!). For this course:

type	size (bytes)
`char`	1
`short`	2
`int`	4
`long`	8

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C Data Types

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type	size (bytes)
`char`	1
`short`	2
`int`	4
`long`	8
`float`	4
`double`	8
`void *`	8
`anything` `*`	8

truth

~~bool~~

truth

~~`bool`~~

`x == 4` is an `int`

`1` if true; `0` if false

false values in C

`0`

including null pointers — `0` cast to a pointer

short-circuit (||)

```
1 #include <stdio.h>
2 int zero() { printf("zero()\n"); return 0; }
3 int one() { printf("one()\n"); return 1; }
4 int main() {
5     printf("> %d\n", zero() || one());
6     printf("> %d\n", one() || zero());
7     return 0;
8 }
```

```
zero()
one()
> 1
one()
> 1
```

OR	**false**	**true**
false	false	true
true	true	true

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zero()
one()
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int zero() { printf("zero()\n"); return 0; }
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int main() {
    printf("> %d\n", zero() || one());
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    return 0;
}
```

```
zero()
one()
> 1
one()
> 1
```

OR	**false**	**true**
false	false	true
true	true	true

short-circuit (&&)

```c
#include <stdio.h>
int zero() { printf("zero()\n"); return 0; }
int one() { printf("one()\n"); return 1; }
int main() {
    printf("> %d\n", zero() && one());
    printf("> %d\n", one() && zero());
    return 0;
}
```

```
zero()
> 0
one()
zero()
> 0
```

AND	**false**	**true**
false	false	false
true	false	true

short-circuit (&&)

```
1 #include <stdio.h>
2 int zero() { printf("zero()\n"); return 0; }
3 int one() { printf("one()\n"); return 1; }
4 int main() {
5     printf("> %d\n", zero() && one());
6     printf("> %d\n", one() && zero());
7     return 0;
8 }
```

```
zero()
> 0
one()
zero()
> 0
```

AND	**false**	**true**
false	false	false
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short-circuit (&&)

```
1 #include <stdio.h>
2 int zero() { printf("zero()\n"); return 0; }
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7     return 0;
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```

```
zero()
> 0
one()
zero()
> 0
```

AND	**false**	**true**
false	false	false
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AND	**false**	**true**
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zero()
> 0
one()
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> 0
```

AND	**false**	**true**
false	false	false
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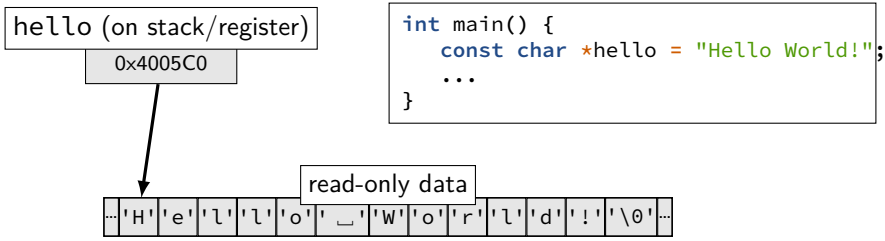
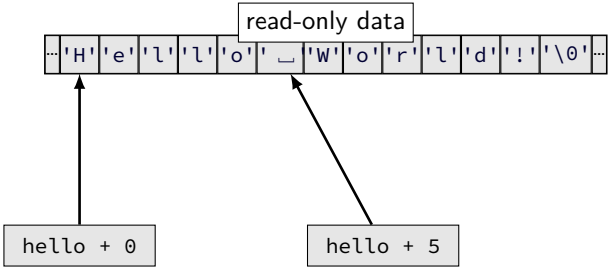
strings in C

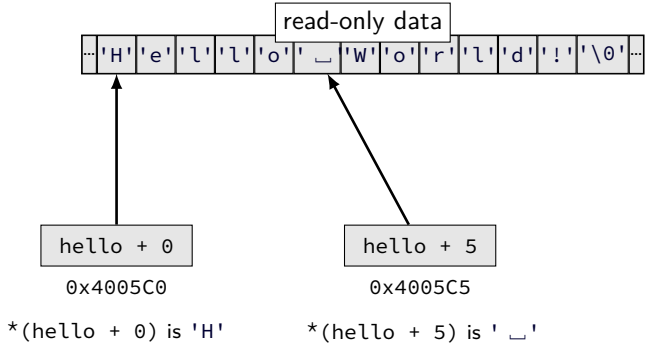



```c
int main() {
   const char *hello = "Hello World!";
   ...
}
```

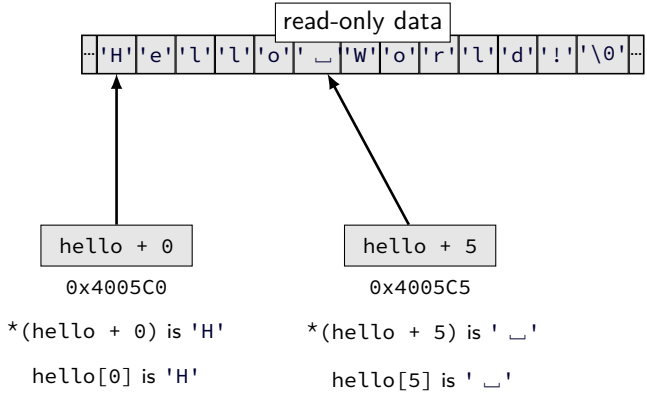
pointer arithmetic



pointer arithmetic



pointer arithmetic



arrays and pointers

`*(foo + bar)` exactly the same as `foo[bar]`

arrays ‘decay’ into pointers

arrays of non-bytes

`array[2]` and `*(array + 2)` still the same

```
1 int numbers[4] = {10, 11, 12, 13};
2 int *pointer;
3 pointer = numbers;
4 *pointer = 20;  // numbers[0] = 20;
5 pointer = pointer + 2;
6 /* adds 8 (2 ints) to address */
7 *pointer = 30;  // numbers[2] = 30;
8 // numbers is 20, 11, 30, 13
```

arrays of non-bytes

`array[2]` and `*(array + 2)` still the same

```
1  int numbers[4] = {10, 11, 12, 13};
2  int *pointer;
3  pointer = numbers;
4  *pointer = 20;  // numbers[0] = 20;
5  pointer = pointer + 2;
6  /* adds 8 (2 ints) to address */
7  *pointer = 30;  // numbers[2] = 30;
8  // numbers is 20, 11, 30, 13
```

a note on precedence

`&foo[1]` is the same as `&(foo[1])` (*not* `(&foo)[1]`)

`*foo[0]` is the same as `*(foo[0])` (*not* `(*foo)[0]`)

`*foo++` is the same as `*(foo++)` (*not* `(*foo)++`)

exercise

```
1 char foo[4] = "foo";
2     // {'f', 'o', 'o', '\0'}
3 char *pointer;
4 pointer = foo;
5 *pointer = 'b';
6 pointer = pointer + 2;
7 pointer[0] = 'z';
8 *(foo + 1) = 'a';
```

Final value of `foo`?

A. `"fao"`

B. `"zao"`

C. `"baz"`

D. `"bao"`

E. something else/crash

exercise

```
1 char foo[4] = "foo";
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exercise explanation

```
1  char foo[4] = "foo";
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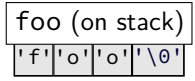

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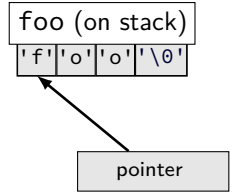
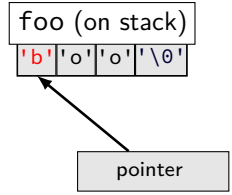
6  pointer = pointer + 2;
7  pointer[0] = 'z';
8  *(foo + 1) = 'a';
```



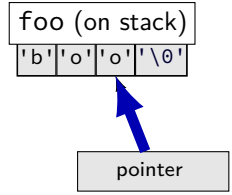
exercise explanation

```
1  char foo[4] = "foo";
2      // {'f', 'o', 'o', '\0'}
3  char *pointer;
4  pointer = foo;
5  *pointer = 'b';
6  pointer = pointer + 2;
7  pointer[0] = 'z';
8  *(foo + 1) = 'a';
```



exercise explanation

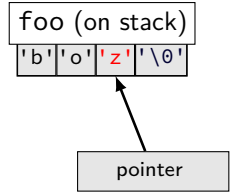
```
1  char foo[4] = "foo";
2      // {'f', 'o', 'o', '\0'}
3  char *pointer;
4  pointer = foo;
5  *pointer = 'b';
6  pointer = pointer + 2;
7  pointer[0] = 'z';
8  *(foo + 1) = 'a';
```



exercise explanation

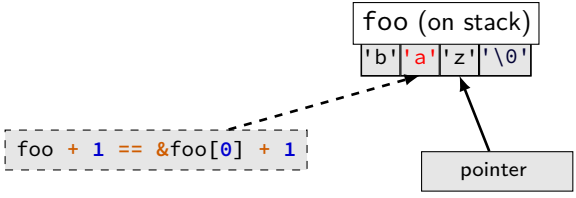
```
1  char foo[4] = "foo";
2      // {'f', 'o', 'o', '\0'}
3  char *pointer;
4  pointer = foo;
5  *pointer = 'b';
6  pointer = pointer + 2;
7  pointer[0] = 'z';
8  *(foo + 1) = 'a';
```

better style: `*pointer = 'z';`



exercise explanation

```
1  char foo[4] = "foo";
2      // {'f', 'o', 'o', '\0'}
3  char *pointer;
4  pointer = foo;
5  *pointer = 'b';
6  pointer = pointer + 2;
7  pointer[0] = 'z';      better style: *pointer = 'z';
8  *(foo + 1) = 'a';      better style: foo[1] = 'a';
```



backup slides

example: C that is not C++

valid C and invalid C++:

```
char *str = malloc(100);
```

valid C and valid C++:

```
char *str = (char *) malloc(100);
```

valid C and invalid C++:

```
int class = 1;
```