

heap vulnerabilities 2 / bounds checking

last time

ASLR details

- libraries/executables stay together

- cost of position-independent code, esp. on 32-bit x86 (little relative addressing)

- Windows choice to edit code for relocations

sudo exploit — defeating ASLR by only changing low bits of pointer

“heap smashing”

- pointer subterfuge using pointers used internally by malloc/free

use-after-free vulnerabilities (started)

vulnerable code

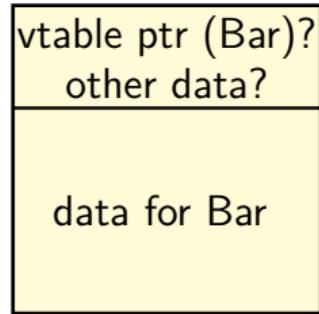
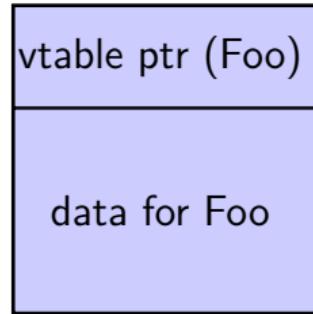
```
class Foo {  
    ...  
};  
Foo *the_foo;  
the_foo = new Foo;  
...  
delete the_foo;  
...  
something_else = new Bar(...);  
the_foo->something();
```

something_else likely where the_foo was

vulnerable code

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exploiting use after-free

trigger many “bogus” frees; then

allocate many things of same size with “right” pattern
pointers to shellcode?

pointers to pointers to system()?

objects with something useful in VTable entry?

trigger use-after-free thing

exercise

vuln. code

```
std::istream *in =  
    new std::ifstream("in.txt");  
...  
delete in;  
...  
char *other_buffer =  
    new char[strlen(INPUT) + 1];  
strcpy(other_buffer, INPUT);  
...  
char c = in->get();
```

ifstream internals

```
class istream {  
    ...  
    int get() { ... buf->uflow(); ... }  
    streambuf *buf;  
    ~istream() { delete buf; }  
};  
class streambuf {  
    ...  
protected:  
    virtual type_for_char uflow() = 0;  
        /* called to get next char */  
};  
class _File_streambuf : public streambuf { ... }
```

attacker goal: change what uflow() call does

Q1: assuming same size → likely to get same address, what size for attacker to choose for INPUT?

Q2: where in INPUT to place pointer to code to run?

real UAF exploitable bug

2012 bug in Google Chrome

exploitable via JavaScript

discovered/proof of concept by PinkiePie

allowed arbitrary code execution via VTable manipulation

UAF triggering code

```
// in HTML near this JavaScript:  
// <video id="vid"> (video player element)  
function source_opened() {  
    buffer = ms.addSourceBuffer('video/webm; codecs="vorbis,vp8"');  
    vid.parentNode.removeChild(vid);  
    gc(); // force garbage collector to run now  
    // garbage collector frees unreachable objects  
    // (would be run automatically, eventually, too)  
    // buffer now internally refers to delete'd player object  
    buffer.timestampOffset = 42;  
}  
ms = new WebKitMediaSource();  
ms.addEventListener('webkitsourceopen', source_opened);  
vid.src = window.URL.createObjectURL(ms);
```

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```

UAF triggering code

```
// implements JavaScript buffer.timestampOffset = 42
// ...
void SourceBuffer::setTimestampOffset(...) {
    if (m_source->setTimestampOffset(...))
        ...
}
bool MediaSource::setTimestampOffset(...) {
    // m_player was deleted when video player element deleted
    // but this call does *not* use a VTable
    if (!m_player->sourceSetTimestampOffset(id, offset))
        ...
}
bool MediaPlayer::sourceSetTimestampOffset(...) {
    // m_private deleted when MediaPlayer deleted
    // this *is* a VTable-based call
    return m_private->sourceSetTimestampOffset(id, offset);
}
```

UAF triggering code

```
// implements JavaScript buffer.timestampOffset = 42
// ...
void SourceBuffer::setTimestampOffset(...) {
    if (m_source->setTimestampOffset(...))
        ...
}
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bool MediaPlayer::sourceSetTimestampOffset(...) {
    // m_private deleted when MediaPlayer deleted
    // this *is* a VTable-based call
    return m_private->sourceSetTimestampOffset(id, offset);
}
```

UAF exploit (approx. pseudocode)

```
... /* use information leaks to find relevant addresses */
buffer = ms.addSourceBuffer('video/webm; codecs="vorbis,vp8"');
vid.parentNode.removeChild(vid);
vid = null;
gc();
// allocate object to replace m_private
var array = new Uint32Array(168/4);
// allocate object to replace m_player
// type chosen to keep m_private pointer unchanged
rtc = new webkitRTCPeerConnection({'iceServers': []});
array[0] = ... /* fill in array with chosen values */
// trigger VTable Call that uses chosen address
buffer.timestampOffset = 42;
```

type confusion

MediaPlayer (deleted but used)

m_private (pointer to PlayerImpl)
m_timestampOffset (double)

PlayerImpl (deleted but used)

VTable pointer
...

webkitRTC... (replacement)

(something not changed)
m_??? (pointer)
...

array of 32-bit ints (replacement)

array[0], array[1]
array[2], array[3]
...

missing pieces: information disclosure

need to learn address to set VTable pointer to
(and other addresses to use)

allocate types other than Uint32Array

rely on confusing between different types, e.g.

MediaPlayer (deleted but used)

<code>m_private</code> (pointer to PlayerImpl)
<code>m_timestampOffset</code> (double)

Something (replacement)

...
<code>m_buffer</code> (pointer)

allows reading timestamp value to get a pointer's address

use-after-free easy cases

common problem for JavaScript implementations

use-after-free'd object often some complex C++ object

example: representation of video stream

exploits can **choose type of object that replaces**

allocate that kind of object in JS

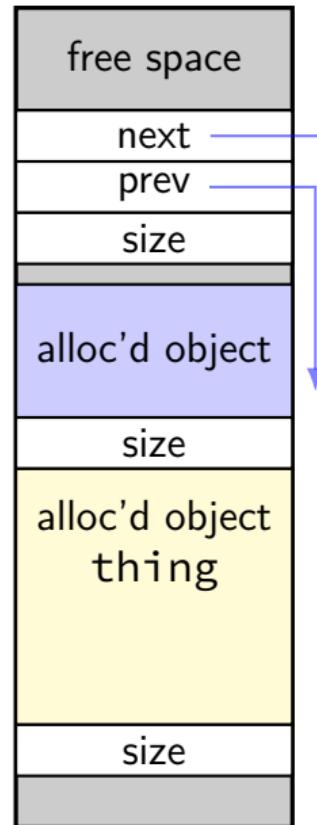
can often arrange to read/write vtable pointer

depends on layout of thing created

easy examples: string, array of floating point numbers

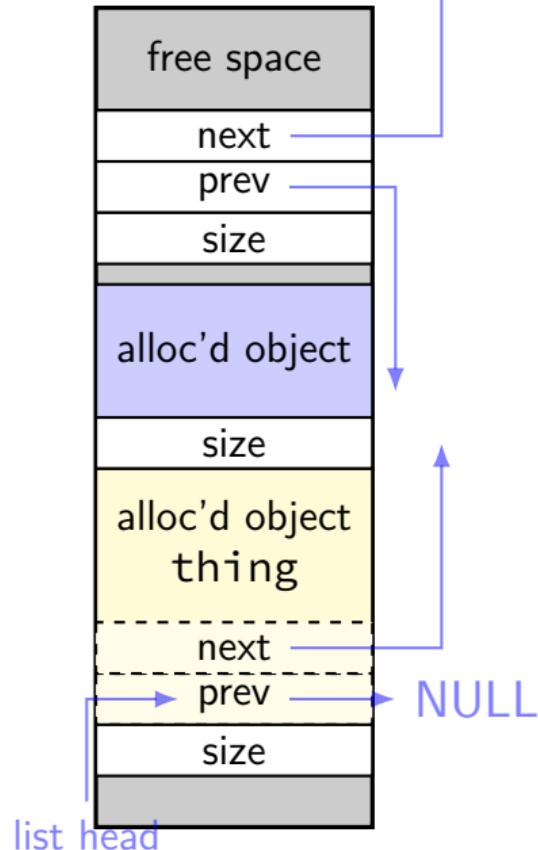
double-frees

```
free(thing);  
free(thing);  
char *p = malloc(...);  
// p points to next/prev  
// on list of avail.  
// blocks  
strcpy(p, attacker_controlled);  
malloc(...);  
char *q = malloc(...);  
// q points to attacker-  
// chosen address  
strcpy(q, attacker_controlled2);  
...
```



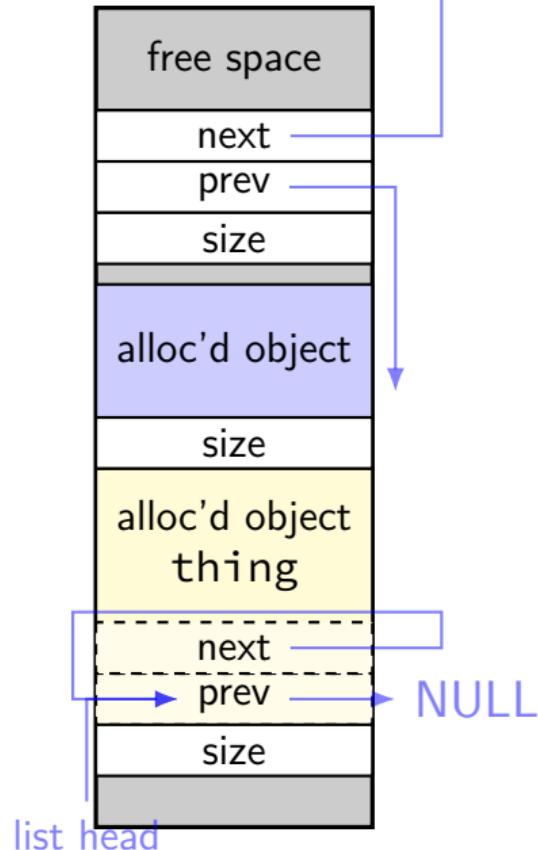
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// chosen address  
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...
```

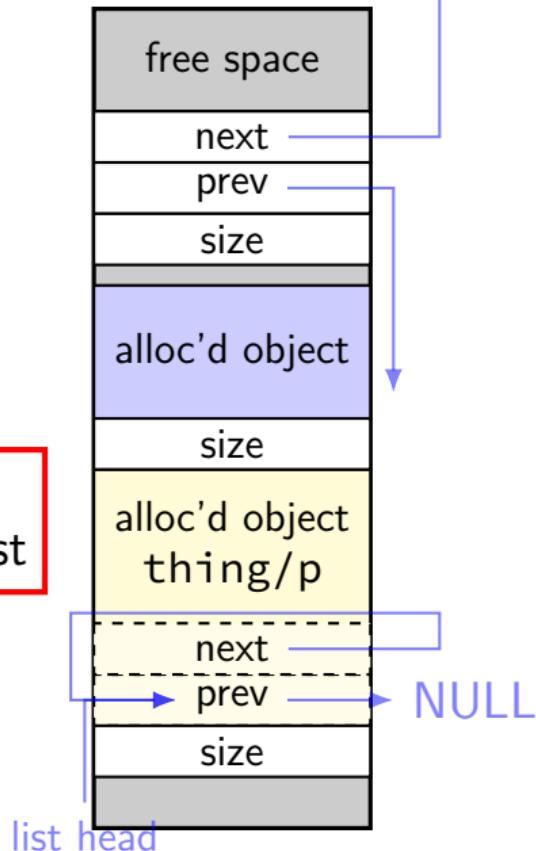


double-frees

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free(thing);  
free(thing);  
char *p = malloc(...);  
// p points to next/prev  
// on list of avail.  
// blocks
```

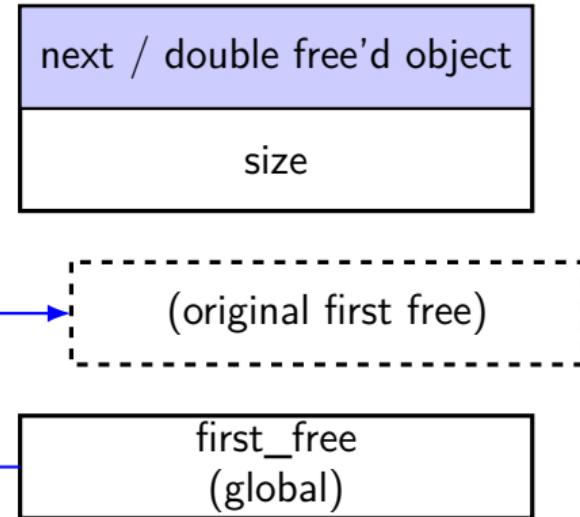
malloc returns something **still on free list**
because double-free made **loop** in linked list

```
// q points to attacker-  
// chosen address  
strcpy(q, attacker_controlled2);  
...
```



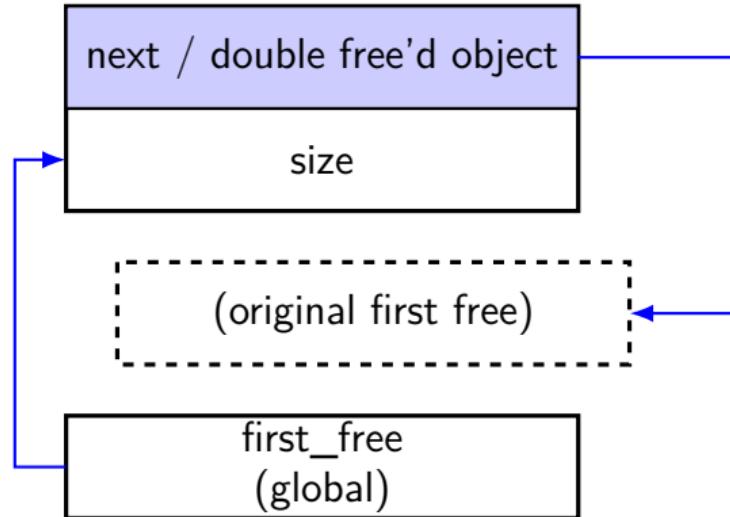
double-free expansion

```
// free/delete 1:  
double_freed->next = first_free;  
first_free = chunk;  
// free/delete 2:  
double_freed->next = first_free;  
first_free = chunk  
// malloc/new 1:  
result1 = first_free;  
first_free = first_free->next;  
// + overwrite:  
strcpy(result1, ...);  
// malloc/new 2:  
first_free = first_free->next;  
// malloc/new 3:  
result3 = first_free;  
strcpy(result3, ...);
```



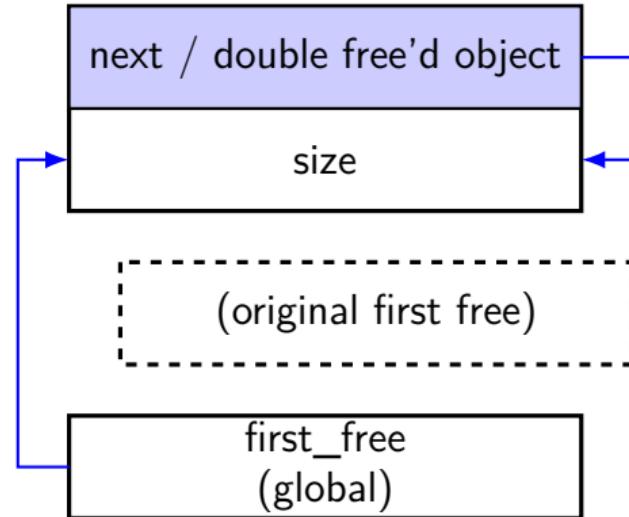
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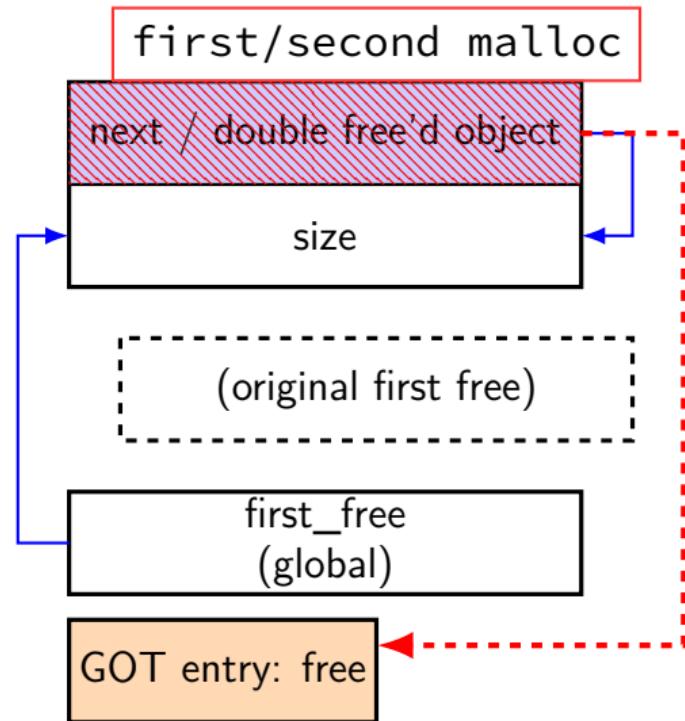
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result1 = first_free;  
first_free = first_free->next;  
// + overwrite:  
strcpy(result1, ...);  
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first_free = first_free->next;  
// malloc/new 3:  
result3 = first_free;  
strcpy(result3, ...);
```



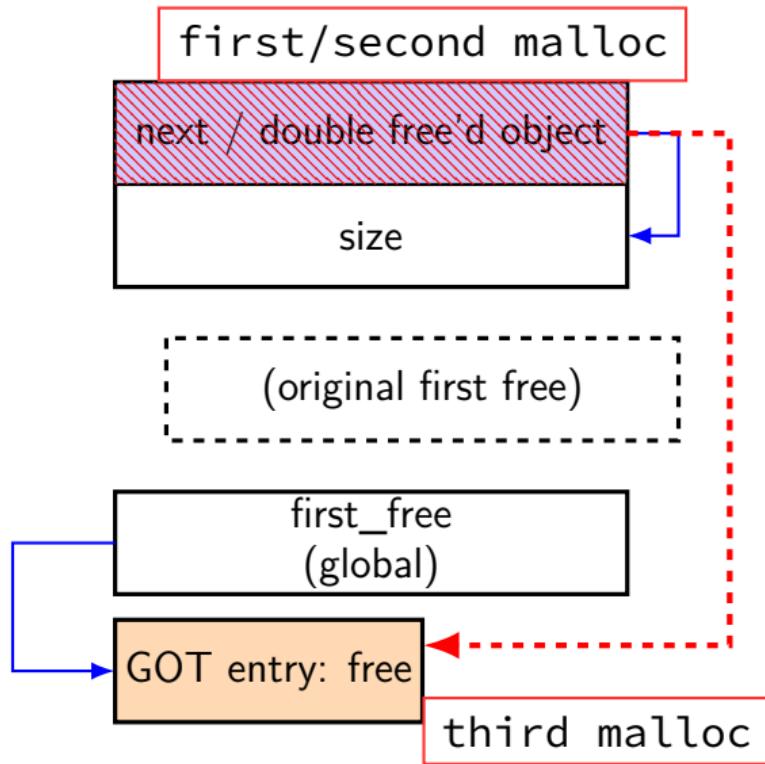
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```



double-free expansion

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first_free = chunk  
// malloc/new 1:  
result1 = first_free;  
first_free = first_free->next;  
// + overwrite:  
strcpy(result1, ...);  
// malloc/new 2:  
first_free = first_free->next;  
// malloc/new 3:  
result3 = first_free;  
strcpy(result3, ...);
```



double-free notes

this attack has apparently not been possible for a while

most malloc/new's **check for double-frees** explicitly
(e.g., look for a bit in size data)

prevents this issue — also catches programmer errors

pretty cheap

double-free exercise

```
free(...) {
    freed->next = first_free
    first_free = freed;
}
malloc(...) {
    if (can use first free) {
        void *to_return = first_free;
        first_free = first_free->next;
        return to_return;
    }
}
vulnerable() {
    char *p = malloc(100);
    free(p);
    free(p);
    char *q = malloc(100);
    char *r = malloc(100);
    strlcpy(q, attacker_input1, 100);
    char *s = malloc(100);
    strlcpy(r, attacker_input2, 100);
    strlcpy(s, attacker_input3, 100);
}
```

Exercise: which input should contain address to overwrite?
which input should contain value to put in that address?

logistical aside

originally was planning to cover more techniques like stack canaries

instead: we'll do later

will start with more principled approaches first

so far

many vulnerabilities we looked at due to poor bounds checking
one exception: use-after-free and related

can we just fix this?

adding bounds checking

```
char buffer[42];  
memcpy(buffer, attacker_controlled, len);
```

couldn't compiler add check for len

modern Linux: it does

added bounds checking

```
char buffer[42];
memcpy(buffer, attacker_controlled, len);
```

```
subq    $72, %rsp
leaq    4(%rsp), %rdi
movslq  len, %rdx
movq    attacker_controlled, %rsi
movl    $42, %ecx
call    __memcpy_chk
```

length 42 passed to __memcpy_chk

_FORTIFY_SOURCE

Linux C standard library + GCC features

adds automatic checking to a bunch of string/array functions

also printf (disable %n unless format string is a constant)

often enabled by default

GCC options:

- D_FORTIFY_SOURCE=1 — enable (backwards-compatible only)
- D_FORTIFY_SOURCE=2 — enable (full)
- U_FORTIFY_SOURCE — disable

bounds checking will happen...

will add checks (gcc 9.3 -O2)

```
void example1() {
    char dest1[1024]; memcpy(dest1, ...); ...
}
char dest2[1024];
void example2() {
    memcpy(dest2, ...); ...
}
void example3() {
    char *p = &dest2[4]; memcpy(p, ...); ...
}
```

bounds checking won't happen...

will not add check (gcc 9.3 -O2)

```
char dest2[1024];
void example4() {
    char *p = &dest2[mystery()]; memcpy(p, ...); ...
}
```

adds check for size 1024 (max possible size):

```
char dest2[1024];
void example5() {
    char dest3[128];
    char *p = dest2;
    if (mystery()) p = dest3;
    memcpy(p, ...); ...
}
```

non-checking library functions

some C library functions make bounds checking hard:

```
strcpy(dest, source);  
strcat(dest, source);  
sprintf(dest, format, ...);
```

bounds-checking versions (**added to library later**):

```
/* might not add \0 (!) */  
strncpy(dest, source, size);  
strncat(dest, source, size);  
snprintf(dest, size, format, ...);
```

poor bounds-checking APIs

```
char dest[100];
/* THIS CODE IS BROKEN */
strncpy(dest, source1, sizeof dest);
strncat(dest, source2, sizeof dest);
printf("result was %s\n", dest)
```

the above can access memory of out of bounds

...in a bunch of ways

Linux's strncpy manual

```
strncpy(dest, source1, sizeof dest);
```

“Warning: If there is no null byte among the first n bytes of src, the string placed in dest will not be null-terminated.”

exercise: what should the call have been?

Linux's strncat manual

```
strncat(dest, source2, sizeof dest);
```

“If src contains n or more bytes, strncat() writes n+1 bytes to dest (n from src plus the terminating null byte). Therefore, the size of dest must be at least strlen(dest)+n+1.”

exercise: what should the call have been?

better versions?

FreeBSD (and Linux via libbsd): `strlcpy`, `strlcat`

“Unlike [strncat and strncpy], `strlcpy()` and `strlcat()` take the full size of the buffer and guarantee to NUL-terminate the result...”

```
strlcpy(dest, source1, sizeof dest);  
strlcat(dest, source2, sizeof dest);
```

Windows: `strcpy_s`, `strcat_s` (same idea, different name)

C++ bounds checking

```
#include <vector>
...
std::vector<int> data;
data.resize(50);
// undefined behavior:
data[60] = 0;
// throws std::out_of_range exception
data.at(60) = 0;
```