

**CS216: Program and Data Representation**  
 University of Virginia Computer Science  
**Spring 2006** **David Evans**

## Lecture 18: Code Safety and Virtual Machines



(Duke suicide picture by Gary McGraw)

<http://www.cs.virginia.edu/cs216>

## JVM Instruction Set

|   |    |  |   |
|---|----|--|---|
| pushing constants   | 20 | getstatic, putstatic                                 | 2 |
| loads, stores (0-3 for each iload,<br>lload, fload, dload, aload) | 66 | newarray, anewarray,<br>multianewarray, arraylength  | 4 |
| pop, dup, swap, etc.  | 9  | invoke methods, throw                                | 5 |
| arithmetic  | 37 | new  | 1 |
| conversion (e.g., i2l)  | 15 | getfield, putfield                                   | 2 |
| comparisons (lcmp)  | 5  | checkcast  | 1 |
| goto, jsr, goto_w, jsr_w, ret                                     | 5  | instanceof   | 1 |
| tableswitch, lookupswitch   | 2  | monitorenter, monitorexit                            | 2 |
| returns (e.g., ireturn)   | 6  | wide   | 1 |
| conditional jumps (ifeq, ifnull,<br>ifnonnull)                    | 16 | nop, breakpoint, unused, implementation<br>dependent | 5 |

(205 out of 256 possible opcodes used)

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## How to get more than 256 local variables!

wide <opcode> <byte1> <byte2>

- Opcode is one of *iload*, *fload*, *aload*, *lload*, *dload*, *istore*, *fstore*, *astore*, *istore*, *dstore*, or *ret*
- Modifies instruction to take 2 byte operand ( $\text{byte1} \ll 8 | \text{byte2}$ )

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## Method Calls

- **invokevirtual <method>**
  - Invokes the method <method> on the parameters and object on the top of the stack.
  - Finds the appropriate method at run-time based on the actual type of the this object.

```
invokevirtual <Method void println(java.lang.String)>
```

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## Method Calls

- **invokestatic <method>**
  - Invokes a static (class) method <method> on the parameters on the top of the stack.
  - Finds the appropriate method at run-time based on the actual type of the this object.

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## Example

```
public class Sample1 {
    static public void main (String args[]) {
        System.out.println ("Hello!");
        System.exit (1);
    }
}
```

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```

> javap -c Sample1
Compiled from Sample1.java
public class Sample1 extends java.lang.Object {
    public Sample1();
    public static void main(java.lang.String[]);
}

Method Sample1()
0 aload_0
1 invokespecial #1 <Method java.lang.Object()>
4 return
Method void main(java.lang.String[])
0 getstatic #2 <Field java.io.PrintStream err>
3 ldc #3 <String "Hello!">
5 invokevirtual #4 <Method void println(java.lang.String)>
8 iconst_1
9 invokestatic #5 <Method void exit(int)>
12 return

```

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## Cast Instruction

```

public class Cast {
    static public void main (String args[]) {
        Object x;
        x = (Object) args[0];
        System.out.println ("result: " + (String) x);
    }
}

```

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```

Method void main(java.lang.String[])
0 aload_0
1 iconst_0
2 aaload
3 astore_1
4 getstatic #2 <Field java.lang.String valueOf>
7 new #3 <Class java.lang.String>
10 dup
11 invokespecial #4 <Method String valueOf(int)>
14 ldc #5 <String "result: ">
16 invokevirtual #6 <Method java.lang.StringBuffer append(java.lang.String)>
19 aload_1
20 checkcast #7 <Class java.lang.String>
23 invokevirtual #6 <Method java.lang.StringBuffer append(java.lang.String)>
26 invokevirtual #8 <Method java.lang.String toString()>
29 invokevirtual #9 <Method void println(java.lang.String)>
32 return

```

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## JVML Instruction Set

|  |    |   |   |
|--|----|---|---|
| pushing constants  | 20 | getstatic, putstatic                              | 2 |
| loads, stores (0-3 for each iload, iload, fload, dload, aload) | 66 | newarray, anewarray, multianewarray, arraylength  | 4 |
| pop, dup, swap, etc.   | 9  | invoke methods, throw                             | 5 |
| arithmetic   | 37 | new   | 1 |
| conversion (e.g., i2l)   | 15 | getfield, putfield                                | 2 |
| comparisons (lcmp)   | 5  | checkcast   | 1 |
| goto, jsr, goto_w, jsr_w, ret                                  | 5  | instanceof  | 1 |
| tableswitch, lookupswitch                                      | 2  | monitorenter, monitorexit                         | 2 |
| returns (e.g., ireturn)  | 6  | wide  | 1 |
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(205 out of 256 possible opcodes used)

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**The Worst Instruction**  
<http://java.sun.com/docs/books/vmspec/2nd-edition/html/Instructions2.doc7.html>

**jsr [branchbyte1] [branchbyte2]**

**Forms**  
 $jsr = 168 (0xa8)$

**Operand Stack**  
 $\dots \Rightarrow \dots, address$

**Description**  
The address of the opcode of the instruction immediately following this *jsr* instruction is pushed onto the operand stack as a value of type *returnAddress*. The unsigned *branchbyte1* and *branchbyte2* are used to construct a signed 16-bit offset, where the offset is  $(branchbyte1 \ll 8) | branchbyte2$ . Execution proceeds at that offset from the address of this *jsr* instruction. The target address must be that of an opcode of an instruction within the method that contains this *jsr* instruction.

**Notes**  
The *jsr* instruction is used with the *ret* instruction in the implementation of the finally clauses of the Java programming language. Note that *jsr* pushes the address onto the operand stack and *ret* gets it out of a local variable. This asymmetry is intentional.

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## Try-Catch-Finally

```

public class JSR {
    static public void main (String args[]) {
        try {
            System.out.println("hello");
        } catch (Exception e) {
            System.out.println ("There was an exception!");
        } finally {
            System.out.println ("I am finally here!");
        }
    }
}

```

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```

Method void main(java.lang.String[])
  0 getstatic #2 <Field java.io.PrintStream
  3 ldc #3 <String "hello">
  5 invokevirtual #4 <Method void println
  8 jsr 35
 11 goto 46
14 astore_1
15 getstatic #2 <Field java.io.PrintStream
18 ldc #6 <String "There was an exception!">
20 invokevirtual #4 <Method void println>
23 jsr 35
26 goto 46
29 astore_2
30 jsr 35
33 aload_2
34 athrow
35 astore_3
36 getstatic #2 <Field java.io.PrintStream
39 ldc #7 <String "I am finally!">
41 invokevirtual #4 <Method void println>
44 ret 3
46 return

```

Exception table:

| from | to | target type                 |
|------|----|-----------------------------|
| 0    | 8  | <Class java.lang.Exception> |
| 0    | 11 | any                         |
| 14   | 26 | any                         |
| 29   | 33 | any                         |

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**Java™: Programming Language**  
 compared to C++, not to C      sort of  
**"A simple, object-oriented,  
 distributed, interpreted, robust,  
 secure, architecture neutral,  
 portable, high-performance,  
 multithreaded, and dynamic  
 language."** [Sun95]

Java: int is 32 bits  
 C: int is >= 16 bits

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## What is a secure programming language?



- Language is designed so it cannot express certain computations considered insecure.  
 A few attempt to do this: PLAN, packet filters
- Language is designed so that (accidental) program bugs are likely to be caught by the compiler or run-time environment instead of leading to security vulnerabilities.

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## Safe Programming Languages

- Type Safety
  - Compiler and run-time environment ensure that bits are treated as the type they represent
- Memory Safety
  - Compiler and run-time environment ensure that program cannot access memory outside defined storage
- Control Flow Safety
  - Can't jump to arbitrary addresses
  - Which of these does C/C++ have?
  - Is Java the first language to have them?
  - No way! LISP had them all in 1960.

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## Java™ Safety

- Type Safety
  - Most types checked statically
  - Coercions, array assignments type checked at run time
- Memory Safety
  - No direct memory access (e.g., pointers)
  - Primitive array type with mandatory run-time bounds checking
- Control Flow Safety
  - Structured control flow, no arbitrary jumps

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## Malicious Code

Can a safe programming language protect you from malcode?



- Code your servers in it to protect from buffer overflow bugs
- Only allow programs from untrustworthy origins to run if they are programmed in the safe language

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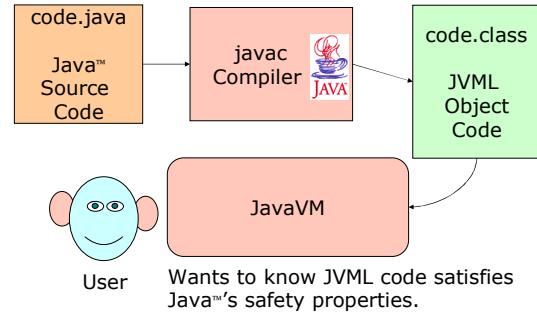
## Safe Languages?

- But how can you tell program was written in the safe language?
  - Get the source code and compile it (most vendors, and all malicious attackers refuse to provide source code)
  - Special compilation service cryptographically signs object files generated from the safe language (SPIN, [Bershad96])
  - Verify object files preserve safety properties of source language (Java) 

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## JVML



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## Does JVML satisfy Java™'s safety properties?

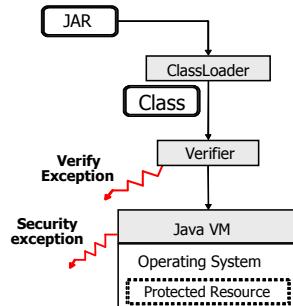
**iconst\_2**      push integer constant 2 on stack  
**istore\_0**      store top of stack in variable 0 as int  
**aload\_0**      load object reference from variable 0

No! This code violates Java™'s type rules. 

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## Java Security Architecture

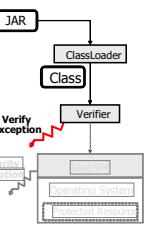


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## Mistyped Code

```
.method public static main([Ljava/lang/String;)V
...
iconst_2
istore_0
aload_0
iconst_2
iconst_3
iadd
...
.end method
```



> java Simple  
Exception in thread  
"main"  
java.lang.VerifyError:  
(class: Simple, method:  
main signature:  
([Ljava/lang/String;)V)  
**Register 0 contains  
wrong type**

Verifier error before any code runs

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## Runtime Error

```
public class Cast {
    static public void main (String args[]) {
        Object o = new Object ();
        Method void main(java.lang.String[])
        0 new #2 <Class java.lang.Object>
        3 dup
        4 invokespecial #1 <Method
        java.lang.Object()>
        7 astore_1
        8 aload_1
        9 checkcast #3 <Class java.lang.String>
        12 astore_2
        13 getstatic #4 <Field java.io.PrintStream
        out>
        16 aload_2
        17 invokevirtual #5 <Method void
        println(java.lang.String)>
        20 return
    }
}
```

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## Bytecode Verifier

- Checks class file is formatted correctly
  - Magic number: class file starts with 0xCAFEBAE
  - String table, code, methods, etc.
- Checks JVM code satisfies safety properties
  - Simulates program execution to know types are correct, but doesn't need to examine any instruction more than once

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## Verifying Safety Properties

- Type safe
  - Stack and variable slots must store and load as same type
  - Only use operations valid for the data type
- Memory safe
  - Must not attempt to pop more values from stack than are on it
  - Doesn't access private fields and methods outside class implementation
- Control flow safe
  - Jumps must be to valid addresses within function, or call/return

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## Charge

- PS6 will be out (electronically) on Friday
- If you would like to be assigned a partner for PS6, send me email as soon as possible

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