

cs2220: Engineering Software

Class 8: Implementing Data Abstractions

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Menu

Implementing Data Abstractions

Abstraction Function

Representation Invariant

Recap: Abstract Data Types

Separate *what* you can do with data from *how* it is represented

Client interacts with data through provided operations according to their specifications

Implementation chooses how to represent data and implement its operations

Data Abstraction in Java

A **class** defines a new data type

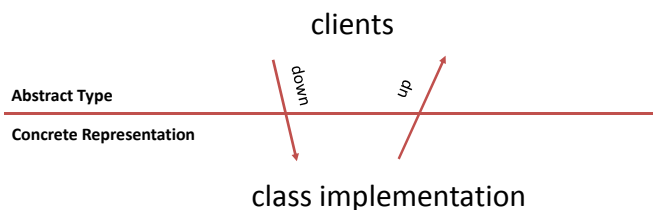
Use **private** instance variables to hide the choice of representation

private variables are only visible inside the class

```
public class StringStack {  
    // Rep:  
    private List<String> rep;  
}
```

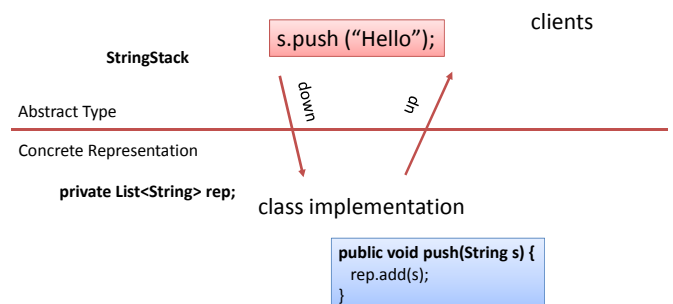
Up and Down

Clients manipulate an abstract data type by calling its operations (methods and constructors)



The representation of an abstract data type is visible only in the class implementation.

StringStack



Advantages/Disadvantages of Data Abstraction

- More code to write and maintain
- Run-time overhead (time to call method, lost opportunities because of abstraction)
- + Client doesn't need to know about representation
- + Can change rep without changing clients
- + Can reason about clients at abstract level

Choosing a Representation

Representation must **store the abstract state**

Think about how methods will be implemented

A good representation choice should:

Enable **easy implementations** of all methods

Allow **performance-critical methods** to be implemented efficiently

Use memory efficiently (if the data may be large)

Choosing the rep is the most important decision in implementing a data abstraction

StringStack Representation

Option 1: **private String [] rep;**

- Recall Java arrays are bounded
- Hard to implement **push**

Option 2: **private List<String> rep;**

- Easy to implement all methods
- Performance may be worse than for array

Implementing StringStack

```
public class StringStack {
    // Rep:
    private List<String> rep;

    /**
     * MODIFIES: this
     * EFFECTS: Pushes s on the top of this.
     * For example, if this_pre = [ e_n-1, e_n-2, ..., e_1, e_0 ],
     * this_post = [ s, e_n-1, e_n-2, ..., e_1, e_0 ]
     */
    public void push(String s) {
        rep.add(s);
    }
}
```

Is this implementation of **push** correct?

Could this implementation of **push** be correct?

It depends...

```
/**
 * MODIFIES: this
 * EFFECTS: If this is empty, throws EmptyStackException. Otherwise,
 * returns the element on top of this and removes that element from this.
 * For example, if this_pre = [ e_n-1, e_n-2, ..., e_1, e_0 ],
 * this_post = [ e_n-2, ..., e_1, e_0 ] and the result is e_n-1.
 */
public String pop() throws EmptyStackException {
    try {
        return rep.remove(0);
    } catch (IndexOutOfBoundsException e) {
        assert rep.size() == 0;
        throw new EmptyStackException();
    }
}
```

```
public String pop() throws EmptyStackException {
    try {
        return rep.remove(rep.size() - 1);
    } catch (IndexOutOfBoundsException e) {
        assert rep.size() == 0;
        throw new EmptyStackException();
    }
}
```

Is it correct?

- How can we possibly implement data abstractions correctly if correctness of one method depends on how other methods are implemented?
- How can we possibly test a data abstraction implementation if there are complex interdependencies between methods?

What must we know to know
if **pop** is correct?

```
/**
 * MODIFIES: this
 * EFFECTS: If this is empty, throws EmptyStackException. Otherwise,
 * returns the element on top of this and removes that element from this.
 * For example, if this_pre = [ e_n-1, e_n-2, ..., e_1, e_0 ],
 * this_post = [ e_n-2, ..., e_1, e_0 ] and the result is e_n-1.
 */
public String pop() throws EmptyStackException {
    try {
        return rep.remove(0);
    } catch (IndexOutOfBoundsException e) {
        assert rep.size() == 0;
        throw new EmptyStackException();
    }
}
```

Abstraction Function

The **Abstraction Function** maps a concrete state to an abstract state:

$$\mathcal{AF}: C \rightarrow \mathcal{A}$$

Function from concrete representation to the abstract notation introduced in Overview specification.

Abstraction Function for StringStack

```
/**
 * OVERVIEW: A StringStack represents a last-in-first-out stack where all
 * elements are Strings.
 * A typical stack is [ e_n-1, e_n-2, ..., e_1, e_0 ] where e_n-1 is the top
 * of the stack.
 */
public class StringStack {
    // Rep:
    private List<String> rep;
```

Handwritten notes:

- `void push(String s) { rep.add(s); }`
- `rep: [1, 2, 3]`
- `add(4) → [1, 2, 3, 4]`
- $AF(c)$
- $c.rep \rightsquigarrow [e_{n-1}, e_{n-2}, \dots, e_1, e_0]$
- $[e_0, e_1, \dots, e_{n-1}]_i = rep.get(i)$
- $(rep.size() - i - 1)$

Correctness of Push

```
/**
 * MODIFIES: this
 * EFFECTS: Pushes s on the top of this.
 * For example, if this_pre = [ e_n-1, e_n-2, ..., e_1, e_0 ],
 * this_post = [ s, e_n-1, e_n-2, ..., e_1, e_0 ]
 */
public void push(String s) {
    rep.add(s);
}
```

Use abstraction function to show push satisfies its specification:

$$\mathcal{AF}(rep_post) = [\mathcal{AF}_{String}(s)] + \mathcal{AF}(rep_pre)$$

```
/**
 * MODIFIES: this
 * EFFECTS: Pushes s on the top of this.
 * For example, if this_pre = [ e_n-1, e_n-2, ..., e_1, e_0 ],
 * this_post = [ s, e_n-1, e_n-2, ..., e_1, e_0 ]
 */
public void push(String s) {
    rep.add(s);
}
```

How do we know **rep** is non-null?

Rep Invariant

Representation Invariant: properties all legitimate objects of the ADT must satisfy

$$I: C \rightarrow \text{boolean}$$

Function from **concrete representation** to **boolean**.

Helps us reason about correctness of methods independently

Limits the range of inputs for the Abstraction Function

Reasoning with Rep Invariants

Prove all objects satisfy the invariant **before** leaving the implementation code

Assume all objects **passed in** satisfy the invariant

REQUIRES: Rep Invariant is true for this (and any other reachable ADT objects)
EFFECTS: Rep Invariant is true for all new and any modified ADT object on exit.

All non-private datatype operations have these specification implicitly!

Rep Invariant for StringStack

```
/**
 * OVERVIEW: A StringStack represents a last-in-first-out stack where all
 * elements are Strings.
 * A typical stack is [ e_n-1, e_n-2, ..., e_1, e_0 ] where e_n-1 is the top
 * of the stack.
 */
```

```
public class StringStack {
    // Rep:
    private List<String> rep;
```

1) Rep Invariant: $I(c) = \text{rep} \neq \text{null}$

Graph

```
public class Graph
// OVERVIEW: A Graph is a mutable type that represents an undirected graph. It consists of
// nodes that are named by Strings, and edges that connect a pair of nodes.
// A typical Graph is: < Nodes, Edges > where
// Nodes = { n_1, n_2, ..., n_m }
// Edges = { {a_1, b_1}, ..., {a_n, b_n} } (the elements of Edges are unordered sets).
```

```
public Graph ()
// EFFECTS: Initializes this to a graph with no nodes or edges: < {}, {} >.
```

```
public void addNode (String name) throws DuplicateException
// MODIFIES: this
// EFFECTS: If name is in Nodes, throws DuplicateException.
// Otherwise, adds a node named name to this:
// this_post = < Nodes_pre U { name }, Edges_pre >
```

```
public void addEdge (String s, String t) throws NoNodeException, DuplicateException
// MODIFIES: this
// EFFECTS: If s and t are not names of nodes in this, throws NoNodeException. If there is
// already an edge between s and t, throws DuplicateEdgeException. Otherwise, adds an
// edge between s and t to this:
// this_post = < Nodes_pre, Edges_pre U { {s, t} } >
```

Charge

Problem Set 3: Designing and Implementing
Data Abstractions: Due Tuesday