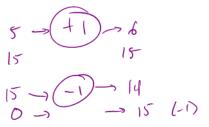
Fetch, Decode, Execute Instruction Set Architecture

CS 2130: Computer Systems and Organization 1

February 10, 2023

Announcements

- · Quiz 3 available today, due Sunday by 11:59pm
- Homework 2 due Monday



Our story so far

- Information modeled by voltage through wires (1 vs 0)
- Transistors
- Gates: **&** ~
- Multi-bit values: representing integers
- Floating point
- Multi-bit operations using circuits
- Storing results using registers
- Memory

Code

How do we run code? What do we need?

Example Code

```
•••
```

```
8: x = 16
```

9:
$$y = x$$

10:
$$x += y$$

•••

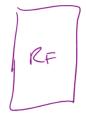
What is the value of x after line 10?

Bookkeeping

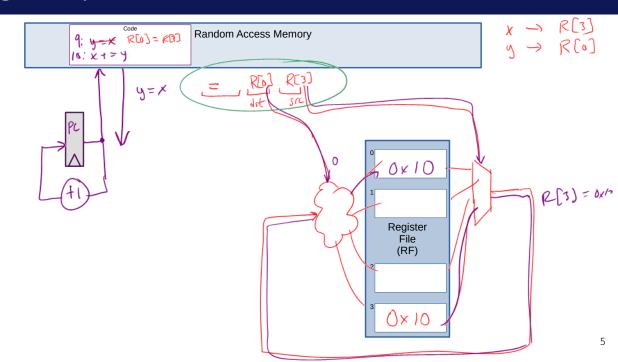


What do we need to keep track of?

- Code the program we are running
 - RAM (Random Access Memory)
- State things that may change value (i.e., variables)
 - · Register file can read and write values each cycle
- Program Counter (PC) where we are in our code
 - · Single register byte number in memory for next instruction



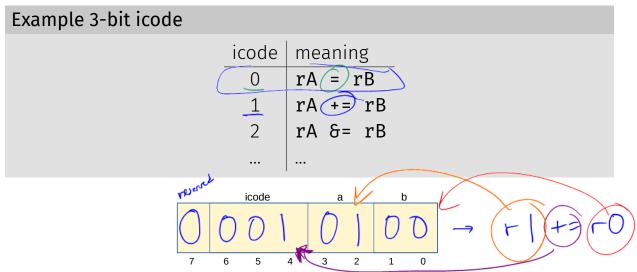
Building a Computer



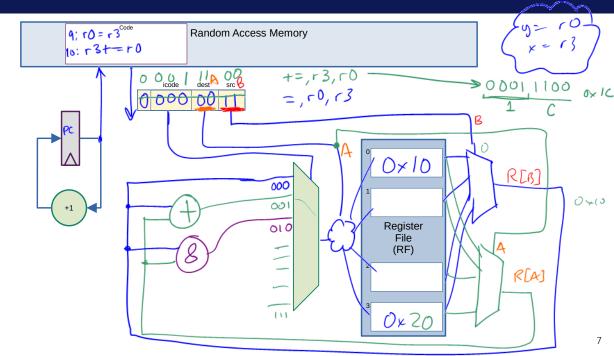
Encoding Instructions

Encoding of Instructions (icode or opcode)

Numeric mapping from icode to operation

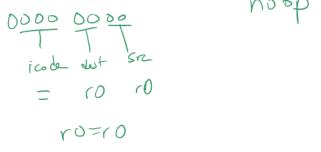


Building a Computer



Question

What happens if we get the 0-byte instruction? 00



Our Computer's Instructions

Toy ISA 3-bit icode

icode	meaning
0	rA = rB
1	rA += rB
2	rA &= rB
3	$\mathbf{r}\mathbf{A}$ = read from memory at address $\mathbf{r}\mathbf{B}$
4	write ${f r}{f A}$ to memory at address ${f r}{f B}$
•••	
7	Compare rA as 8-bit 2's-complement to 0
	if $rA \le 0$ set $pc = rB$
	else increment pc as normal

Our Computer's Instructions

Toy ISA 3-bit icode		
icode	b	action
5	0	rA = ~rA
	1	rA = -rA
	2	rA = !rA
	3	rA = pc
6	0	rA = read from memory at pc + 1
	1	rA += read from memory at pc + 1
	2	rA &= read from memory at pc + 1
	3	rA = read from memory at the address stored at $pc + 1$
		For icode 6, increase pc by 2 at end of instruction

High-level Instructions

In general, 3 kinds of instructions

- moves move values around without doing "work"
- math broadly doing "work"
- jumps jump to a new place in the code

Moves

Few forms

- Register to register (icode 0), x = y
- Register to/from memory (icodes 3-4), x = M[b], M[b] = x

Memory

- · Address: an index into memory.
 - · Addresses are just (large) numbers
 - Usually we will not look at the number and trust it exists and is stored in a register

Moves

Toy ISA 3-bit icode			
icode	b	action	
0		rA = rB	
3		$\mathbf{r}\mathbf{A}$ = read from memory at address $\mathbf{r}\mathbf{B}$	
4		write ${f r}{f A}$ to memory at address ${f r}{f B}$	
5	3	rA = pc	
6	0	rA = read from memory at pc + 1	
	3	rA = read from memory at the address stored at $pc + 1$	

Math

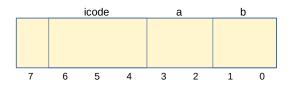
Broadly doing work

Toy ISA 3-bit icode			
	icode	b	meaning
	1		rA += rB
	2		rA &= rB
	5	0	rA = ~rA
		1	rA = -rA
		2	rA = !rA
	6	1	rA += read from memory at pc + 1
		2	rA &= read from memory at pc + 1

Note: We can implement other operations using these things!

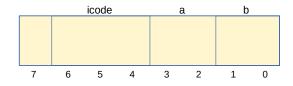
icodes 5 and 6

Special property of icodes 5-6: only one register used



icodes 5 and 6

Special property of 5-6: only one register used

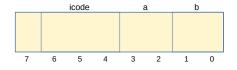


- · Side effect: all bytes between 0 and 127 are valid instructions!
- As long as high-order bit is 0
- · No syntax errors, any instruction given is valid

Immediate values

icode 6 provides literals, **immediate** values

Example 3-bit icode		
icode	b	action
6		rA = read from memory at pc + 1
	1	rA += read from memory at pc + 1
	2	${ m rA}$ &= read from memory at ${ m pc}$ + 1
	3	rA = read from memory at the address stored at $pc + 1$
		For icode 6, increase pc by 2 at end of instruction





Encoding Instructions

Example 1: r1 += 19

Instructions

icode	b	meaning
0		rA = rB
1		rA += rB
2		rA &= rB
3		${f r}{f A}$ = read from memory at address ${f r}{f B}$
4		write ${f r}{f A}$ to memory at address ${f r}{f B}$
5	0	rA = ~rA
	1	rA = -rA
	2	rA = !rA
	3	rA = pc
6	0	rA = read from memory at pc + 1
	1	rA += read from memory at pc + 1
	2	rA &= read from memory at $pc + 1$
	3	rA = read from memory at the address stored at $pc + 1$
		For icode 6, increase pc by 2 at end of instruction
7		Compare rA as 8-bit 2's-complement to 0
		if $rA \le 0$ set $pc = rB$
		else increment pc as normal

Encoding Instructions

Example 2: M[0x82] += r3

Read memory at address 0x82, add r3, write back to memory at same address

Instructions

icode	b	meaning
0		rA = rB
1		rA += rB
2		rA &= rB
3		${f r}{f A}$ = read from memory at address ${f r}{f B}$
4		write ${f r}{f A}$ to memory at address ${f r}{f B}$
5	0	rA = ~rA
	1	rA = -rA
	2	rA = !rA
	3	rA = pc
6	0	rA = read from memory at $pc + 1$
	1	rA += read from memory at pc + 1
	2	rA &= read from memory at $pc + 1$
	3	rA = read from memory at the address stored at $pc + 1$
		For icode 6, increase pc by 2 at end of instruction
7		Compare rA as 8-bit 2's-complement to 0
		if $rA \le 0$ set $pc = rB$
		else increment pc as normal

Jumps

- Moves and math are large portion of our code
- · We also need **control constructs**
 - · Change what we are going to do next
 - if, while, for, functions, ...
- Jumps provide mechanism to perform these control constructs
- We jump by assigning a new value to the program counter PC

Jumps

For example, consider an **if**

Jumps

Example 3-bit icode		
	icode	meaning
	7	Compare rA as 8-bit 2's-complement to 0
		if rA <= 0 set pc = rB
		else increment pc as normal

Instruction icode 7 provides a conditional jump

 Real code will also provide an unconditional jump, but a conditional jump is sufficient

Writing Code

We can now write any* program!

- · When you run code, it is being turned into instructions like ours
- Modern computers use a larger pool of instructions than we have (we will get there)

*we do have some limitations, since we can only represent 8-bit values and some operations may be tedious.

Our code to this machine code

How do we turn our control constructs into jump statements?

if/else to jump

while to jump

Function Calls

Encoding Instructions

Example 3: if r0 < 9 jump to 0x42

Instructions

icode	b	meaning
0		rA = rB
1		rA += rB
2		rA &= rB
3		${f r}{f A}$ = read from memory at address ${f r}{f B}$
4		write ${f r}{f A}$ to memory at address ${f r}{f B}$
5	0	rA = ~rA
	1	rA = -rA
	2	rA = !rA
	3	rA = pc
6	0	rA = read from memory at $pc + 1$
	1	rA += read from memory at pc + 1
	2	rA &= read from memory at pc + 1
	3	rA = read from memory at the address stored at $pc + 1$
		For icode 6, increase pc by 2 at end of instruction
7		Compare rA as 8-bit 2's-complement to 0
		if rA <= 0 set pc = rB
		else increment pc as normal