Fetch, Decode, Execute Instruction Set Architecture

CS 2130: Computer Systems and Organization 1 February 10, 2023

- Quiz 3 available today, due Sunday by 11:59pm
- Homework 2 due Monday

$$\begin{array}{c} \cancel{5} \\ \cancel{$$

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Our story so far

- Information modeled by voltage through wires (1 vs 0)
- Transistors
- Gates: & | ~ ^
- Multi-bit values: representing integers
- Floating point
- Multi-bit operations using circuits
- Storing results using registers
- Memory

How do we run code? What do we need?

Example Code ... 8: x = 16 9: y = x 10: x += y ...

What is the value of *x* after line 10?

What do we need to keep track of?

- Code the program we are running
 - RAM (Random Access Memory)
- State things that may change value (i.e., variables)
 - Register file can read and write values each cycle
- Program Counter (PC) where we are in our code
 - Single register byte number in memory for next instruction



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Building a Computer



Encoding Instructions

Encoding of Instructions (icode or opcode)

 \cdot Numeric mapping from icode to operation



Building a Computer



Question



Our Computer's Instructions

| Toy ISA 3-bit icode | | | | |
|---------------------|-------|---|--|--|
| | icode | meaning | | |
| | 0 | rA = rB | | |
| | 1 | rA += rB | | |
| | 2 | $rA \delta = rB$ | | |
| | 3 | rA = read from memory at address rB | | |
| | 4 | write \mathbf{rA} to memory at address $\mathbf{rB} - \mathcal{M}[rB] = rA$ | | |
| | | | | |
| | 7 | Compare rA as 8-bit 2's-complement to 0 | | |
| | | if $rA \leq = 0$ set $pc = rB$ | | |
| | | else increment pc as normal | | |

Our Computer's Instructions

| Toy ISA 3-bit icode | | | | |
|---------------------|----------------|--|--|--|
| icode | icode b action | | | |
| 5 | 0 | $rA = \sim rA$ | | |
| | 1 | rA = -rA | | |
| | 2 | rA = !rA | | |
| | 3 | rA = pc | | |
| 6 | 0 | rA = read from memory at pc + 1 | | |
| | 1 | rA += read from memory at pc + 1 | | |
| | 2 | rA &= read from memory at pc + 1 | | |
| | 3 | rA = read from memory at the address stored at pc + 1 | | |
| | | For icode 6, increase pc by 2 at end of instruction | | |

In general, 3 kinds of instructions

- moves move values around without doing "work"
- math broadly doing "work"
- jumps jump to a new place in the code

Few forms

- Register to register (icode 0), x = y
- Register to/from memory (icodes 3-4), x = M[b], M[b] = x

Memory

- Address: an index into memory.
 - Addresses are just (large) numbers
 - Usually we will not look at the number and trust it exists and is stored in a register

| Toy ISA 3-bit icode | | | |
|---------------------|---------------------------------------|---|--|
| icode | icode b action | | |
| 0 | | rA = rB | |
| 3 | 3 rA = read from memory at address rB | | |
| 4 | | write ${f r}{f A}$ to memory at address ${f r}{f B}$ | |
| 5 | 3 | rA = pc | |
| 6 | 0 | rA = read from memory at pc + 1 | |
| | 3 | rA = read from memory at the address stored at pc + 1 | |

Math

Broadly doing work

Toy ISA 3-bit icode

| icode | b | meaning |
|-------|---|------------------------------------|
| 1 | | rA += rB |
| 2 | | rA δ= rB |
| 5 | 0 | $rA = \sim rA$ |
| | 1 | rA = -rA |
| | 2 | rA = !rA |
| 6 | 1 | rA += read from memory at pc + 1 |
| | 2 | rA &= read from memory at pc + 1 |

Note: We can implement other operations using these things!

icodes 5 and 6

Special property of icodes 5-6: only one register used



Toy ISA 3-bit icode

| icode | b | action |
|-------|---|----------------|
| 5 | 0 | $rA = \sim rA$ |
| | 1 | rA = -rA |
| | 2 | rA = !rA |
| | 3 | rA = pc |

Special property of 5-6: only one register used



- Side effect: all bytes between 0 and 127 are valid instructions!
- As long as high-order bit is 0
- \cdot No syntax errors, any instruction given is valid

icode 6 provides literals, **immediate** values

Example 3-bit icode

| icode | b | action |
|-------|---|--|
| 6 | 0 | rA = read from memory at pc + 1 |
| | 1 | rA += read from memory at pc + 1 |
| | 2 | rA &= read from memory at pc + 1 |
| | 3 | rA = read from memory at the address stored at pc + 1 |
| | | For icode 6, increase pc by 2 at end of instruction |





Encoding Instructions

Example 1: r1 += 19

Instructions

| icode | b | meaning |
|-------|---|--|
| 0 | | rA = rB |
| 1 | | rA += rB |
| 2 | | rA &= rB |
| 3 | | ${f r}{f A}$ = read from memory at address ${f r}{f B}$ |
| 4 | | write ${f r}{f A}$ to memory at address ${f r}{f B}$ |
| 5 | 0 | $rA = \sim rA$ |
| | 1 | rA = -rA |
| | 2 | rA = !rA |
| | 3 | rA = pc |
| 6 | 0 | rA = read from memory at pc + 1 |
| | 1 | rA += read from memory at pc + 1 |
| | 2 | rA &= read from memory at $pc + 1$ |
| | 3 | rA = read from memory at the address stored at pc + 1 |
| | | For icode 6, increase pc by 2 at end of instruction |
| 7 | | Compare rA as 8-bit 2's-complement to 0 |
| | | if rA <= 0 set pc = rB |
| | | else increment pc as normal |

Example 2: M[0x82] += r3

Read memory at address 0x82, add r3, write back to memory at same address

Instructions

| icode | b | meaning |
|-------|---|--|
| 0 | | rA = rB |
| 1 | | rA += rB |
| 2 | | rA &= rB |
| 3 | | ${f r}{f A}$ = read from memory at address ${f r}{f B}$ |
| 4 | | write ${f r}{f A}$ to memory at address ${f r}{f B}$ |
| 5 | 0 | $rA = \sim rA$ |
| | 1 | rA = -rA |
| | 2 | rA = !rA |
| | 3 | rA = pc |
| 6 | 0 | rA = read from memory at pc + 1 |
| | 1 | rA += read from memory at pc + 1 |
| | 2 | rA &= read from memory at $pc + 1$ |
| | 3 | rA = read from memory at the address stored at pc + 1 |
| | | For icode 6, increase pc by 2 at end of instruction |
| 7 | | Compare rA as 8-bit 2's-complement to 0 |
| | | if rA <= 0 set pc = rB |
| | | else increment pc as normal |

- \cdot Moves and math are large portion of our code
- $\cdot\,$ We also need control constructs
 - \cdot Change what we are going to do next
 - if, while, for, functions, ...
- Jumps provide mechanism to perform these control constructs
- We jump by assigning a new value to the program counter **PC**



For example, consider an **if**



Example 3-bit icode

| icode | meaning |
|-------|--|
| 7 | Compare \mathbf{rA} as 8-bit 2's-complement to 0 |
| | if $rA <= 0$ set $pc = rB$ |
| | else increment pc as normal |

Instruction icode 7 provides a **conditional** jump

• Real code will also provide an **unconditional** jump, but a conditional jump is sufficient

We can now write any* program!

- When you run code, it is being turned into instructions like ours
- Modern computers use a larger pool of instructions than we have (we will get there)

*we do have some limitations, since we can only represent 8-bit values and some operations may be tedious.

How do we turn our control constructs into jump statements?

if/else to jump

while to jump

Function Calls

Example 3: if r0 < 9 jump to 0x42

Instructions

| icode | b | meaning |
|-------|---|---|
| 0 | | rA = rB |
| 1 | | rA += rB |
| 2 | | rA &= rB |
| 3 | | $\mathbf{r}\mathbf{A}$ = read from memory at address $\mathbf{r}\mathbf{B}$ |
| 4 | | write ${f r}{f A}$ to memory at address ${f r}{f B}$ |
| 5 | 0 | $rA = \sim rA$ |
| | 1 | rA = -rA |
| | 2 | rA = !rA |
| | 3 | rA = pc |
| 6 | 0 | rA = read from memory at pc + 1 |
| | 1 | rA += read from memory at $pc + 1$ |
| | 2 | rA &= read from memory at pc + 1 |
| | 3 | rA = read from memory at the address stored at pc + 1 |
| | | For icode 6, increase pc by 2 at end of instruction |
| 7 | | Compare rA as 8-bit 2's-complement to 0 |
| | | if rA <= 0 set pc = rB |
| | | else increment pc as normal |