Toy Instruction Set Architecture

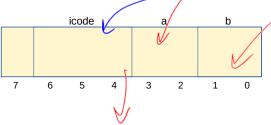
CS 2130: Computer Systems and Organization 1 February 17, 2023

- Homework 3 due **Wednesday** at 11pm on Gradescope
- Quiz 4 available today, due Sunday at 11:59pm (submit early)
- Exam 1 next Friday in class, Review on Wednesday

Encoding Instructions

Encoding of Instructions

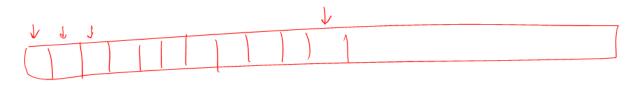
- 3-bit icode (which operation to perform)
 - \cdot Numeric mapping from icode to operation
- Which registers to use (2 bits each)
- Reserved bit for future expansion



In general, 3 kinds of instructions

- moves move values around without doing "work"
- math broadly doing "work"
- jumps jump to a new place in the code

- \cdot Moves and math are large portion of our code
- $\cdot\,$ We also need control constructs
 - \cdot Change what we are going to do next
 - if, while, for, functions, ...
- Jumps provide mechanism to perform these control constructs
- $\cdot\,$ We jump by assigning a new value to the program counter PC





Toy ISA 3-bit icode

| icode | meaning |
|-------|---|
| 7 | Compare ${f r}{f A}$ as 8-bit 2's-complement to ${f 0}$ |
| | if $rA <= 0$ set $pc = rB$ |
| | else increment pc as normal |

Instruction icode 7 provides a **conditional** jump

• Real code will also provide an **unconditional** jump, but a conditional jump is sufficient

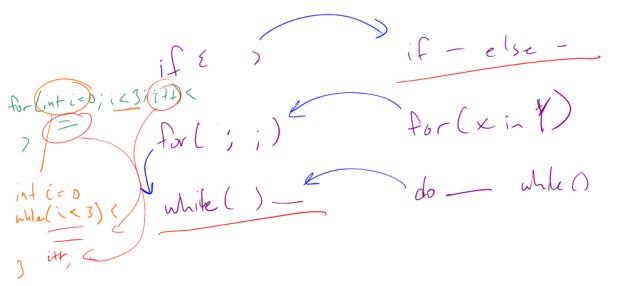
We can now write any* program!

- \cdot When you run code, it is being turned into instructions like ours
- Modern computers use a larger pool of instructions than we have (we will get there)

*we do have some limitations, since we can only represent 8-bit values and some operations may be tedious.

Our code to this machine code

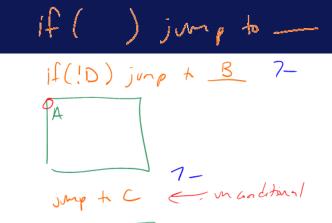
How do we turn our control constructs into jump statements?

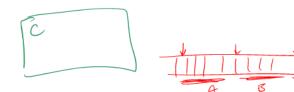


if/else to jump

John E

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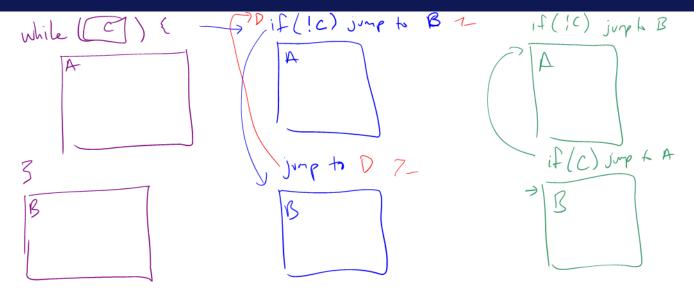


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while to jump

option 2

Option 2



Encoding Instructions

| | i | icode | b | meaning | | | | | |
|---------------------|----|-----------|---|--|-----------------|-----|-----|---|----|
| A C | (- | 0 | | rA = rB | Fx 3: if | r0 | < | 9 jump | to |
| | | 1 | | rA += rB | | - • | | Jamp | |
| | | 2 | | rA δ= rB | 0x42 | | | | |
| | 1 | 3 | | rA = read from memory at address rB | | | | | |
| | | 4 | | write ${f r}{f A}$ to memory at address ${f r}{f B}$ | | | | | |
| (| | 5 | 0 | $rA = \sim rA$ | | | | | |
| | | | 1 | rA = -rA | | | | | |
| | | | 2 | $r\Delta = !rA$ | | | | | |
| | | \langle | 3 | rA = pc | | | | | |
| | | 6 | Û | rA = read from memory at pc + 1 | | 1 | ٨ | 1. | |
| | | (| 1 | rA += read from memory at pc + 1 | Jbyte | + | 124 | The second se | |
| | | | 2 | rA &= read from memory at pc + 1 | \ | | Im | ente | |
| | | | 3 | rA = read from memory at the address stored at pc + 1 |) | | | | |
| | | | | For icode 6, increase pc by 2 at end of instruction | | | | | |
| | | 7 | | Compare rA as 8-bit 2's-complement to 0 | | | | | |
| Jose Association 20 | | | | if rA <= 0 set pc = rB | | | | | |
| 18 | | | | else increment pc as normal | | | | | |
| 6 | | | | | | | | | |

Example

+ Dx17 + 0x17 $x = 0x17 \times 3$ 00 D x = D -0#= O fr-(i=0; i<3; i++) ь r1\$ = × 4 6 _3 $\chi += 0 \times 17;$ -3=PC 00 01 Pox += 0×17 X=0 1 =0 d\$ +=. while (i < 3) { 1(r1< 3) -F3 jung to $X + = 0 \times 17;$ var reg i+=1; -1 <= 2 x r D 1 と=0 11

Example

| icode | b | meaning |
|-------|---|--|
| 0 | | rA = rB |
| 1 | | rA += rB |
| 2 | | rA &= rB |
| 3 | | ${f r}{f A}$ = read from memory at address ${f r}{f B}$ |
| 4 | | write ${f r}{f A}$ to memory at address ${f r}{f B}$ |
| 5 | 0 | $rA = \sim rA$ |
| | 1 | rA = -rA |
| | 2 | rA = !rA |
| | 3 | rA = pc |
| 6 | 0 | rA = read from memory at pc + 1 |
| | 1 | rA += read from memory at pc + 1 |
| | 2 | rA &= read from memory at pc + 1 |
| | 3 | rA = read from memory at the address stored at pc + 1 |
| | | For icode 6, increase pc by 2 at end of instruction |
| 7 | | Compare rA as 8-bit 2's-complement to 0 |
| | | if rA <= 0 set pc = rB |
| | | else increment pc as normal |

Function Calls

What kinds of things do we put in memory?

- Code: binary code like instructions in our example ISA
 - Intel/AMD compatible: x86_64
 - Apple Mx and Ax, ARM: ARM
 - And others!
- Variables: we may have more variables that will fit in registers
- Data Structures: organized data, collection of data
 - Arrays, lists, heaps, stacks, queues, ...

What if we have many variables? Compute: x += y

Array: a sequence of values (collection of variables)

In Java, arrays have the following properties:

- Fixed number of values
- Not resizable
- All values are the same type

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- Not resizable
- \cdot All values are the same type

How do we store them in memory?



In memory, store array sequentially

- Pick address to store array
- Subsequent elements stored at following addresses
- Access elements with math

Example: Store array arr at **0x90**

• Access arr[3] as **0x90** + **3** assuming 1-byte values

What are we missing?

- Nothing says "this is an array" in memory
- Nothing says how long the array is