

Exam Review

CS 2130: Computer Systems and Organization 1

February 22, 2023

Announcements

- Exam 1 Friday (in class)
 - Closed book, closed notes, closed neighbor, closed internet, closed smart-watch
 - Please bring pen or pencil, we will have scratch paper if needed
 - For SDAC accommodations, please schedule a time with their testing center
- Homework 3 due tonight
- Homework 4 available after the exam

Topics

So far, we have electricity on wire to building a computer and programming it!

- Logic: Operations, Gates, Truth tables
- Numbers: Binary, Octal, Decimal, Hexadecimal
- Bitwise Operations: and, or, bitwise not, logical not, xor, ...
- Binary Arithmetic: addition, subtraction
- Binary Representations: biased integers, two's complement, floating point (8-bit)
- Circuits: adder, subtractor, incrementer, registers, clocks
- High-level how these pieces fit together to form a computer
- Instruction Set Architectures (ISAs) and how to write instructions with our ISA

Today's Agenda

1-bit Logic Gates

- and, or, not
- nand, nor, xor
- Transistors and how to make these gates (high level)

Trinary operator - Mux

- Python: `x = b if a else c`
- Java: `x = a ? b : c`

From our oldest cultures, how do we mark numbers?

- Arabic numerals
 - Positional numbering system
 - The **10** is significant:
 - 10 symbols, using 10 as base of exponent
 - The **10** is *arbitrary*
 - *We can use other bases!* π , 2130, 2, ...

Base-8 Example

Try to turn 134_8 into base-10:

Long Numbers in Binary

Making binary more readable

- Typical to group by 3 or 4 bits
- No need for commas *Why?*
- We can use a separate symbol per group
- How many do we need for groups of 3?
- Turn each group into decimal representation
- Converts binary to **octal**

100001010010

Long Numbers in Binary

Making binary more readable

- Groups of 4 more common
- How many symbols do we need for groups of 4?
- Converts binary to **hexadecimal**
- Base-16 is very common in computing

100001010010

Negative Integers

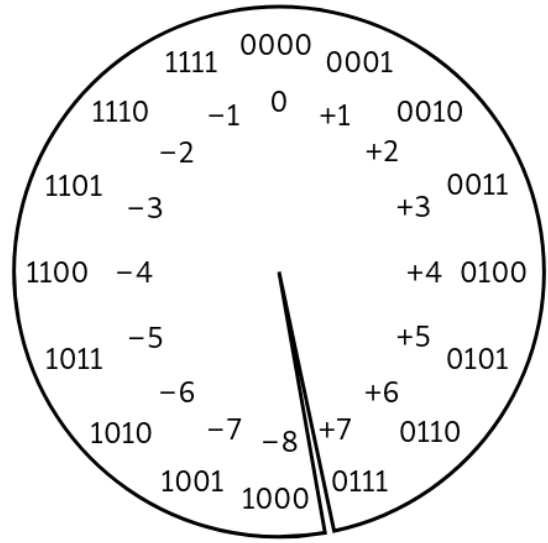
Representing negative integers

- Computers store numbers in fixed number of wires
- Ex: consider 4-digit decimal numbers
- Throw away the last borrow:
 - $0000 - 0001 = 9999$
 - $9999 - 0001 = 9998$
 - Normal subtraction/addition still works
- This works the same in binary

Two's Complement

This scheme is called **Two's Complement**

- More generically, a *signed* integer
- There is a break as far away from 0 as possible
- First bit acts vaguely like a minus sign
- Works as long as we do not pass number too large to represent



Values of Two's Complement Numbers

Consider the following 8-bit two's complement binary number:

11010011

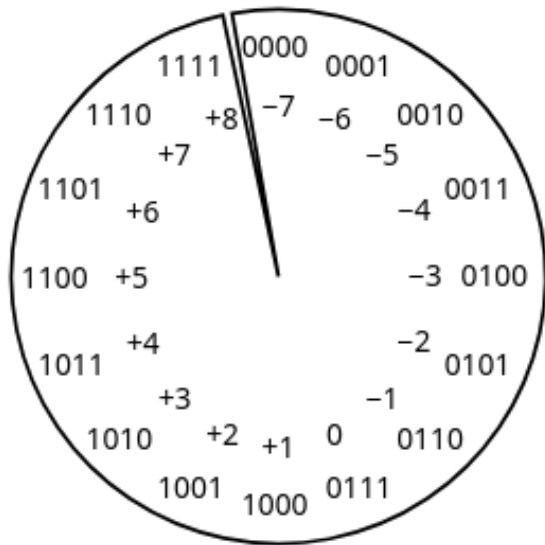
What is its value in decimal?

1. Flip all bits
2. Add 1

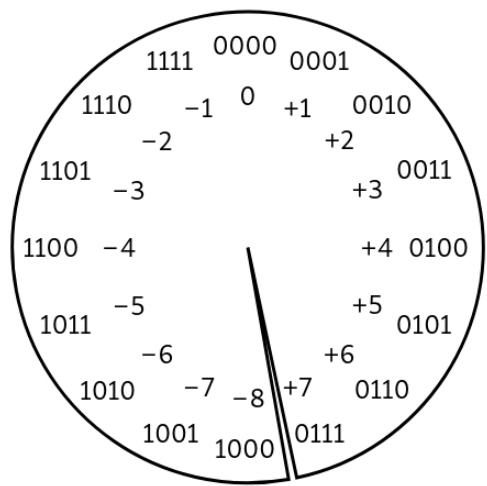
Biased Integers

Similar to Two's Complement, but add **bias**

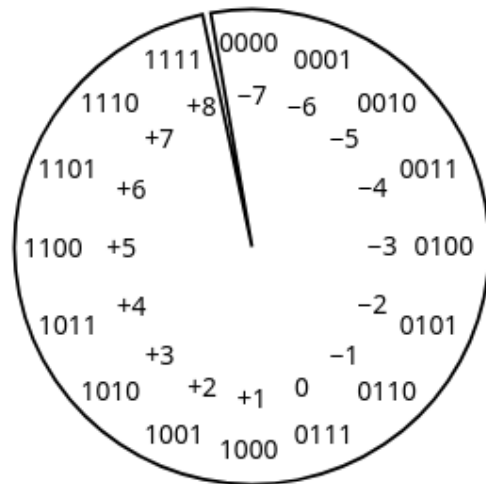
- **Two's Complement:** Define 0 as 00...0
- **Biased:** Define 0 as 0111...1
- Biased wraps from 000...0 to 111...1



Biased Integers



Two's Complement



Biased

Non-Integer Numbers

Floating point numbers

- Decimal: 3.14159
- Binary: 11.10110
- With integers, the point is always fixed after all digits
- With floating point numbers, the point can move!

Challenge! only 2 symbols in binary

Floating Point in Binary

We must store 3 components

- **sign** (1-bit): 1 if negative, 0 if positive
- **fraction** or **mantissa**: (?-bits): bits after binary point
- **exponent** (?-bits): how far to move binary point

We do not need to store the value before the binary point. Why?

Floating Point in Binary

How do we store them?

- Originally many different systems
- IEEE standardized system (IEEE 754 and IEEE 854)
- Agreed-upon order, format, and number of bits for each

$$1.01101 \times 2^5$$

Exponent

How do we store the exponent?

- Exponents *can* be negative

$$2^{-3} = \frac{1}{2^3} = \frac{1}{8}$$

- Need positive and negative ints (but no minus sign)
- *Don't we always use Two's Complement?* **Unfortunately Not**
- **Biased integers**
 - Make comparison operations run more smoothly
 - Hardware more efficient to build
 - Other valid reasons

Floating Point Numbers

Four cases:

- **Normalized:** What we have seen today

$$seeeffff = \pm 1.ffff \times 2^{eeee - \text{bias}}$$

- **Denormalized:** Exponent bits all 0

$$seeeffff = \pm 0.ffff \times 2^{1 - \text{bias}}$$

- **Infinity:** Exponent bits all 1, fraction bits all 0
- **Not a Number (NaN):** Exponent bits all 1, fraction bits not all 0

Operations So Far

So far, we have discussed:

- Addition: $x + y$
 - Can get multiplication
- Subtraction: $x - y$
 - Can get division, but more difficult
- Unary minus (negative): $-x$
 - Flip the bits and add 1

Operations (on Integers)

Bit vector: fixed-length sequence of bits (ex: bits in an integer)

- Manipulated by bitwise operations

Bitwise operations: operate over the bits in a bit vector

- Bitwise not: $\sim x$ - flips all bits (unary)
- Bitwise and: $x \& y$ - set bit to 1 if x, y have 1 in same bit
- Bitwise or: $x | y$ - set bit to 1 if either x or y have 1
- Bitwise xor: $x \wedge y$ - set bit to 1 if x, y bit differs

Example: Bitwise AND

```
    11001010  
& 01111100  
-----
```

Example: Bitwise OR

```
    11001010  
|  01111100  


---


```


Example: Bitwise XOR

$$\begin{array}{r} 11001010 \\ \wedge 01111100 \\ \hline \end{array}$$

Your Turn!

What is: $0x1a \wedge 0x72$

Operations (on Integers)

- Logical not: $!x$
 - $!0 = 1$ and $!x = 0, \forall x \neq 0$
 - Useful in C, no booleans
 - Some languages name this one differently
- Left shift: $x \ll y$ - move bits to the left
 - Effectively multiply by powers of 2
- Right shift: $x \gg y$ - move bits to the right
 - Effectively divide by powers of 2
 - Signed (extend sign bit) vs unsigned (extend 0)

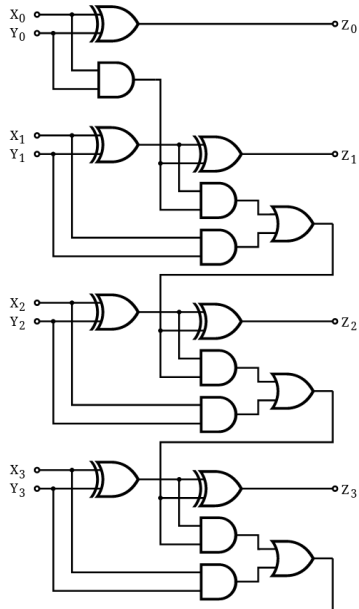
Right Bit-shift Example 2

For **signed** integers, extend the sign bit (1)

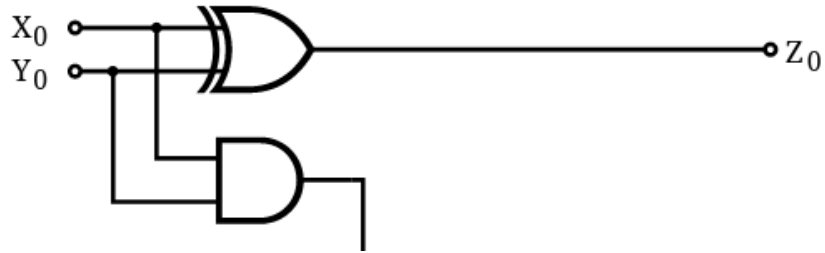
- Keeps negative value (if applicable)
- Approximates divide by powers of 2

11001010 >> 1

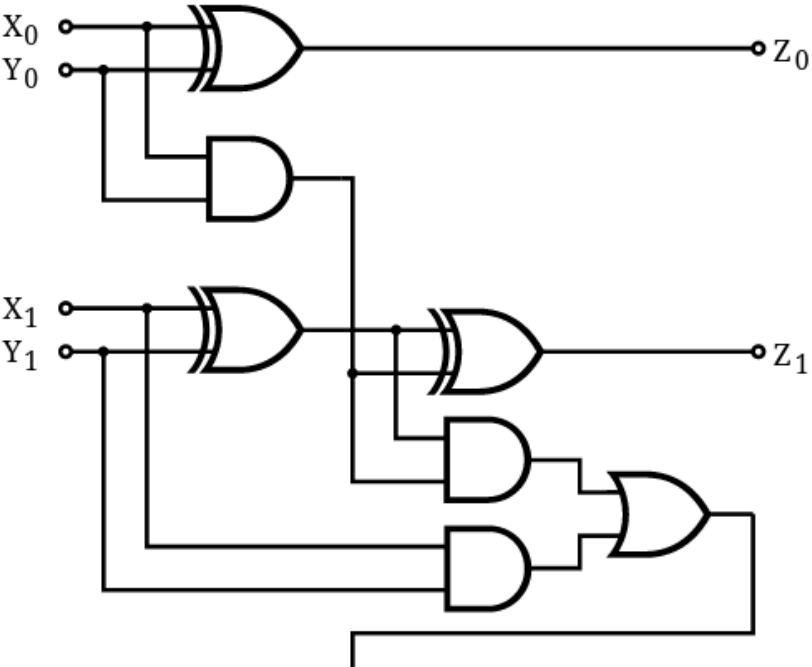
Ripple-Carry Adder



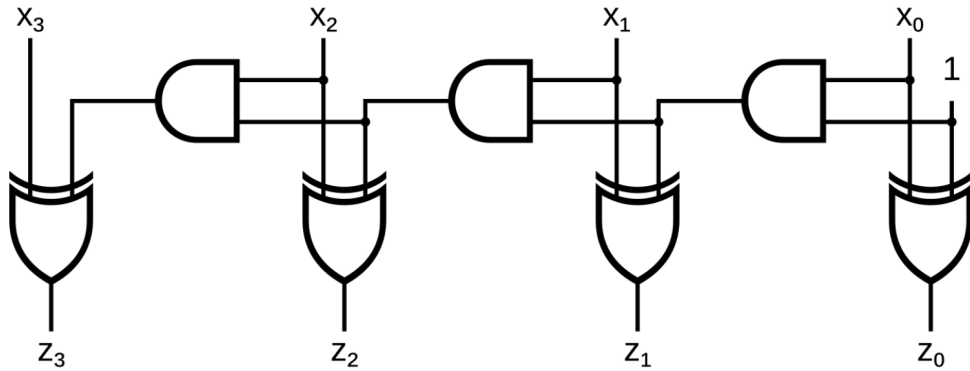
Ripple-Carry Adder



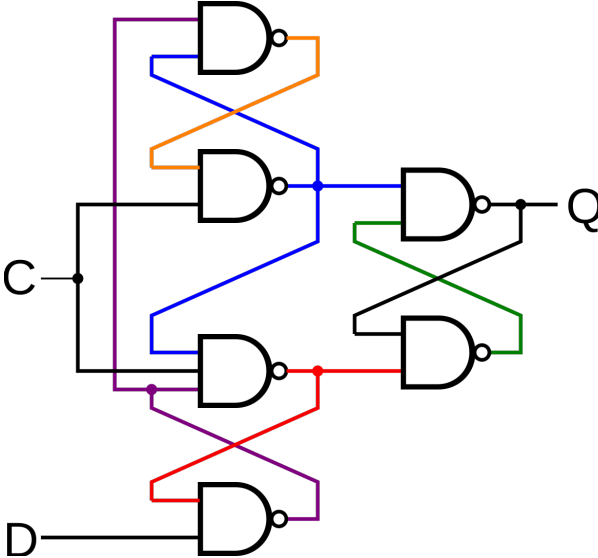
Ripple-Carry Adder



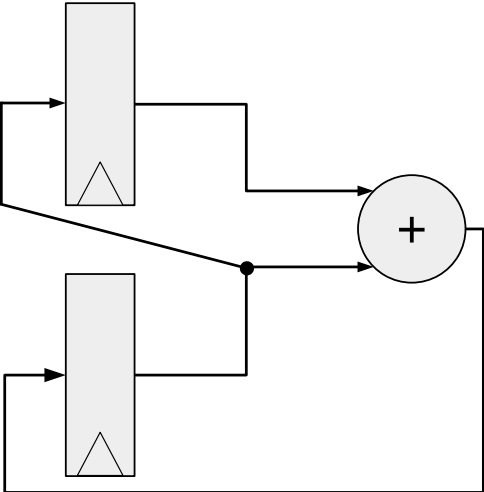
Increment Circuit



1-bit Register Circuit



Another Circuit



Code to Build Circuits from Gates

Write code to build circuits from gates

- Gates we *already* know: $\&$, $|$, \wedge , \sim
- Operations we can build from gates: $+$, $-$
- Others we can build:
- Ternary operator: $?$ $:$

Equals

Equals: =

- Attach with a wire (i.e., connect things)
- Ex: $z = x * y$
- What about the following?

$$x = 1$$

$$x = 0$$

- **Single assignment:** each variable can only be assigned a value once

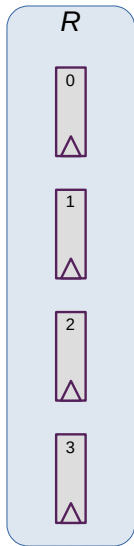
Indexing

Indexing with square brackets: []

- **Register bank** (or **register file**) - an array of registers
 - Can programmatically pick one based on index
 - I.e., can determine which register while running
- Two important operations:
 - $x = R[i]$ - Read from a register
 - $R[j] = y$ - Write to a register

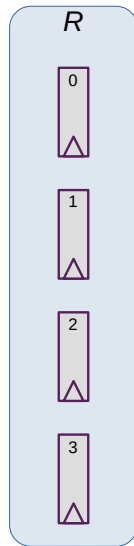
Reading

$x = R[i]$ - connect output of registers to x based on index i



Writing

$R[j] = y$ - connect y to input of registers based on index j



Memory and Storage

Registers

≈ KiB

- 6 gates each, ≈ 24 transistors
- Efficient, fast
- Expensive!
- Ex: local variables

Memory

≈ GiB

- Two main types: SRAM, DRAM
- DRAM: 1 transistor, 1 capacitor per bit
- DRAM is cheaper, simpler to build
- Ex: data structures, local variables

These do not persist between power cycles

Memory and Storage

Disk

≈ GiB-TiB

- Two main types: flash (solid state), magnetic disk
- Magnetic drive
 - Platter with physical arm above and below
 - Cheap to build
 - Very slow! Physically move arm while disk spins

- Ex: files

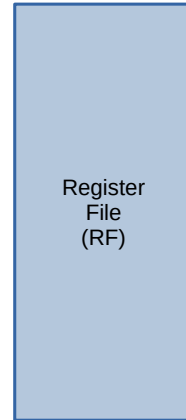
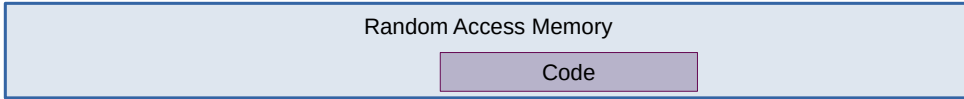
Data on disk does persist between power cycles

Bookkeeping

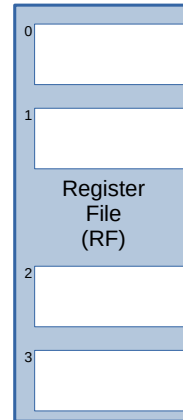
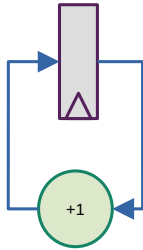
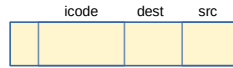
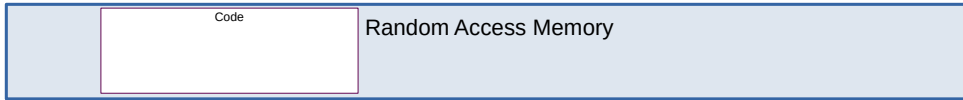
What do we need to keep track of?

- **Code** - the program we are running
 - RAM (Random Access Memory)
- **State** - things that may change value (i.e., variables)
 - Register file - can read and write values each cycle
- **Program Counter (PC)** - where we are in our code
 - Single register - byte number in memory for next instruction

Building a Computer

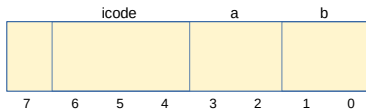


Building a Computer



Our Instruction Set Architecture

icode	b	meaning
0		$rA = rB$
1		$rA += rB$
2		$rA \&= rB$
3		$rA =$ read from memory at address rB
4		write rA to memory at address rB
5	0	$rA = \sim rA$
	1	$rA = -rA$
	2	$rA = !rA$
	3	$rA = pc$
6	0	$rA =$ read from memory at $pc + 1$
	1	$rA +=$ read from memory at $pc + 1$
	2	$rA \&=$ read from memory at $pc + 1$
	3	$rA =$ read from memory at the address stored at $pc + 1$ For icode 6, increase pc by 2 at end of instruction
7		Compare rA as 8-bit 2's-complement to θ if $rA \leq \theta$ set $pc = rB$ else increment pc as normal



High-level Instructions

In general, 3 kinds of instructions

- **moves** - move values around without doing “work”
- **math** - broadly doing “work”
- **jumps** - jump to a new place in the code

Moves

Few forms

- Register to register (icode 0), $x = y$
- Register to/from memory (icodes 3-4), $x = M[b], M[b] = x$

Memory

- **Address:** an index into memory.
 - Addresses are just (large) numbers
 - Usually we will not look at the number and trust it exists and is stored in a register

Broadly doing work

Example 3-bit icode

icode	b	meaning
1		$rA += rB$
2		$rA \&= rB$
5	0	$rA = \sim rA$
	1	$rA = -rA$
	2	$rA = !rA$
6	1	$rA +=$ read from memory at $pc + 1$
	2	$rA \&=$ read from memory at $pc + 1$

Note: I can implement other operations using these things!

Jumps

- Moves and math are large portion of our code
- We also need **control constructs**
 - Change what we are going to do next
 - **if, while, for**, functions, ...
- Jumps provide mechanism to perform these control constructs
- We jump by assigning a new value to the program counter **PC**

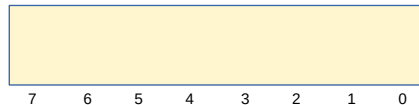
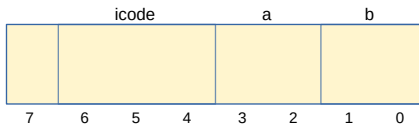
Immediate values

icode 6 provides literals, **immediate** values

Example 3-bit icode

icode	b	action
6	0	$rA = \text{read from memory at } pc + 1$
	1	$rA += \text{read from memory at } pc + 1$
	2	$rA \&= \text{read from memory at } pc + 1$
	3	$rA = \text{read from memory at the address stored at } pc + 1$

For icode 6, increase pc by 2 at end of instruction



Jumps

Example 3-bit icode

icode	meaning
7	Compare rA as 8-bit 2's-complement to θ if $rA \leq \theta$ set $pc = rB$ else increment pc as normal

Instruction icode 7 provides a **conditional** jump

- Real code will also provide an **unconditional** jump, but a conditional jump is sufficient

Writing Code

We can now write any* program!

- When you run code, it is being turned into instructions like ours
- Modern computers use a larger pool of instructions than we have (we will get there)

*we do have some limitations, since we can only represent 8-bit values and some operations may be tedious.

Arrays

Array: a sequence of values (collection of variables)

In Java, arrays have the following properties:

- Fixed number of values
- Not resizable
- All values are the same type

How do we store them in memory?

Storing Arrays

In memory, store array sequentially

- Pick address to store array
- Subsequent elements stored at following addresses
- Access elements with math

Example: Store array *arr* at `0x90`

- Access *arr*[3] as `0x90 + 3` assuming 1-byte values

What's Missing?

What are we missing?

- Nothing says “this is an array” in memory
- Nothing says how long the array is

