

Endianness, Assembly

CS 2130: Computer Systems and Organization 1

March 1, 2023

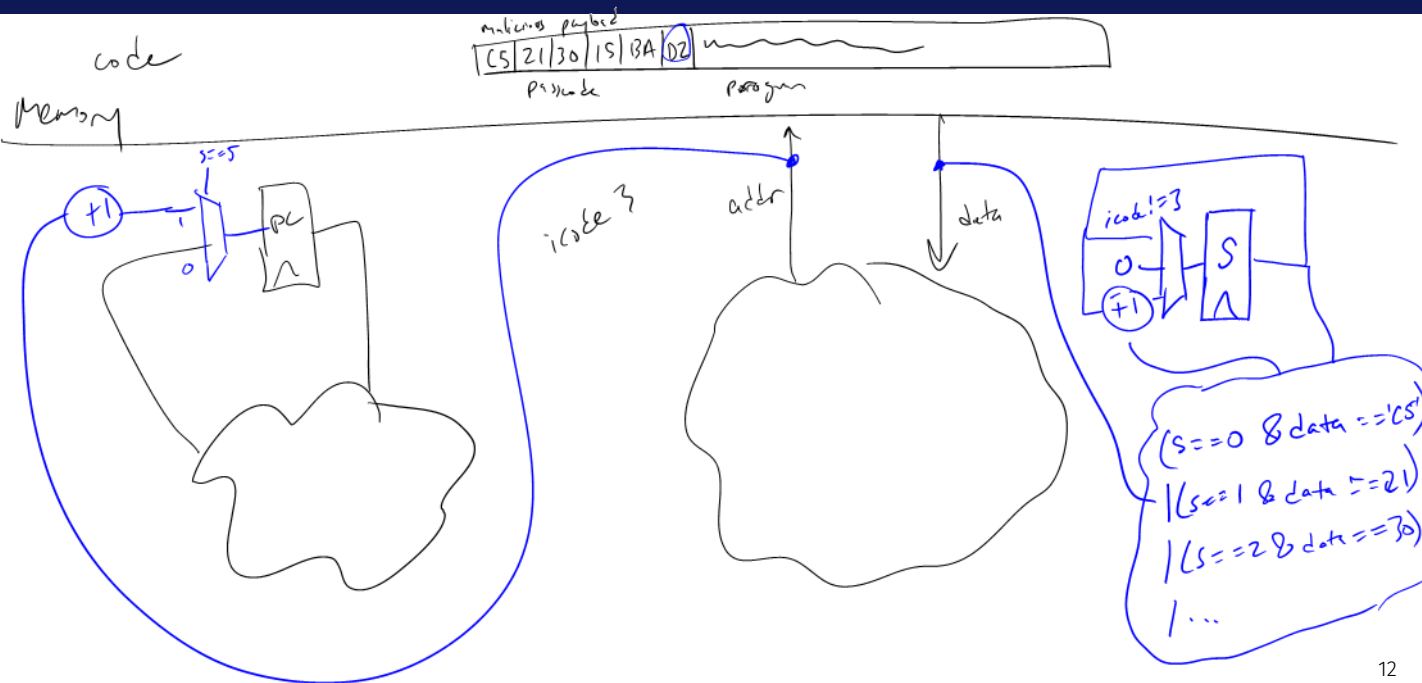
Announcements

- Homework 4 due **Friday** at 11pm on Gradescope
- Exam 1 scores released

Statistics

Mean	75.2
Median	78.0
Std. Dev.	18.66

Our Hardware Backdoor



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- Modern chips have **billions** of transistors
- We're talking adding a few hundred transistors
- *Maybe with a microscope? But you'd need to know where to look!*

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Our Hardware Backdoor

Have you heard about something like this before?

- Sounds like something from the movies
- People claim this might be happening
- To the best of my knowledge, no one has ever *admitted* to falling in this trap

Are there reasons to do this? Not to do this?

- No technical reason not to, it's easy to do!

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Can we make a system where one bad actor can't break it?

- Code reviews, double checks, verification systems, automated verification systems, ...

Why does this work?

Why?

Why does this work?

- **It's all bytes!**
- Everything we store in computers are bytes
- We store code and data in the same place: memory

It's all bytes

Memory, Code, Data... It's all bytes!

- **Enumerate** - pick the meaning for each possible byte
- **Adjacency** - store bigger values together (sequentially)
- **Pointers** - a value treated as address of thing we are interested in

Enumerate

Enumerate - pick the meaning for each possible byte

What is 8-bit 0x54?

Unsigned integer	eighty-four
Signed integer	positive eighty-four
Floating point w/ 4-bit exponent	twelve
ASCII	capital letter T: T
Bitvector sets	The set {2, 3, 5}
Our example ISA	Flip all bits of value in r1

Adjacency

Adjacency - store bigger values together (sequentially)

- An array: build bigger values out of many copies of the same type of small values

$a[i]$

- Store them next to each other in memory
- Arithmetic to find any given value based on index

$$\text{addr} + (i \times \text{size of item})$$

- Records, structures, classes

- Classes have fields! Store them adjacently
- Know how to access (add offsets from base address)
- If you tell me where object is, I can find fields

b, x

$$\text{addr} + \text{offset of } x$$

Pointers

Pointers - a value treated as address of thing we are interested in

- A value that really points to another value
- Easy to describe, hard to use properly
- *We'll be talking about these a lot in this class!*
- Give us strange new powers (represent more complicated things), e.g.,
 - Variable-sized lists
 - Values that we don't know their type without looking
 - Dictionaries, maps

Programs Use These!

How do our programs use these?

- Enumerated icodes, numbers
- Adjacenty stored instructions (PC+1)
- Pointers of where to jump/goto (addresses in memory)

Moving On

icode	b	meaning
0		$rA = rB$
1		$rA += rB$
2		$rA \&= rB$
3		$rA =$ read from memory at address rB
4		write rA to memory at address rB
5	0	$rA = \sim rA$
	1	$rA = -rA$
	2	$rA = !rA$
	3	$rA = pc$
6	0	$rA =$ read from memory at $pc + 1$
	1	$rA +=$ read from memory at $pc + 1$
	2	$rA \&=$ read from memory at $pc + 1$
	3	$rA =$ read from memory at the address stored at $pc + 1$ For icode 6, increase pc by 2 at end of instruction
7		Compare rA as 8-bit 2's-complement to θ if $rA \leq \theta$ set $pc = rB$ else increment pc as normal

So far, we've been dealing with an 8-bit machine!

64-bit Machines

64-bit machine: The **registers** are 64-bits

- i.e., r0, but also PC

Important to have large values. Why?

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4 GiB

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- Today's processors - 64 bits: 2^{64} addresses

Aside: Powers of Two

Powers of Two

Value	base-10	Short form	Pronounced
2^{10}	1024	Ki	Kilo
2^{20}	1,048,576	Mi	Mega
2^{30}	1,073,741,824	Gi	Giga
2^{40}	1,099,511,627,776	Ti	Tera
2^{50}	1,125,899,906,842,624	Pi	Peta
2^{60}	<u>1,152,921,504,606,846,976</u>	Ei	Exa

Example: 2^{27} bytes

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Example: 2^{27} bytes = $2^7 \times 2^{20}$ bytes = 2^7 MiB = 128 MiB

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- 16 EiB (2^{64} addresses = $2^4 \times 2^{60}$)

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- 16 EiB (2^{64} addresses = $2^4 \times 2^{60}$)
- But I only have 8 GiB of RAM

A Challenge

There is a disconnect:

- Registers: 64-bits values
- Memory: 8-bit values (i.e., **1 byte** values)
 - Each address addresses an 8-bit value in memory
 - Each address points to a 1-byte slot in memory

A Challenge

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- Registers: 64-bits values
- Memory: 8-bit values (i.e., **1 byte** values)
 - Each address addresses an 8-bit value in memory
 - Each address points to a 1-byte slot in memory
- How do we store a 64-bit value in an 8-bit spot?