Bitwise Operations, Floating Point Numbers

CS 2130: Computer Systems and Organization 1 January 27, 2023

- TA Office Hours starting very soon
- Discord link coming this afternoon
- Quiz 1 opens this afternoon, due Sunday night
- Homework 1 due Feb 6 (Mon)

Two's Complement

Two's Complement

This scheme is called Two's Complement

- More generically, a *signed* integer
- There is a break as far away from 0 as possible
- First bit acts vaguely like a minus sign
- Works as long as we do not pass number too large to represent



Two's Complement

Consider the following 8-bit binary number in Two's Complement:

What is its value in decimal?

$$\int_{3} \frac{1111111}{0010011} = -45$$

$$\int_{3} \frac{100001}{0010100} = -45$$

$$\frac{1}{32}$$

$$\frac{1}{3263421}$$

Consider the following 8-bit binary number in Two's Complement:

11010011



What is its value in decimal?

- 1. Flip all bits
- 2. Add 1



So far, we have discussed:

- Addition: x + y
 - Can get multiplication
- Subtraction: x y
 - $\cdot\,$ Can get division, but more difficult
- Unary minus (negative): -x
 - + Flip the bits and add 1

Bit vector: fixed-length sequence of bits (ex: bits in an integer)

• Manipulated by bitwise operations

Bitwise operations: operate over the bits in a bit vector

- Bitwise not: ~x flips all bits (unary)
- Bitwise and: $\mathbf{x} \cdot \mathbf{\delta} \mathbf{y}$ set bit to 1 if x, y have 1 in same bit
- Bitwise or: $x \mid y$ set bit to 1 if either x or y have 1
- Bitwise xor: $\mathbf{x} \circ \mathbf{y}$ set bit to 1 if x, y bit differs

Example: Bitwise AND

11001010 × 84 × 8 01111100 Y 01001000

Example: Bitwise OR

11001010 × × / y | 01111100 y | | | | | | 0

Example: Bitwise XOR

11001010 ^ 01111100 | 0 | | 0 | | 0

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Your Turn!

What is:
$$0x1a ^{0}x72$$

 $0x1a = 0001100$
 $0x72 = 01110010$
 $010100 = 0x68$

Operations (on Integers)

- Logical not: **!x**
 - $!0 = 1 \text{ and } !x = 0, \forall x \neq 0$
 - Useful in C, no booleans
 - \cdot Some languages name this one differently
- Left shift: x << y move bits to the left
 - Effectively multiply by powers of 2
- Right shift: x >> y move bits to the right
 - Effectively divide by powers of 2
 - Signed (extend sign bit) vs unsigned (extend 0)

Left Bit-shift Example

01011010 << 2 01101000

Right Bit-shift Example



Bit-shift

Computing bit-shift effectively multiplies/divides by powers of 2

Consider decimal:

$$2130 <<_{10} 2 = 213000 = 2130 \times 100$$

 $2130 >>_{10} 1 = 213 = 2130 / 10$

Right Bit-shift Example 2



Right Bit-shift Example 2

For signed integers, extend the sign bit (1)

- Keeps negative value (if applicable)
- Approximates divide by powers of 2

110110 -54 ~ 11001010 >> 1

$$-27 \leftarrow 1||00|0|$$

11001010 >>>3 []] []00]

Bit fiddling example

What about other kinds of numbers?

• Decimal: 3.14159

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- Binary: 11.10110

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- With floating point numbers, the point can move!

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Challenge! only 2 symbols in binary

Convert the following decimal to scientific notation:

2130

Convert the following binary to scientific notation:

101101

An interesting phenomenon:

• Decimal: first digit can be any number *except* 0

 2.13×10^3

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• Decimal: first digit can be any number *except* 0

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• Binary: first digit can be any number *except* 0 Wait!

 1.01101×2^{5}

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• Decimal: first digit can be any number *except* 0

2.13×10^3

• Binary: first digit can be any number *except* 0 Wait!

 1.01101×2^5

• First digit can only be 1

We must store 3 components

- sign (1-bit): 1 if negative, 0 if positive
- fraction or mantissa: (?-bits): bits after binary point
- exponent (?-bits): how far to move binary point

We do not need to store the value before the binary point. Why?

How do we store them?

- Originally many different systems
- IEEE standardized system (IEEE 754 and IEEE 854)
- Agreed-upon order, format, and number of bits for each

1.01101×2^{5}



A rough example in Decimal:

6.42×10^{3}

Exponent

How do we store the exponent?

• Exponents *can* be negative

$$2^{-3} = \frac{1}{2^3} = \frac{1}{8}$$

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- Biased integers
 - Make comparison operations run more smoothly
 - Hardware more efficient to build
 - Other valid reasons

Biased Integers

Similar to Two's Complement, but add bias

- Two's Complement: Define 0 as 00...0
- Biased: Define 0 as 0111...1
- Biased wraps from 000...0 to 111...1

0000 0001 1110 0010 +8 -6 +7 1101 0011 +6-3 0100 1100 +5-2 +41011 0101 +3 1010 0110 1001 0111 1000

Biased Integers

0000 1111 0001 0 1110 0010 -1 +1 +2 -2 1101 0011 +3 -3 +4 0100 1100 -4 +5 -5 1011 0101 +6 -6 1010 -7 +7 0110 -8 1001 1000 0111

Two's Complement

0000 0001 1111 -7 1110 0010 +8 -6 +7 -5 1101 0011 -4 +6 1100 -3 0100 +5 -2 +4 1011 0101 +3 -1 +2 1010 0 0110 +11000 0111 1001 Biased

Calculate value of biased integers (4-bit example)

0010

Biased Integers

101.011₂

101.011₂

What does the following encode?



What does the following encode?



What about 0?

Four cases:

• Normalized: What we have seen today

s eeee ffff = $\pm 1.ffff \times 2^{eeee-bias}$

• **Denormalized**: Exponent bits all 0

s eeee ffff =
$$\pm 0.ffff \times 2^{1-bias}$$

- Infinity: Exponent bits all 1, fraction bits all 0 (i.e., $\pm\infty$)
- Not a Number (NaN): Exponent bits all 1, fraction bits not all 0