Pipeline Simulator Exercise #2
SimpleScalar Exercise #3

Due Friday, September 28, 2001, 4:30 p.m.

This exercise is meant to extend your how pipelines function from the previous exercise and to further familiarize you with the SimpleScalar 3.0 toolset. In this phase, you will add bimodal branch prediction, single level caches and data forwarding to your pipeline. Your simulator will appropriately handle branch mispredictions and cache misses. There are 2 parts: a programming assignment and a few written exercises.

Assumptions

For the purposes of this exercise, we are assuming the following:

- Single level caches
  - Instruction cache
    - Probed by the fetch unit
    - Miss latency: determined by the cache access function
  - Data cache
    - Probed by the MEM stage if executing a load or store
    - Miss latency: determined by the cache access function
    - Infinite write buffer, so no stalls necessary for writes to memory

- Bimodal branch predictor
  - Branch predictor will be probed on a branch in the decode stage and provide IF with a PC to fetch
  - When the branch reaches the EX, determine whether the prediction was correct, if it’s incorrect, squash the wrongly fetched instructions, fetch the correct instructions

- Data forwarding:
  - If there is a data hazard, but the result is available in the EX/MEM latch or the MEM/WB latch, no stalling should occur

- Split-phase register access (writes occur in first half of clock cycle, reads in second half)

Simulator Skeleton Code

0. Download the code distribution assign3.tar.gz from: ~cs654/fall2001/assign3
You MUST use the code distributed here as your baseline for starting the assignment. Please rename the solution file to sim-pipe2.c, so as not to interfere with the grading of pipeline assignment part 1. (This means you’ll need to edit the Makefile again).

1. Include sim-pipe.h, bpred.h and cache.h to your simulator. This will give you access to the functions in the branch prediction and cache modules. The contents of sim-pipe.h are an assortment of stuff needed in order to do branch prediction and figure out cache miss latency.

2. You’ll have to register options (sim_reg_options) and check them (sim_check_options). For hints on how to do this, look to sim-bpred.c and sim-cache.c

3. More detailed discussion of code specifics will happen at the recitation: Monday, September 24, 2001, 6:30-7:30 p.m., MEC 339

Sample Output

Will be provided soon. Watch your email.

The Programming Assignment

1. Implement data forwarding/bypassing according to the assumptions given above
2. Add instruction cache, data cache, and branch predictor to the pipeline according to the assumptions given above
3. Your simulator should do the following:
   - Detect data and control hazards
   - Stall as appropriate, forward data as appropriate
   - Detect instruction cache misses and data cache misses and appropriately stall when a miss occurs for the latency required to get the instruction from memory
   - Predict branches in the decode stage. Use the prediction to tell fetch the next PC to fetch
   - Count the number of branches executed
   - Calculate the correct prediction rate (number of correct predictions / total number of branches)

The Written Assignment

1. Let us consider a single-issue, in-order system like the DLX pipeline in Chapter 3, except our pipeline only has four stages: IF, ID, EX, and WB. Loads compute their effective address and access the data cache in the EX stage. Let’s assume that all operations take 1 cycle to execute, except for loads. Without a cache, loads take 10 cycles (including the cycle spent in EX to compute the effective address). Loads represent 20% of the dynamic instruction mix.
a.) (5) What is the maximum speedup that we can attain by minimizing the load latency? Assume that, no matter what improvements are made, a load cannot take less than 1 cycle to execute (i.e., it completes at the end of the EX stage).

b.) (10) Suppose that 15% of the dynamic instruction mix consists of arithmetic operations that directly use the loaded values, and that the compiler can convert these sequences into arithmetic operations with a memory operand. These arithmetic-memory operations are able to compute the effective address, access the data cache, and perform the arithmetic operation in the EX stage. Since 20% of the instruction mix is loads, this means that we can convert 75% of the loads. Further suppose that this system now has a cache with a 95% hit rate (i.e., 95% of loads complete at the end of the MEM stage, while 5% still take 10 cycles). The ratio of hits to misses is the same for arithmetic-memory operations and for remaining load operations. The hit rate is the same for the base system and for the new system that uses arithmetic-memory operations. Finally, assume that arithmetic-memory operations complete at the end of the MEM stage. Unfortunately, accommodating the arithmetic-memory operations increases the cycle time by 10%.

Is this change worthwhile? Which system is faster and by how much?

2. Consider the following code sequence:
   lw R1 ← 0(R2)
   add R4 ← R1 + R3
   add R1 ← R1 + R4

   a.) (5) Identify all dependences.
   b.) (5) Using a pipeline diagram (format of your choice), show this code sequence progressing through a five-stage pipeline (IF, ID, EX, MEM, WB) with no bypassing but with split-phase register access (i.e., a register can be written in the first half of the cycle, and the result can then be read in the second half of the cycle). Loads complete at the end of the MEM phase; adds take only one cycle to complete. Registers are read in the ID stage.

Homework Submission

For the programming part:
( NOTE: We will be compiling your simulators and running our test cases on them, so make sure that your simplesim3.0 directory has group rx permissions set)

In an email to cs654@cs.virginia.edu, include the following information:

1) Group member names and email ids (no email aliases please).
2) README: This should detail how to compile your simulator and the names of files you have modified.
3) Pointer to the location of the files you have modified to write your simulator. Ideally this is something like /home/userid/simplesim3.0

**For the written part:**

Follow the guidelines for labelling your assignment described in the BOC memo and submit your homework to the brown folder in Olsson 204 (Brenda Perkins’ office). The folder is near the printer. *Please staple your assignment if it is more than one page.*

**NOTES:**

Recitation, Monday, September 24, 2001, 6:30-7:30 p.m.